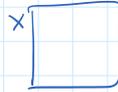


POLITICA QU-DEMAND

```

for ( ) {
    x[i] = f(y[i]);
}

```



start-process lc
 ↑
 @ for

cerchi la 1^a (prima) ^{linea}
 del codice della
 C1

LOAD Rbasey, R_i, R₋

caricamento di
 y[0] --- y[0-1]

CAU R_r, R_{ret}

cerca del codice P

STORE Rbasex, R_i, R₋

caricamento di

x[0] --- x[0-1]

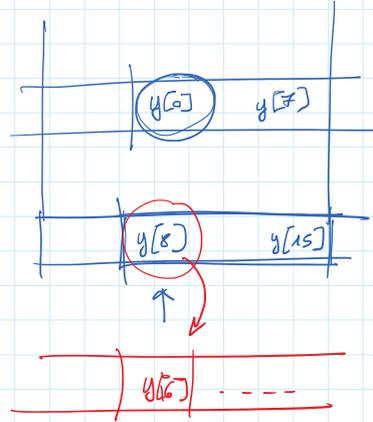
(non c'è un fatto vero
 e proprio)

PREFETCH

se accede ad un indirizzo $X \longrightarrow$ linea di cache
 \Rightarrow faccio in modo esincronamente di cercare
 in cache anche lo rigo successivo

prefetch(x)
(y)

LOAD



for (<N)

x[i] = f(y[i])

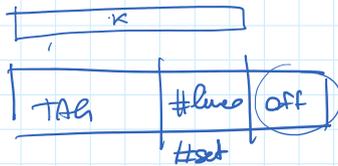
$$\#fault = fault(codice) + 2 \frac{N}{O}$$

con prefetch

$$= fault(codice) + 2$$

y[0] - y[7]
x[0] - x[7]

LOAD
STORE
CLO

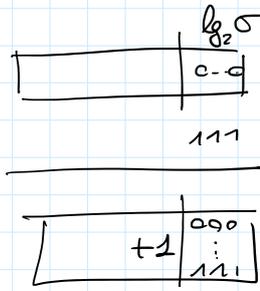


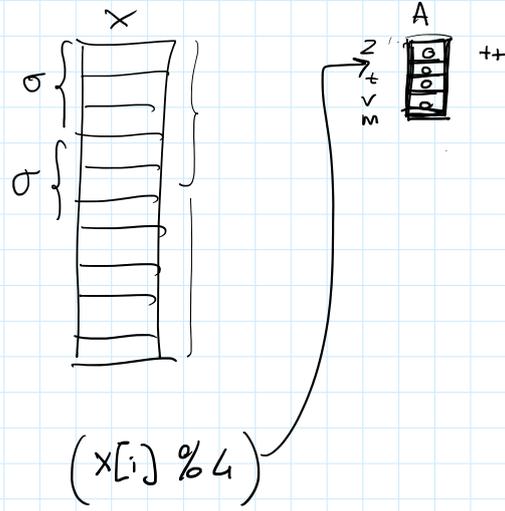
k.. 000
k.. (5-1)



-- prefetch(z[17]);

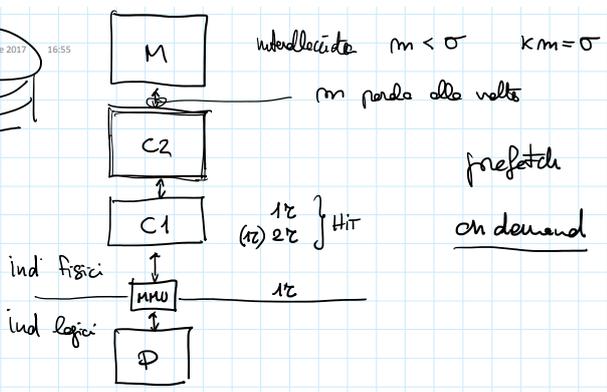
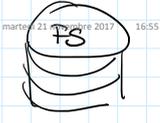
LOAD Rbase, Rindia, Rtarget, "prefetch"
↑
log





TR	BC	0
		cod
		$X[i] \dots$
		$(X[i] \% 4)$
		$X[i]$

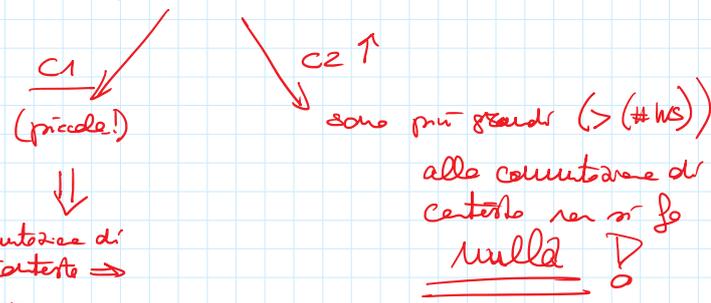
LOAD Rn, A, R12, #4, "non deallocate"
STORE



come si comportano

M-C_i

nei confronti dei
processori.



Primo dello cache è libera ← FIRST

(in caso di WRITE BACK => dopo aver scritto nei livelli superiori le linee modificate)

Esercizio

programma ①: scorre una lista $\boxed{\text{int}} \rightarrow$ e calcola la somma di tutti gli elementi N

vs.

programma ②: scorre un vettore e calcola la somma di tutti gli elementi N

Calcolare T_c
Saperne T_c

Cache 1° livello set assoc 4 vie $\sigma=8$ 1k linee
M interbloccata 4 moduli $T_M = 4\phi\tau$ $T_{tr} = 4\tau$

```
① [ s = 0;
    p = lista;
    while (p != null) {
        s = s + p.info;
        p = p.next;
    }
```

```
② [ s = 0;
    for (int i = 0; i < N; i++)
        s = s + x[i];
    END
```

WORKING SET
↓
#fault

WS = { codice: "while"
dati: s, p (og)
~ }

WS = { codice: "for"
dati: s, i, N
linea con x[i]
(PREFETCH) }

NON È WS!

	località	riuso
codice	✓	✓
x	✓	-
i, s	?	✓

↳ riuso x scalari ⇒ RTM!

```

S = ∅;
p = t.c.a.;
while (p != null) {
    S = S + p.info;
    p = p.next;
}
    
```

```

Hp: null = ∅
MOV R0, R3 ; ADD R0, R0, R3
MOV Rdest, Rp ; ADD Rdest, R0, Rp
while: .IF= Rp, R0, fine
        .LOAD Rp, R0, Rinfo
        .ADD Rinfo, R0, R3 ; s = s + p.info
        .LOAD Rp, #1, Rp ; p = p.next;
        .GOTO while
fine .END
    
```



manca nei conti un $\times 5$ (5 tempi di clock) x 5 (5 istruz. nel corpo del while!)

$$T_{cid} = 2(2\tau + t_a + \tau) + N(5(2\tau + t_a) + 2\tau + 2(2\tau + t_a) + (1\tau) + (1\tau)) + (2\tau)$$

do. ds exec Acc IF= toAD ADD GOTO IF= alto press

$$= \frac{6\tau + 2t_a}{init} + N(10\tau + 3t_a) + \frac{4\tau + t_a}{fine}$$

t_a : Hp: cache 1 set assoc "Hammerly Potters" $\Rightarrow 1\tau$
 $+ 1\tau$
nmw
 $t_a = 2\tau$

$$T_{cid} = (6\tau + 2t_a) + N(10\tau + 3t_a) + (4\tau + t_a)$$

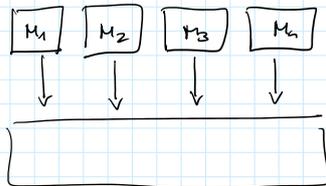
$$= (10\tau) + N(16\tau) + (6\tau) = 16\tau + N(16\tau)$$

$$= (N+1)(16\tau)$$

#fault

$\left. \begin{array}{l} 1 \text{ fault} \times \text{il codice} \\ \text{fault} \times \text{dati?} \\ N \text{ fault} \times \text{dati} \end{array} \right\} \# \text{fault} = (N+1)$

Cache 1° livello set assoc 4 vie $\sigma = 8$ 1k indirizzi
 M interdiscritti 4 moduli $\Sigma M = 4\phi\tau$ $T_{tr} = 4\tau$



$$2(4\tau + T_{tr}) + \frac{\sigma}{m} \Sigma M + \tau$$

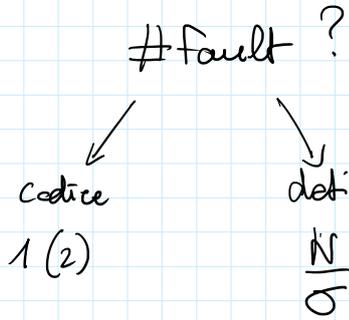
$$= 10\tau + 8\tau + \tau$$

$$T_{traf} = 91\tau$$

$$T_c = T_{cid} + \# \text{fault} \times T_{traf} = (N+1)16\tau + (N+1)91\tau$$

```
for( ) { s = s + x[i]; }
```

WS {
 codice: codice for
 dot: linee commentate di x (x[i])



```
LOAD Rbase, R1, Rtemp, "PREFETCH"  
ADD R1, Rtemp, R1
```

#Fault = 1
(dot)?

~~#Fault~~ = 2