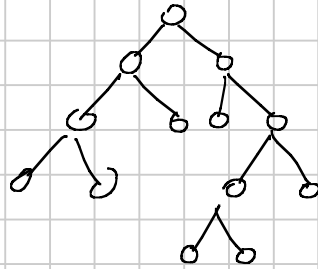
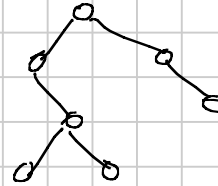


Es 1 : Verificare se un albero binario è completo



SI



NO

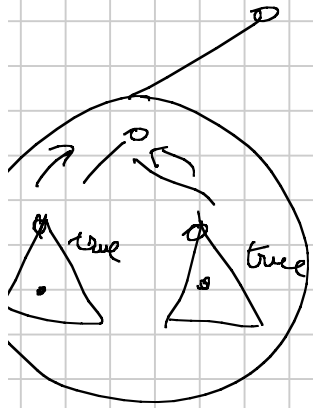
Ogni nodo di un albero completo ha sui figlio sinistro che
figlio destro $\neq \emptyset$

Completo (u):

if (u == NULL || (u.left == NULL && u.right == NULL))
 return true;

if (u.left == NULL || u.right == NULL) return false;

return completo(u.sx) && completo(u.dx);



Perf-Bil (u)

$\langle \underline{\text{bil}}, \text{alt} \rangle$

if ($u == \text{NULL}$) return $\langle \text{true}, -1 \rangle$

else {

$\langle \text{bil}_{sx}, \text{alt}_{sx} \rangle = \text{Perf-Bil}(u.sx);$

$\langle \text{bil}_{dx}, \text{alt}_{dx} \rangle = \text{Perf-Bil}(u.dx);$

bil = $(\text{bil}_{sx} \ \&\& \ \text{bil}_{dx} \ \&\& \ (\text{alt}_{sx} = \text{alt}_{dx}))$;

$\text{alt} = \max(\text{alt}_{sx}, \text{alt}_{dx}) + 1;$

return $(\text{bil}, \text{alt});$

}

$$T(n) = T(n_s) + T(n_d) + \Theta(1)$$

$$T(n) \approx 2T\left(\frac{n}{2}\right) + \Theta(1)$$

$$\Theta(n)$$

Contafoglie (u) :

if $u == \text{NULL}$ return ϕ ;

if ($u.sx == \text{NULL} \ \&\& \ u.dx == \text{NULL}$) return 1;

return Contafoglie(u.sx) + Contafoglie(u.dx);

nodou cardine $h(u) = p(u)$

altessa = profondità

cardine (u, p) 1° chiamata cardine $(u, 0)$

if $(u == NULL)$ return -1;

else {

• altessa_sx = cardine $(u.sx, p+1)$;

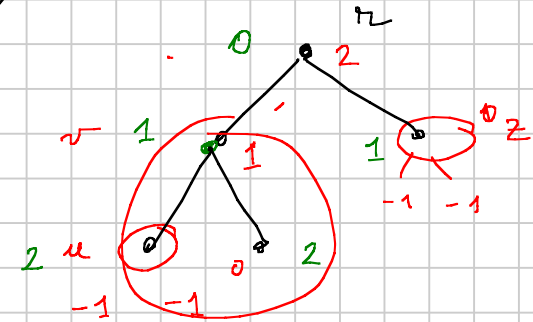
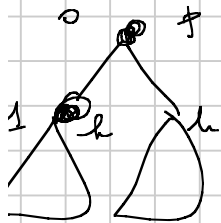
• altessa_dx = cardine $(u.dx, p+1)$;

altessa = $\max(\text{altessa}_{sx}, \text{altessa}_{dx}) + 1$

if $(p == \text{altessa})$ print $u.key$;

return altessa;

}



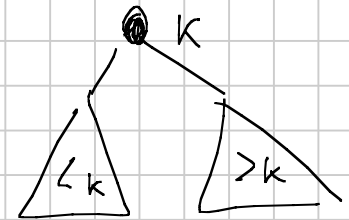
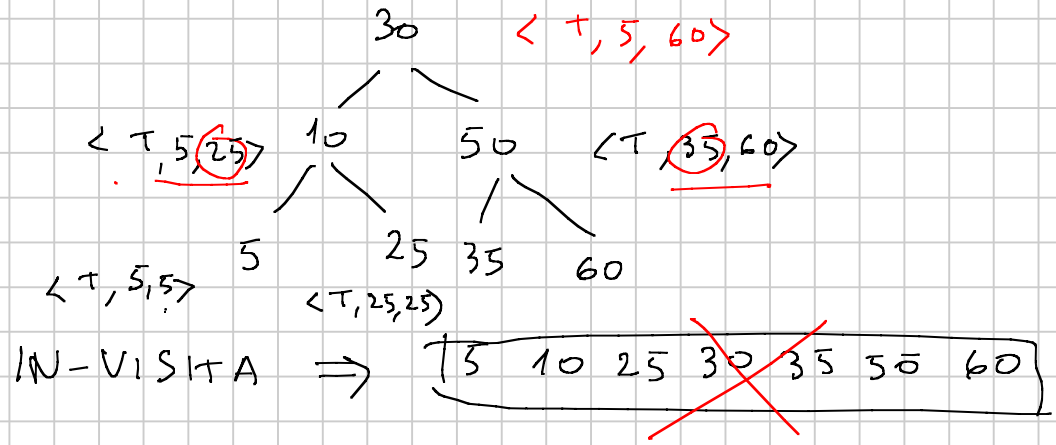
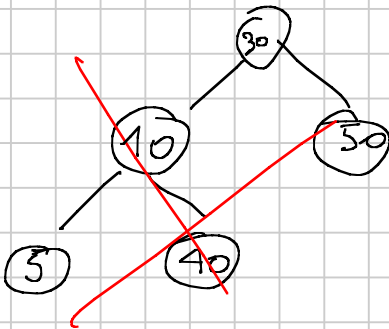
v è cardine

$\Theta(n)$

ABR

Verificare se un albero binario è un albero binario di ricerca

Definizione



Soluz. 1

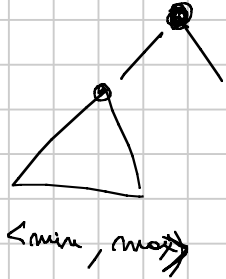
IN-VISITA: memorizzando in array A + controllo

tempo: $\Theta(n)$, $\Theta(n)$ spazio

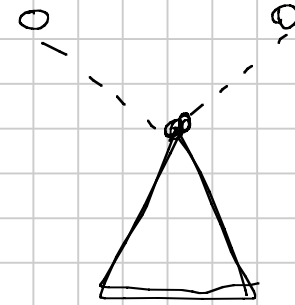
Sol. 2 : Invisita con controllo locale

ricordandoci l'ultimo visitato
 $\Theta(n)$ tempo

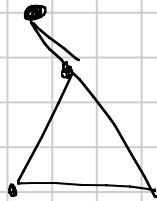
Sol. 3 :



$n. Key > max$



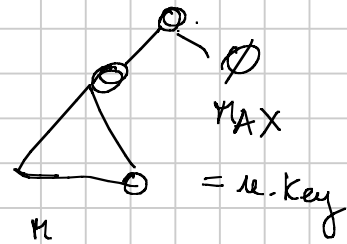
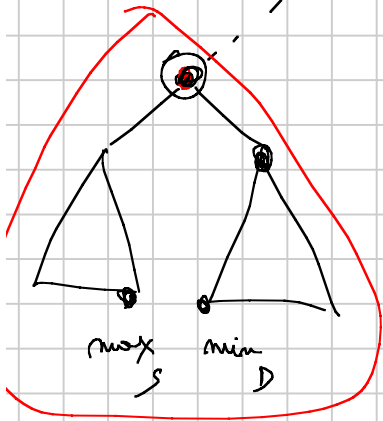
$\langle min, max \rangle$



$n. Key < min$

$\langle min, max \rangle$

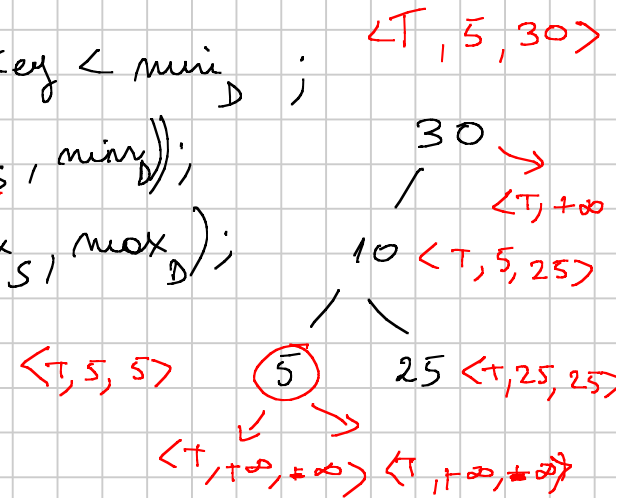
usando puntatore
al padre



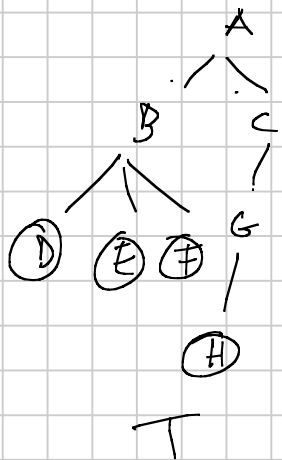
ABR (u) : $\langle isABR, min, max \rangle$
 if (u == NULL) return $\langle true, +\infty, -\infty \rangle$;

$\langle isABR_s, min_s, max_s \rangle = ABR(u.sx)$;
 $\langle isABR_D, min_D, max_D \rangle = ABR(u.dx)$;

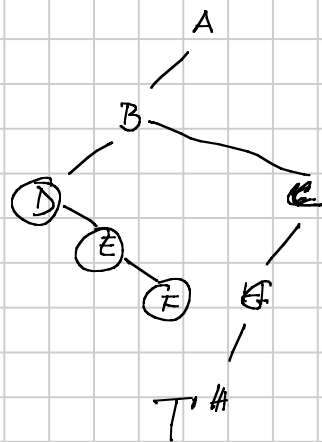
$isABR = isABR_s \& \& isABR_D \& \& max_s < u.Key < min_D$;
 $min = Mathmin(u.Key, Mathmin(min_s, min_D))$;
 $max = Mathmax(u.Key, Mathmax(max_s, max_D))$;
 return $\langle isABR, min, max \rangle$;



Contare le foglie di un albero dalla sua immagine binaria



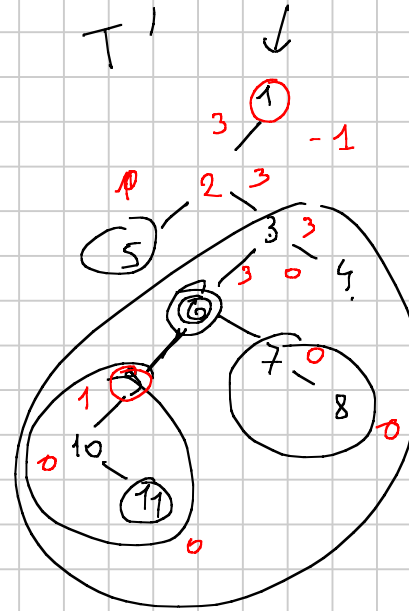
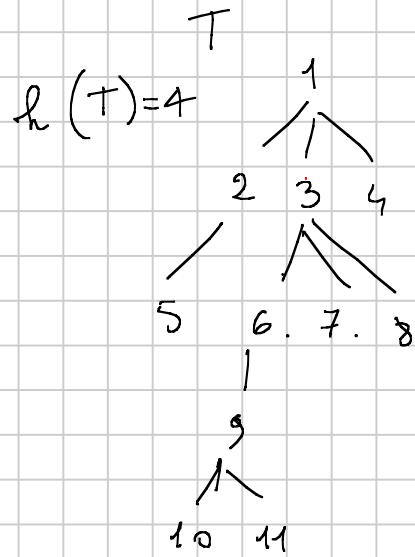
⇒



$h(T') = 4$

foglia di T
se figli suoi = \emptyset

calcolare l'altezza di T da T'



$h(T') = 6$

$= 4$

0

$\max(h_{SIN} + 1, h_{DES})$

Calcolo l'attesa delle radici di T.

$\max(-1 + 1, 0)$

```
Alt(u):  
if (u == NULL) return -1;  
altSIN = Alt(u.sx);  
altDES = Alt(u.dx);  
alt = max(altSIN + 1, altDES);  
return alt;
```