



Preprocessing Mobility Data



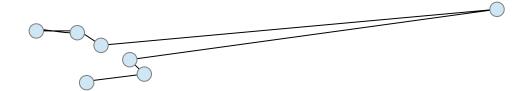
Content of this lesson

- Preprocessing trajectories
 - trajectory filtering
 - point map matching
 - route reconstruction
 - trajectory compression
 - o Semantic enrichment
 - stop detection / trajectory segmentation
 - home location detection (GPS & MobPhones)
 - activity recognition (POI-based)

Trajectory filtering

- Data points are sometimes affected by errors
- Errors can have huge effects on results

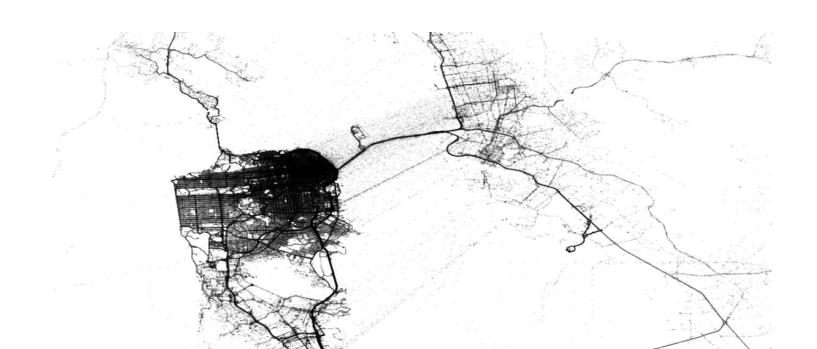
What is the real length of this trip?



- Two families of approaches:
 - Context-based filtering
 - Movement-based filtering

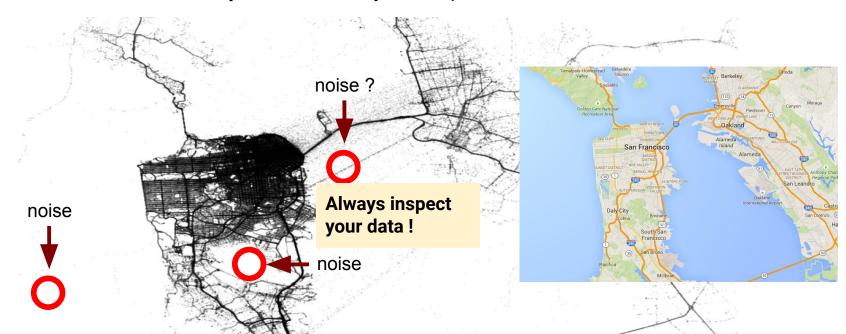
Context-based filtering

• Single points might contain errors of various kinds



Context-based filtering

- Single points might contain errors of various kinds
- Map-based detection: cars on the water or out of roads are noise
 - Caution: do you trust 100% your map?



Movement-based filtering

- No context is used, just the geometry / dynamics of movement
- Speed-based noise filtering approach:
 - The first point of the trajectory is set as valid
 - Scan all remaining points "p" of the trajectory (time order)
 - Compute "v" = average straight-line speed between point "p" and the previous valid one

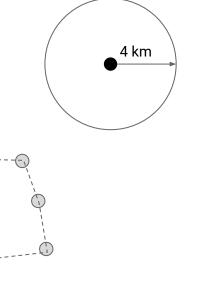


Movement-based filtering

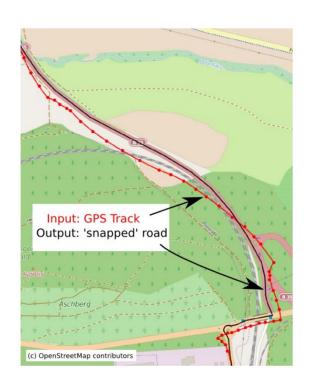
Exercise

start

- What happens in this situation? (Multiple noisy points)
 - Assume constant sampling rate 1 minute
 - Speed threshold = 240 km/h (= 4 km/minute)

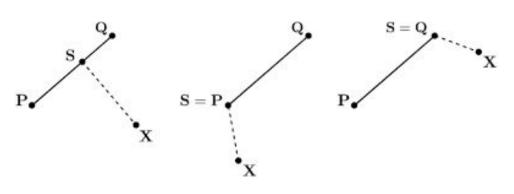


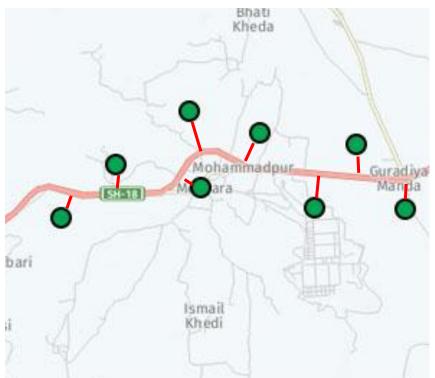
- Points can be aligned to the road network
 - Objective 1: improve accuracy of position
 - Objective 2: remove extreme cases (ref. filtering)
 - Objective 3: translate trajectories to sequences of road IDs
- Idea: project the point to the close location in the network
 - Usually there is a maximum threshold
 - Points farther than the threshold from any road are removed as noise



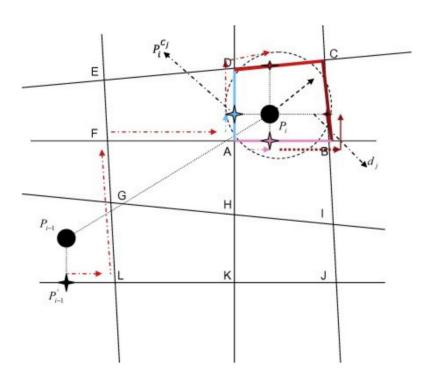
- Point projection
 - Requires to compare each point to each road segment

Refresher on point-to-segment distance computation



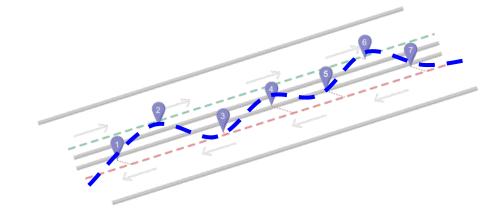


- In some contexts there can be multiple choices at reasonable distance
 - Simply taking the closest one is "risky"

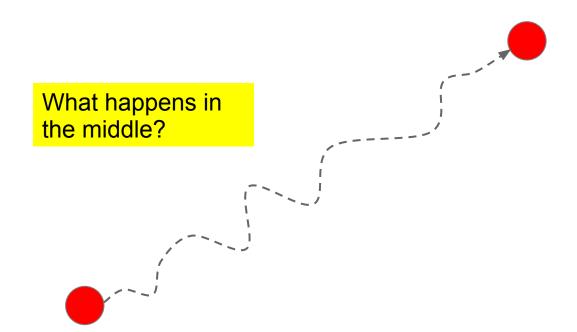


- Matching points separately can lead to inconsistent results
 - Mainly road-dense areas with position accuracy comparable to road separation
- Need a trajectory-level matching
 - Linked to route reconstruction

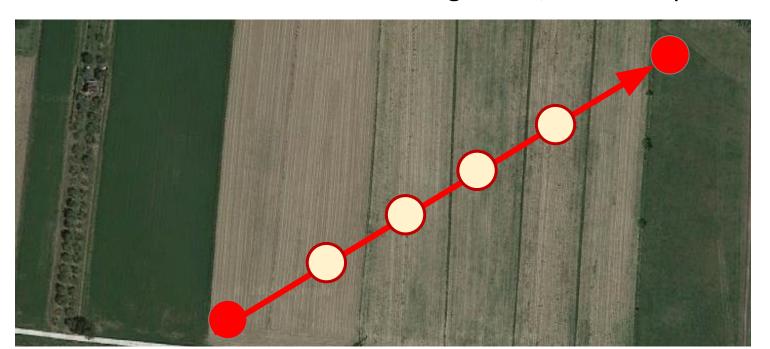




Sometimes the space/time gap between consecutive points is significant



- Typical solutions:
 - Free movement => straight line, uniform speed





- Typical solutions:
 - Constrained movement => shortest path



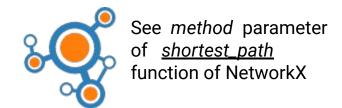


Shortest paths can be replaced by alternative "optimal paths"

- Based on a notion of path cost
- Typical ones: path **length**, path **duration** (requires to know typical traversal times of roads)
- Alternative ones: fuel consumption, EV battery consumption, CO2 emissions, mixed costs

Algorithms applied are standard graph path optimization methods:

- ullet Dijkstra's algorithm o efficient, requires that costs are non-negative
- Bellman-Ford algorithm → less efficient, can work with negative weights (but no cycles)

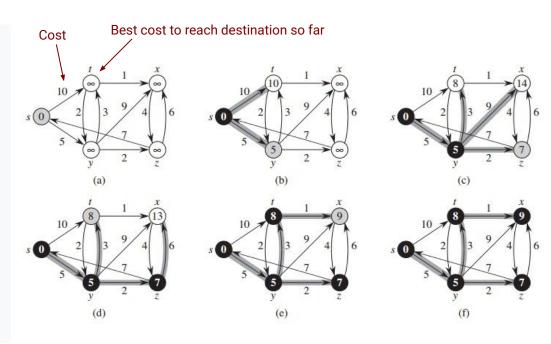


Refresher: Dijkstra's minimum cost algorithm

• Simple and efficient: O(M + N log N) time complexity

(M = |edges|, N = |nodes|)

```
function Dijkstra(Graph, source):
          for each vertex v in Graph.Vertices:
               dist[v] \leftarrow INFINITY
               prev[v] \leftarrow UNDEFINED
               add v to 0
          dist[source] \leftarrow 0
          while Q is not empty:
10
               u \leftarrow \text{vertex in } Q \text{ with min dist[u]}
11
               remove u from Q
12
13
               for each neighbor v of u still in Q:
14
                    alt \leftarrow dist[u] + Graph.Edges(u, v)
15
                    if alt < dist[v]:</pre>
16
                         dist[v] \leftarrow alt
17
                         prev[v] \leftarrow u
18
19
          return dist[], prev[]
```



Trajectory Map Matching

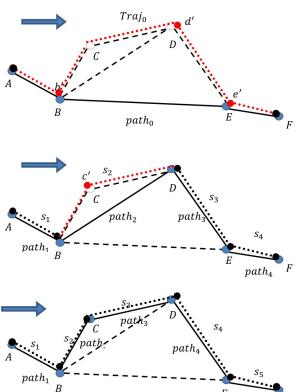
- Assigns points to road segments
- Reconstructs the movement between consecutive points
- Ensures coherence of the overall process

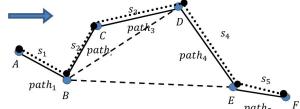
- Two sample approaches:
 - Based on shortest path
 - Based on probabilities

Shortest path-based Map Matching

Used by MappyMatch

- Similar ideas as trajectory simplification
 - Match first and last point
 - Compute shortest path on the network
 - Find farthest point from shortest path
 - If distance > threshold \Rightarrow
 - split into two parts
 - run recursively the process on both



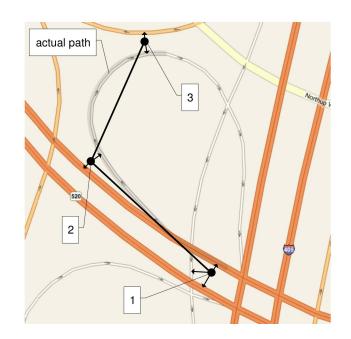


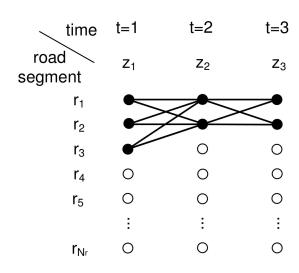
Reference: Zhu, honda & Gonder. A Trajectory Segmentation Map Matching Approach for Large-Scale, High-Resolution GPS Data. TRB 2017.

Probability-based Map Matching

Used by <u>pyTrack</u>

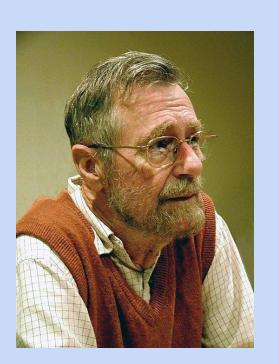
- Consider possible point-to-road assignments, with probabilities
- Compute most likely path that visits all points in the correct sequence





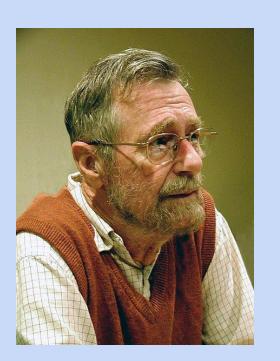
Reference: <u>Newson & Krumm. Hidden Markov Map Matching</u>
<u>Through Noise and Sparseness. ACM GIS'09.</u>

Who's Dijkstra



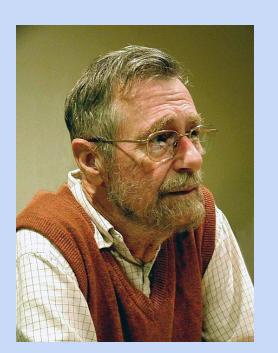
- 1930 2002
- Dutch computer scientist, programmer, software engineer, systems scientist, and science essayist
- 1972 Turing Award for "fundamental contributions to developing programming languages"

Dijkstra is famous for...



- Dijkstra's algorithm, of course
- Contributions to "self-stabilization of program computation"
 - Won him the "ACM PODC Influential Paper Award", later renamed "Dijkstra Prize"
- Hundreds of papers on computational and science philosophy issues

Dijkstra is famous for...



- His habit of writing everything with paper
 & fountain pen
- Hundreds of papers, many unpublished
 - E. W. Dijkstra Archive
- Counting should start from 0, not 1...

When dealing with a sequence of length N, the elements of which we wish to distinguish by subscript, the next vexing question is what subscript value to assign to its starting element. Adhering to convention a) yields, when starting with subscript 1, the subscript range 15 i < N+1; starting with 0, however, gives the nicer range 05 i < N. So let us let our ordinals start at zero: an element's ordinal (subscript) equals the number of elements preceding it in the sequence. And the moral of the story is that we had better regard -after all those centuries! - zero as a most natural number.

Dijkstra the teacher



- Chalk & blackboard, no projectors
- No textbooks
- Improvisation & long pauses
- No references in papers

"For the absence of a bibliography I offer neither explanation nor apology."

- Long exams
 - Each student was examined in Dijkstra's office or home, and an exam lasted several hours

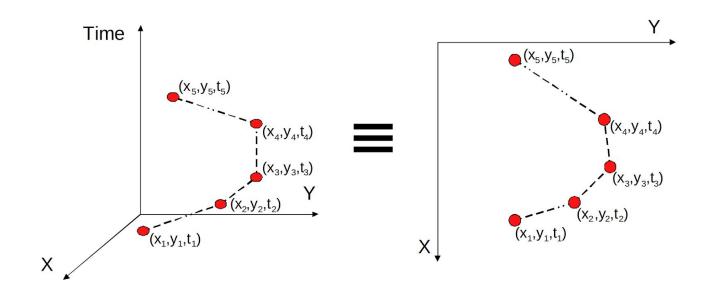
Trajectory compression / simplification

- Many algorithms for trajectories are expensive
 - Their complexity depends on the number of points
 - Sometimes trajectories have more points than needed

- Objective of compression / simplification
 - Reduce the number of points...
 - ... without affecting the quality of results

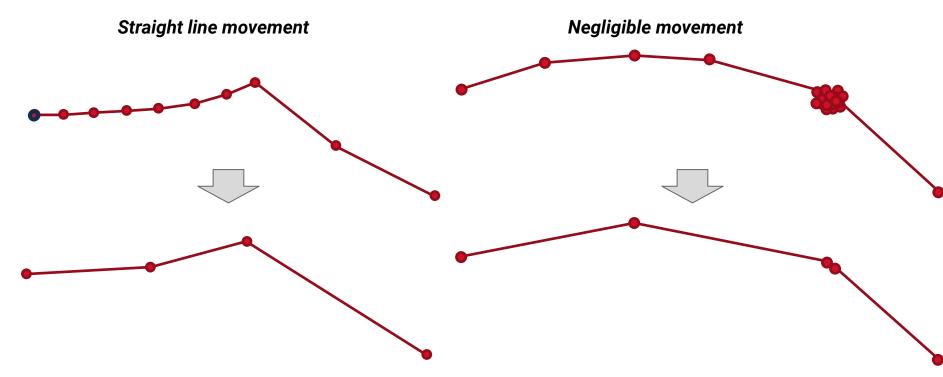
Trajectory data

- A trajectory is a temporal sequence of time-stamped locations
- Most methods focus on the spatial component



Trajectory compression / simplification

• Typical cases where points might be removed



Compression/simplification methods

Some standard methods for simplifying polygonal curves:

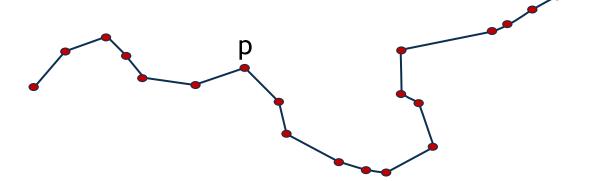
- Ramer-Douglas-Peucker, 1973
- Driemel-HarPeled-Wenk, 2010

1972 by Urs Ramer and 1973 by David Douglas and Thomas Peucker

Mostly known as Douglas-Peucker (DP) algorithm

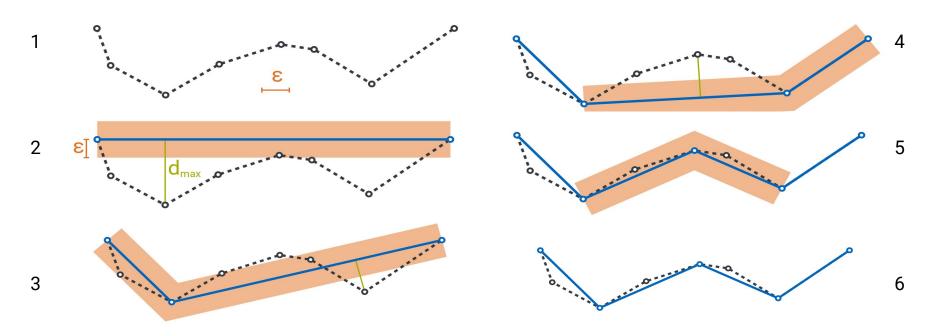
The most successful simplification algorithm. Used in GIS, geography, computer vision, pattern recognition...

Very easy to implement and works well in practice.



```
Input polygonal path P = \langle p_1, ..., p_n \rangle and threshold \varepsilon
Initially i=1 and j=n
Algorithm DP(P,i,j)
    Find the vertex v_f between p_i and p_j farthest from p_i p_j.
    dist := the distance between v_f and p_i p_i.
    if dist > \epsilon then
           DP(P, V_i, V_f)
            DP(P, v_f, v_i)
    else
            Output(v<sub>i</sub>v<sub>i</sub>)
```

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                                                                                dist
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```



Time complexity?

Testing a shortcut between p_i and p_i takes O(j-i) time.

Worst-case recursion?

```
DP(P, v_{i}, v_{i+1})

DP(P, v_{i+1}, v_{i})
```

Time complexity $T(n) = O(n) + T(n-1) = O(n^2)$

But rather fast in practice

```
Algorithm DP(P,i,j)

Find the vertex v_f farthest from p_i p_j.

dist := the distance between v_f and p_i p_j.

if dist > \epsilon then
DP(P, v_i, v_f)
DP(P, v_f, v_j)
else
Output(v_i v_i)
```

```
Simple simplification(P = \langle p_1,...,p_n \rangle, \varepsilon)
```

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i:=1
while i<n do
      q := p_i
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      if no such vertex then set i:=n
      add p<sub>i</sub> to P'
end
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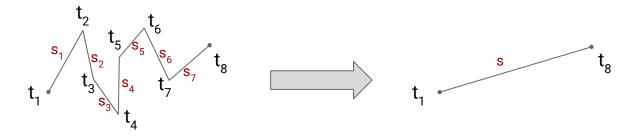
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```

Impact on speed

- What about time and speeds?
 - Time-stamps were never considered in the algorithms
 - They considered on impact on space / geometry of trajectories
 - What impact on time-related aspects, e.g. speed?

Typically, simplification reduces average speed estimates:

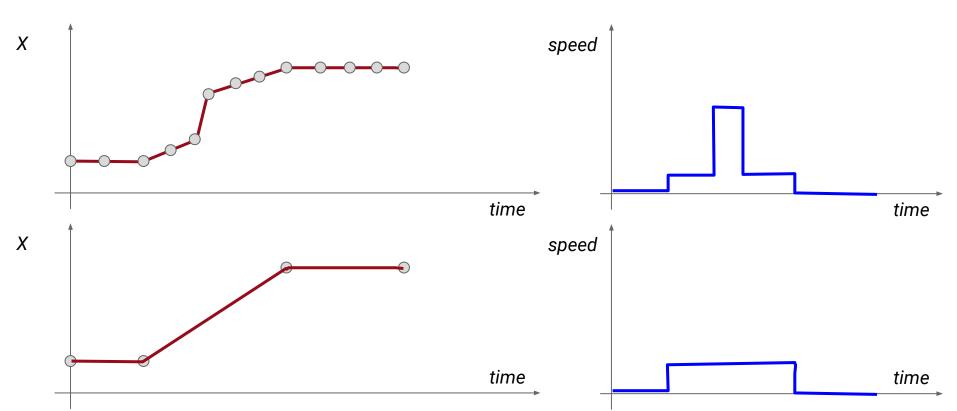


Avg speed =
$$(s_1 + ... + s_7) / (t_8 - t_1)$$

 Avg speed = s / $(t_8 - t_1)$

Impact on speed

Flattening effect



How fast is a cow?

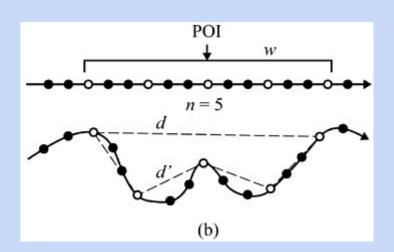


How fast is a cow?

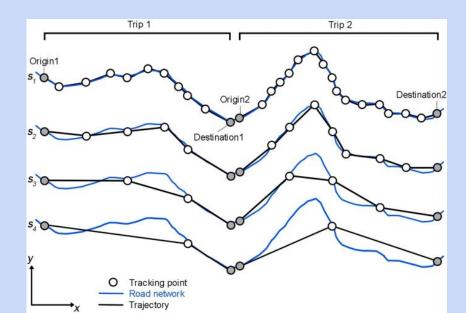
- Trajectory compression / simplification changes the scale of the analysis
 - Simplified data → macroscopic analysis
 - Detailed data → microscopic analysis
- Several movement characteristics can be affected

How fast is a cow?

How fast is a cow? Cross-Scale Analysis of Movement Data Laube P, Purves RS (2011)

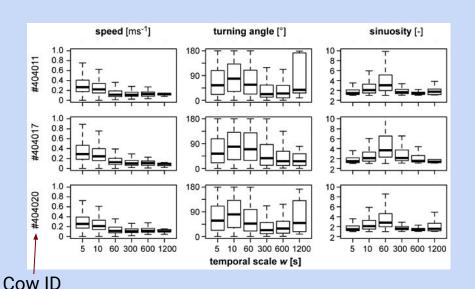


Understanding the impact of temporal scale on human movement analytics Su, R., Dodge, S. & Goulias, K.G (2022)



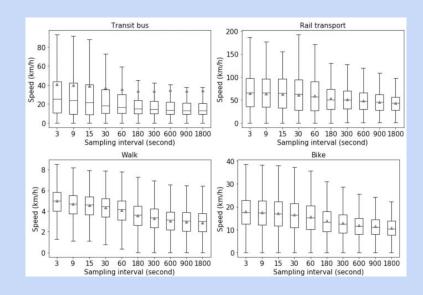
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Understanding the impact of temporal scale on human movement analytics

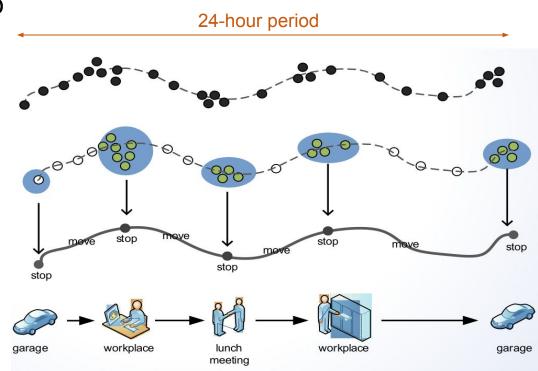
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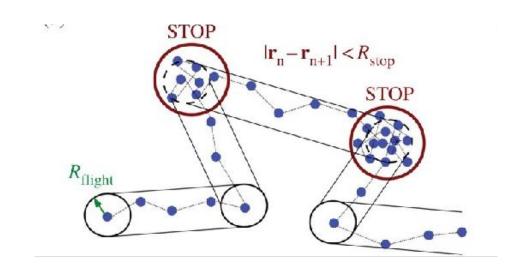
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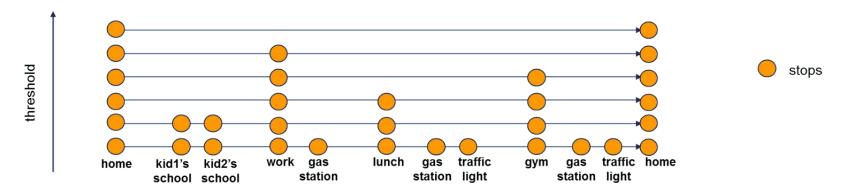
- Raw data forms a continuous stream of points
- Typical unit of analysis: the trip
- How to segment?
 - Basic idea: identify stops



- General criteria based on speed
 - If it **moves very little** (threshold Th_s) over a significant **time interval** (threshold Th_T)
 - => it is practically a stop
 - Trajectory (trip) = contiguous sequence
 of points between two stops
 - Typical values:
 - Th_s within [50, 250] meters
 - o Th_⊤ within [1, 20] minutes



Different time thresholds yield different semantics

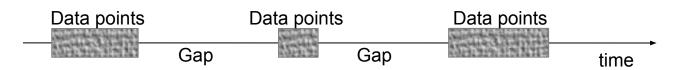


- Which one is the best for you?
 - Application dependent

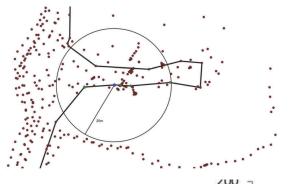
- Special cases, easier to treat
 - Stop explicitly in the data: e.g. engine status on/off
 - Simply "cut" trajectories on status transitions



- Device is off during stops:
 - Typical of cars data
 - A stop results in a time gap in the data
 - Exceptions: short stops might remain undetected

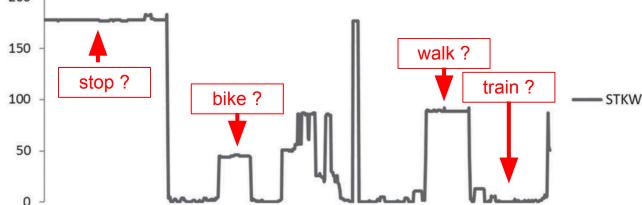


Generalization: transportation means segmentation



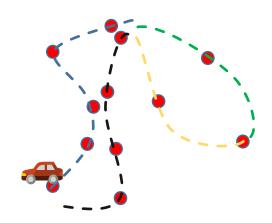
- Speed / density-based approach
- Idea: faster means less of my points around me

Number of points within radius R



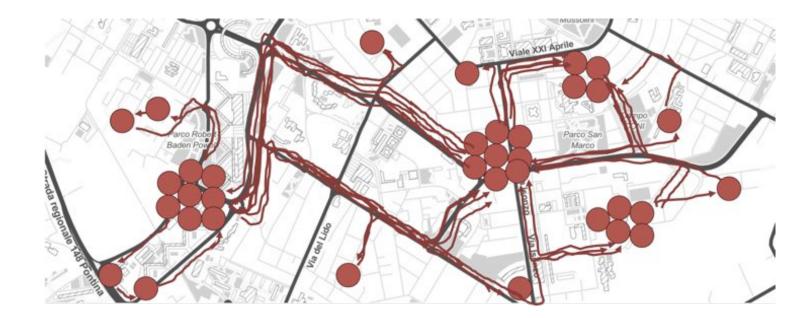
User's Mobility History

- What do we get after segmentation?
- Several trajectories associated to the same subject
- Enables individual-level analyses
 - E.g. explore user's habits, find deviations from usual, etc.



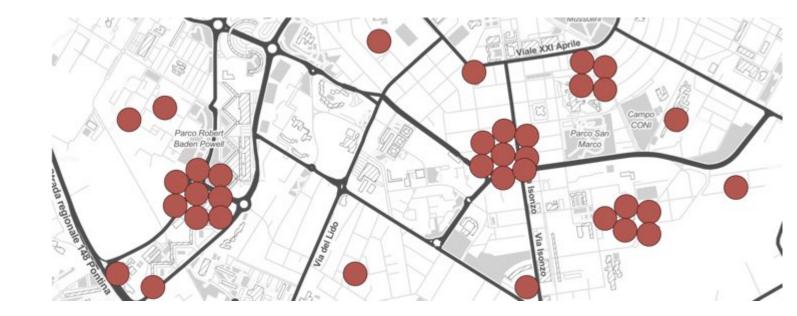
Inferring Home / Work locations

- Take all trips of a vehicle / user
- Build a "Individual Mobility Network"
 - o Graph abstraction of the overall mobility based on locations (nodes) and movements (edges).





- Focus on start and stop points
 - Dense areas represent important places





Cluster points to identify locations



Viale XXI Aprile



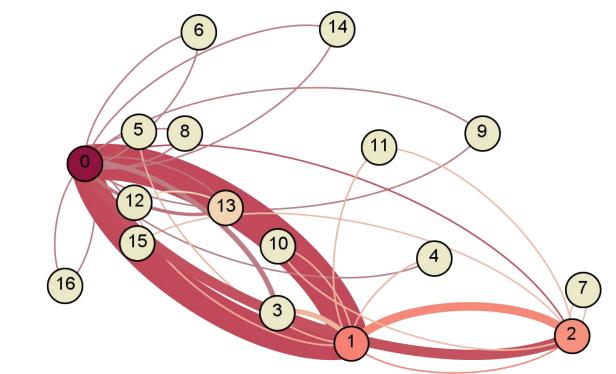
Each location is characterized by its frequency





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• Trips between points area aggregated as edges between nodes/locations

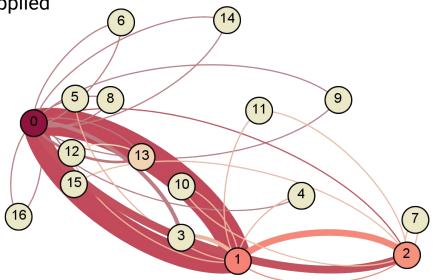


Inferring Home / Work locations

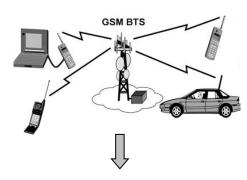
- Basic approach is based on frequency only
 - Most frequent location (L0) := Home
 - Second most frequent location (L1) := Work

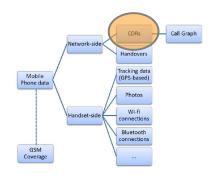
A minimum frequency threshold is applied

- Various alternatives & refinement are possible
 - Check time of stop & stay duration
 - Home: stop at 20-22, stay 8-10 hrs
 - Work: stop at 7-10, stay 6-9 hrs



Data gathered from mobile phone operator for billing purpose

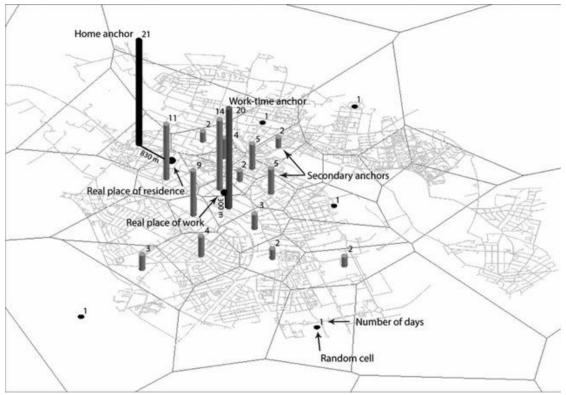




User id	Time start	Cell start	Cell end	Duration
10294595	"2014-02-20 14:24:58"	"PI010U2"	"PI010U1"	48
10294595	"2014-02-20 18:50:22"	"PI002G1"	"PI010U2"	78
10294595	"2014-02-21 09:19:51"	"PI080G1"	"PI016G1"	357

- "Personal Anchor Points": high-frequency visited places of a user
 - Select top 2 cells with max number of days with calls
 - Determine home and work through time constraints:
 - Based on average start time (AST) of calls and its deviation (std)
 - IF AST<17:00 & std<0.175 ⇒ WORK
 - ELSE HOME

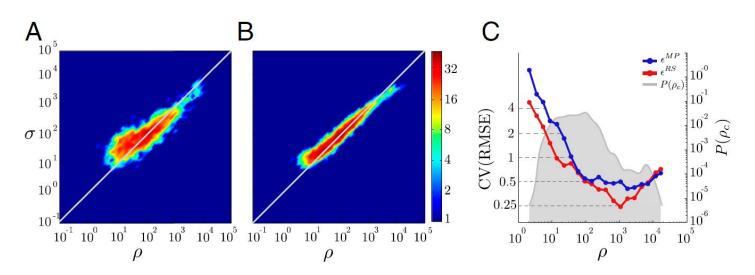
"Personal Anchor Points"



AHAS, R., SILM, S., JARV, O., SALUVEER, E., AND TIRU, M. 2010. Using mobile positioning data to model locations meaningful to users of mobile phones. Journal of Urban Technology 17, 1, 3–27.

- Estimating users' residence through night activity
 - Home = region with highest frequency of calls during nighttime
 - More suitable for larger scales
 - E.g. region = municipality

- Sample results on national level (France)
 - \circ estimate resident density (ρ) vs. real one (σ)



A = GSM data estimates

B = Environment/Infrastructures-based estimates

Objective: adding information to points / locations

Two main ways:

- Assign a single activity
- Assign a distribution of POIs / activity types

Given a dataset of GPS tracks of private vehicles, annotate trajectories with the most probable activities performed by the user.



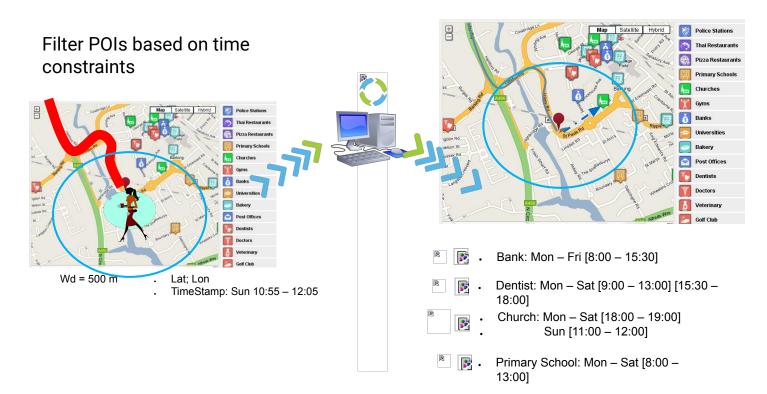
Associates the list of possible <u>POIs</u> (with corresponding probabilities) visited by a user moving by car when he stops.

A mapping between POIs categories and Transportation Engineering activities is necessary.

- POI collection from available sources, e.g. from Google Places.
- Association POI Activity: Each POI is associated to a ``activity". For example Restaurant → Eating/Food, Library → Education, etc.
- Basic elements/characteristics:
 - C(POI) = {category, opening hour, location}
 - C(Trajectory) = {stop duration, stop location, time of the day}
 - C(User) = {max walking distance}
- Computation of the probability to visit a POI/ to make an activity: For each POI, the probability of ``being visited" is a function of the POI, the trajectory and the user features.
- Annotated trajectory: The list of possible activities is then associated to a Stop based on the corresponding probability of visiting POIs



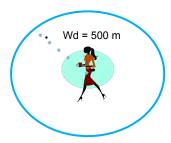
POI assignment

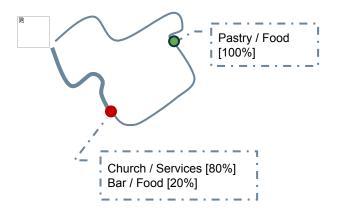


What if multiple POIs match

- Select closest one
- Assign a distribution of probability:



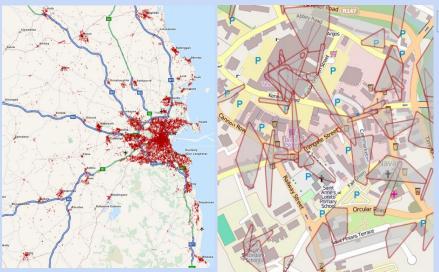




Reading social media to find POIs

An Irish experiment on Twitter

The points of each trajectory taken separately were grouped into spatial clusters of maximal radius 150m. For groups with at least 5 points, convex hulls have been built and spatial buffers of small width (5m) around them. 1,461,582 points belong to the clusters (89% of 1,637,346); 24,935 personal places have been extracted.



1150 1000 250 250

Statistical distribution of the number of places per person

Examples of extracted places

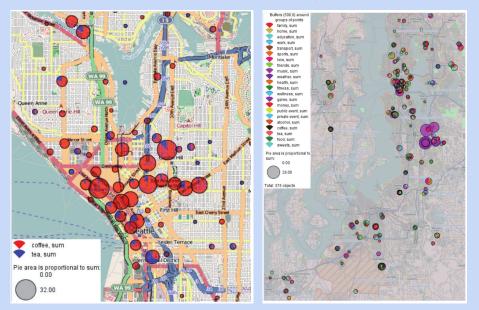
Reading social media to find POIs

Topics have been assigned to 208,391 messages (14.3% of the 1,461,582 points belonging to the personal places)

Message Fe	eatures	topic=family: Occurrences of topic		topic=education: Occurrences of topic	topic=work: Occurrences of topic	
@joe_lennon I usually ed	ducation	0	0	1	0	
@joe_lennon together ed	ducation	0	0	1	0	
@jas_103 deadly; don w	ork	0	0	0	1	
Just got home and see ho	ome	0	1	0	0	
So excited about my nest	veets	0	0	0	0	
@lamtcdizzy I haven't b sh	nopping	0	0	0	0	
Get in from my night ou fa	mily;home;work	1	1	0	1	
Home again at 6pm! Nhr	nme	n	1	0		
Bussing it home forti Get in from my night out, my dad gets home from work 1 0						
Ah shite. It's been a p	minutes later. Gre	at timing :)	0	0	Hacks	
@ronanhutchinson be ed	ducation	0	0	1		

- Some places did not get topic summaries (about 20% of the places)
- 2) In many places the topics are very much mixed
- The topics are not necessarily representative of the place type (e.g., topics near a supermarket: family, education, work, cafe, shopping, services, health care, friends, game, private event, food, sweets, coffee)

In the meanwhile, in Seattle...



G. Andrienko et al. Thematic Patterns in Georeferenced Tweets through Space-Time Visual Analytics. Computing in Science & Engineering, 2013.

Homeworks

How fast are users?

Choose one of the datasets seen at lesson (taxis, Geolife, etc.), select at least 10 users/vehicles and compute distributions of lengths. Remove 10% of points in each trajectory and repeat the distribution. Do the same for 20%, 30%, ... 90%. How does length distribution change?

Prepare a (well commented) python notebook

Implement a "speed-aware" trajectory compression method, that preserves speed, and test it on a dataset of your choice, e.g. a subset of taxis or Geolife users.

- Show the effects of simplification on some sample trajectories
- Study how the lengths of trajectories are affected
- Prepare a (well commented) python notebook

Inferring Home locations is often used to estimate the resident population of geographical areas. What are the existing approaches to face the problem?

- Make a research on Internet on the methods, including big data-based ones (GPS, GSM data, maybe satellite data or others) but also any other approach – e.g. coming from statistics/demography, sociology, etc.
- Prepare a blog (basically a survey) summarizing your discoveries.

Estimating GPS errors. Choose a bounding rectangle covering SF city. Download the road network/graph of that area. Select the GPS points of taxis in the same area. Assign each point P to its closest road segment R. Define pseudo-error(P) as the distance dist(P,R).

- Analyze the overall distribution of the pseudo-errors. Is it coherent with GPS.gov estimates of errors?
- Are pseudo-errors the same downtown vs. out of city?
- Prepare a (well commented) python notebook

The typical filtering algorithm relies on the individual. Define this new algorithm: label the points which have less than X neighbohrs in the set of points of all trips within an y radius as outliers, and discard them.

- Test it on SF taxi data
- Prepare a (well commented) python notebook