

EDIT DISTANCE

X e Y $X = x_1 \dots x_m$ $Y = y_1 \dots y_m$

errori

mismatch

blank in X che corrisponde a un carattere in Y \equiv inserzione

blank in Y " " " " " X \equiv cancellazione

inversione

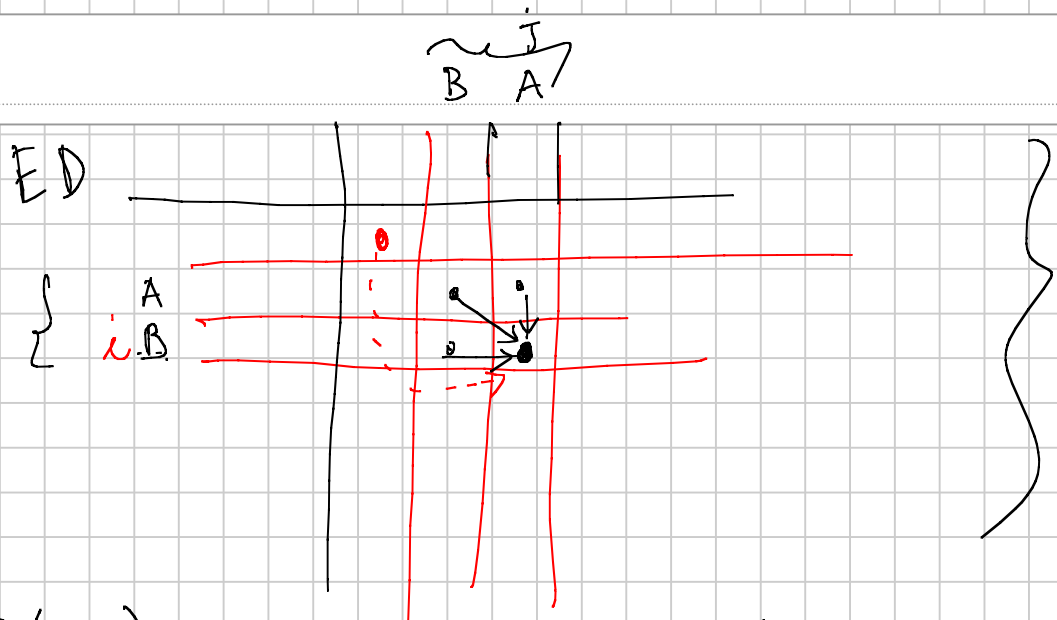
$X = \text{ATLETA}$

$Y = \text{ALTERA}$

$X = \dots ab \dots$

$Y = \dots \underbrace{ba} \dots$

1 errore



$$ED(i, j) = \min (ED(i-1, j-1) + p(i, j), ED(i, j-1) + 1, ED(i-1, j) + 1, ED(i-2, j-2) + 1 \text{ if } (x_i == y_{j-1} \& \& x_{i-1} == y_j))$$

ED

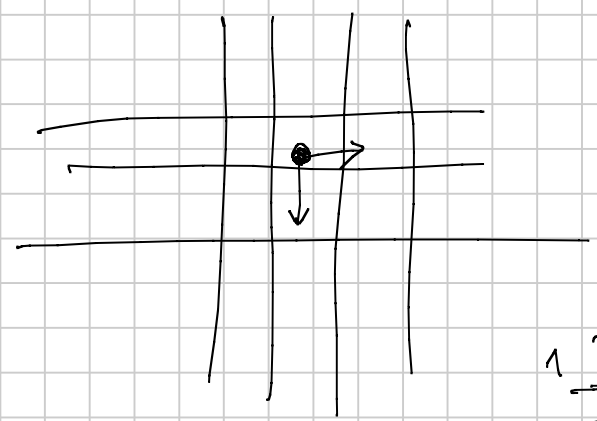
X \ Y	Ø	A	L	I	E	R	A
Ø	Ø	1	2	3	4	5	6
A	1	0	1	2	3	4	5
T	2	1	1	2	3	4	5
L	3	2	1	2	3	4	5
E	4	3	2	2	3	4	5
T	5	4	3	2	2	3	4
A	6	5	4	3	3	3	2

$$\begin{array}{r} \text{ATLETA} \\ \text{ALTERA} \\ \hline = 1 = 1 = \end{array}$$
 2 errori

$$\begin{array}{r} \text{ALTEERA} \\ \text{ATLETA} \\ \hline = 1 = 1 = \end{array}$$
 3 errori

1 unica soluzione ottima con 2 errori

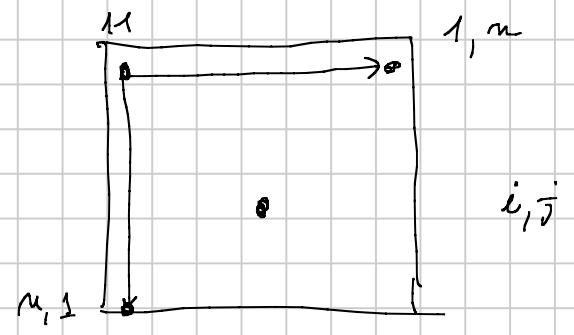
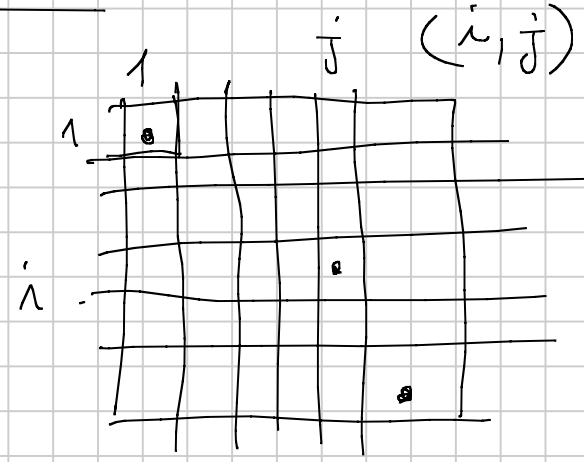
Scegliere di $n \times n$ caselle

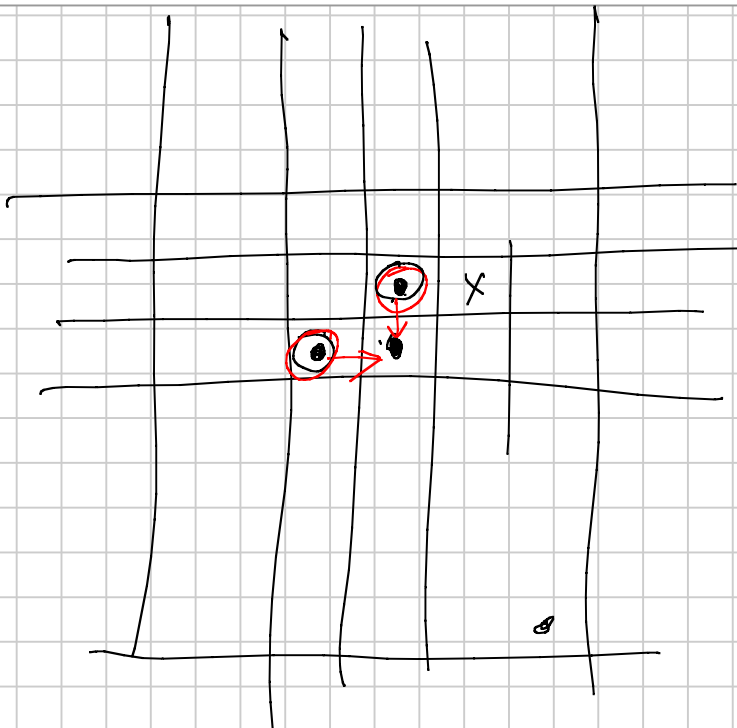


P.D.

$$(i, j) \rightarrow (i+1, j), i < n$$

$$(i, j) \rightarrow (i, j+1), j < n$$

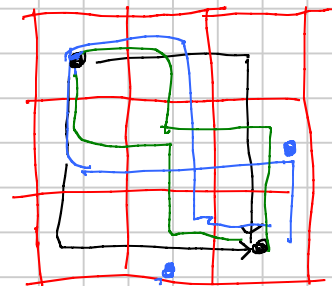




$$\underline{P(i, 0) = P(0, j) = 1}$$

$$P(m, m)$$

$$P(i, j) = P(i-1, j) + P(i, j-1)$$



$$P(3, 3) = 6$$

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Contapercorsi (n);
P = nuova matrice (n x n)
for i = 1 to n      P[i, 1] = 1;
for j = 1 to n      P[1, j] = 1;
for i = 2 to n
  for j = 2 to n    P[i, j] = P[i-1, j] + P[i, j-1];
return P(n, n);

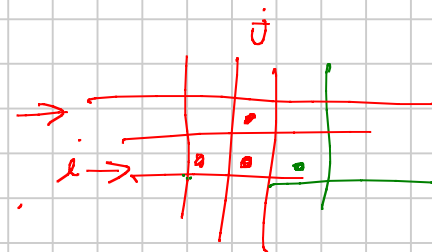
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$\Theta(n^2)$ tempo

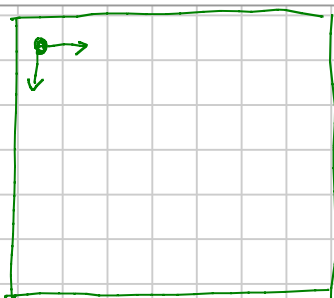
$\Theta(n^2)$ spazio

\Downarrow
 $\Theta(n)$ memorizzando
 le ultime 2 righe

Memorizzare tutta la matrice non serve.



Spazio costante?



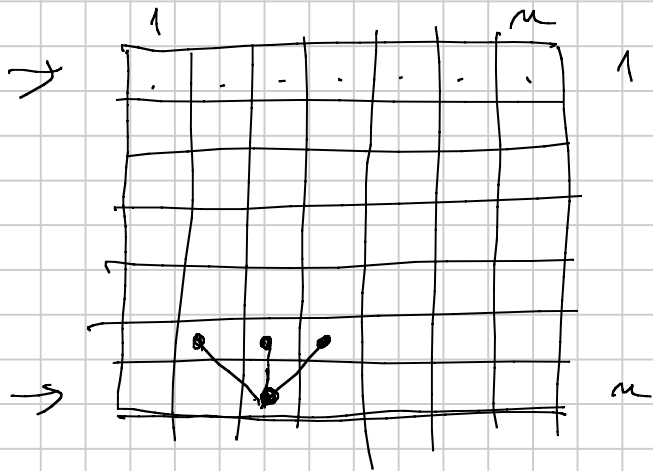
$(i, j) \rightarrow (i, j+1)$ a destra
 $(i, j) \rightarrow (i+1, j)$ in basso

c_{ij} si acquista un profitto c_{ij}

Alg. di PD che calcoli il percorso da $(1, 1)$ a (n, n) col massimo profitto.



$$\max (S[i-1, j], S[i, j-1]) + c[i, j]$$



quando si visita x, y si prende
un profitto $P(x, y) \geq 0$;

Calcolare il percorso da una qualunque cella dell'ultima
riga a una qualunque cella della riga superiore che
produce il profitto max, con un algoritmo di programmazione
dinamica.

$S(x, y)$ il massimo profitto per arrivare a x, y
 $S(1, y)$ " " per arrivare alla riga 1

Regole ricorsive

$$i < m \quad \left\{ \begin{array}{l} S(m, j) = P(m, j) \\ S(i, j) = \max \left(S(i+1, j), S(i+1, j-1), S(i+1, j+1) + P(i, j) \right); \end{array} \right.$$

Ricerca cammino (P, m):

S = matrix di dimensione $n \times n$;

$M =$ $\text{for } j=1 \text{ to } n \quad S[n, j] = P[n, j];$

$\text{for } i=n-1 \text{ downto } 1$

$\text{for } j=1 \text{ to } n \{$

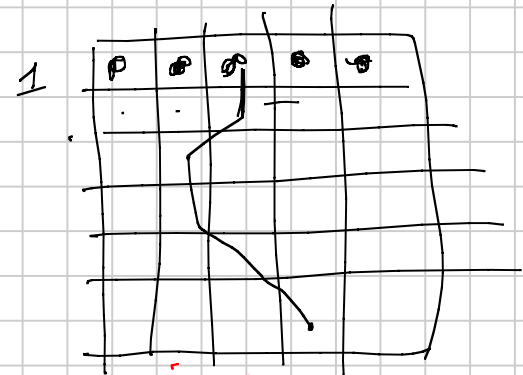
$S[i, j] = S[i+1, j]; \quad \uparrow \quad M[i, j] = 0$

$\text{if } (j > 1) \ \&\& \ (S[i+1, j-1] > S[i, j]) \quad S[i, j] = S[i+1, j-1]; \quad M[i, j] =$

$\text{if } (j < n) \ \&\& \ (S[i+1, j+1] > S[i, j]) \quad S[i, j] = S[i+1, j+1]; \quad M[i, j] =$

$S[i, j] + \phi[i, j]$

$\}$



aggiunte per ricostruire il cammino

Percece casello (n):

$max = S[1,1]; j_{max} = 1;$

for $j = 2$ to n {

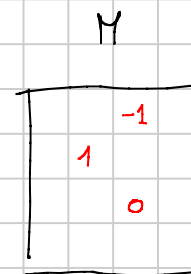
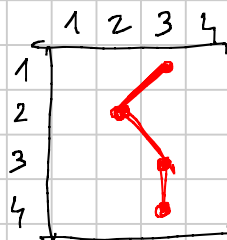
if $S[1,j] > max$ {

$max = S[1,j]$

$j_{max} = j$

}

} return $j_{max};$



max in $S(1,3)$