

TEOREMA DELL' ESPERTO

$$\begin{cases} T(n) = a T\left(\frac{n}{b}\right) + f(n) & a \geq 1 \quad b > 1 \\ T(n) = \Theta(1) & \text{per } n=1 \end{cases}$$

$n = b^i$

metodo iterativo

$$T(n) = a T\left(\frac{n}{b}\right) + f(n) \qquad T\left(\frac{n}{b}\right) = a T\left(\frac{n}{b^2}\right) + f\left(\frac{n}{b}\right)$$

$$T(n) = a \left(a T\left(\frac{n}{b^2}\right) + f\left(\frac{n}{b}\right) \right) + f(n) \qquad T\left(\frac{n}{b^2}\right) = a T\left(\frac{n}{b^3}\right) + f\left(\frac{n}{b^2}\right)$$

$$T(n) = a \left(a \left(a T\left(\frac{n}{b^3}\right) + f\left(\frac{n}{b^2}\right) \right) + f\left(\frac{n}{b}\right) + f(n) \right)$$

$$T(n) = a \left(a \left(a \dots \left(a T\left(\frac{n}{b^i}\right) + f\left(\frac{n}{b^{i-1}}\right) \right) + f\left(\frac{n}{b^{i-2}}\right) \right) \dots + f\left(\frac{n}{b}\right) \right) + f(n)$$

$$T\left(\frac{n}{b^i}\right) = \Theta(1)$$

$$T(n) = a^i T\left(\frac{n}{b^i}\right) \quad \Sigma$$

$$\frac{n}{b^i} = 1$$

$$n = b^i$$

$$\log_b n = \log_b b^i$$

$$i = \log_b n$$

$$= a^i \Theta(1) + a^{i-1} f\left(\frac{n}{b^{i-1}}\right) + a^{i-2} f\left(\frac{n}{b^{i-2}}\right) + \dots + a f\left(\frac{n}{b}\right) + a^0 f\left(\frac{n}{b^0}\right)$$

$$\frac{n \log_b a}{\log_b a} : \frac{f(n)}{1}$$

$$= a^{\log_b n} \cdot \Theta(1) + \sum_{i=0}^{\log_b n - 1} a^i f\left(\frac{n}{b^i}\right)$$

1) peso di più
 $\Theta(n \log_b a)$

$$= \underbrace{\Theta(n^{\log_b a})}_{\text{Risolti}} + \underbrace{\sum_{i=0}^{\log_b a - 1} a^i f\left(\frac{n}{b^i}\right)}_{\text{Combinare e Dividere}}$$

2) se
 $\Theta(n \log_b a) : f(n)$

3) $f(n)$, peso di più

MOLTIPLICAZIONE

Carta e penna $\Theta(n^2)$

MOLTIGIZIA $\Theta(n^2)$

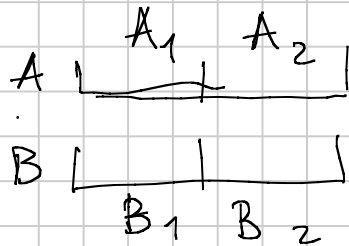
MOL-DI ?

A e B
di n cifre

bit model
(RAM)
 A^n

$$\begin{cases} T(n) = \Theta(1) & \text{se } n = 1 \\ T(n) = 4T\left(\frac{n}{2}\right) + \Theta(n) \end{cases}$$

↓
somme
e traslazioni



Dividi et Impero

- Dividi A e B in n cifre in A_1, B_1, A_2, B_2 di n/2 cifre
- Risoli ricorsivamente
- Combine

$$A = 1587 = 15 \cdot 10^2 + 87$$

$$1500 + 87$$

$$A * B = (A_1 \cdot 10^{n/2} + A_2) * (B_1 \cdot 10^{n/2} + B_2)$$

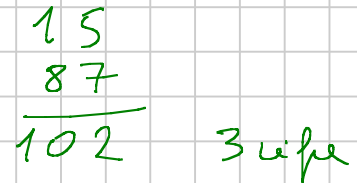
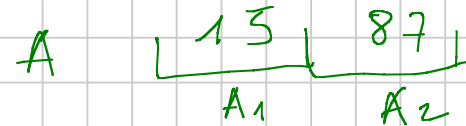
$$\Rightarrow \underbrace{A_1 B_1 \cdot 10^n + A_1 B_2 \cdot 10^{n/2} + A_2 B_1 \cdot 10^{n/2} + A_2 B_2}_{\text{Combine}}$$

$$T(n) = 4T\left(\frac{n}{2}\right) + \Theta(n)$$

$$a=4 \quad b=2 \quad n^{\log_2 4} = n^2$$

$$n : n^2 \quad f(n) = O(n^{2-\epsilon}) \quad \forall \epsilon \leq 1$$

caso 1) $T(n) = \Theta(n^2)$



$$(A_1 + A_2)(B_1 + B_2) = (A_1 B_1 + A_2 B_2 + A_2 B_1 + A_1 B_2)$$

$$A_2 B_1 + A_1 B_2 = (A_1 + A_2)(B_1 + B_2) - A_1 B_1 - A_2 B_2$$

$$A * B = \underbrace{A_1 B_1}_{\text{red}} 10^n + \left(\underbrace{(A_1 + A_2)(B_1 + B_2)}_{\substack{\text{green} \\ \uparrow \\ (n/2+1) \text{ cifre}}} - \underbrace{A_1 B_1}_{\text{red}} - \underbrace{A_2 B_2}_{\text{green}} \right) 10^{n/2} + \underbrace{A_2 B_2}_{\text{green}}$$

MOLT RAPIDA (X, Y, n)

risultati
dir. $\Theta(1)$

if $(n == 1)$ return $X * Y$
else {

$$T(n) = \Theta(1) \quad \text{per } n=1$$

$$T(n) = 3T\left(\frac{n}{2}\right) + \Theta(n)$$

dividi $\Theta(1)$

"dividi X in X_1 e X_2 di $n/2$ cifre"
"dividi Y in Y_1 e Y_2 di $n/2$ cifre"

$$a = 3 \quad b = 2$$

$$n^{\log_2 3} = n^{1.58}$$

$$n = O(n^{1.58 - \epsilon})$$

$$\text{per } \epsilon \leq 0.58$$

caso 1

risultati
ricorsivamente

$A = \text{MOLT RAPIDA}(X_1, Y_1, \underline{n/2});$

$B = \text{MOLT RAPIDA}(X_2, Y_2, \underline{n/2});$

$C = \text{MOLT RAPIDA}(X_1 + X_2, Y_1 + Y_2, \underline{n/2});$

combine

return $(A 10^n + (C - A - B) 10^{n/2} + B);$

}

$$T(n) = \Theta(n^{1.58})$$

MOLTIPLICAZIONE MATRICI

A e B $n \times n$

input size : n

$$\begin{bmatrix} a_{11} & \dots & a_{1n} \end{bmatrix} + \begin{bmatrix} b_{11} \end{bmatrix} =$$

$$S = A + B \quad \Theta(n^2)$$

$C = A \times B$ = prodotto righe \times colonne

$$c_{ij} = \sum_{k=1}^n a_{ik} \cdot b_{kj}$$

$$A = \begin{bmatrix} \text{---} \\ \bullet & \bullet & \bullet \\ \text{---} \end{bmatrix} B = \begin{bmatrix} \bullet \\ \bullet \\ \bullet \end{bmatrix} = \begin{bmatrix} c_{21} \\ \bullet \end{bmatrix} C$$

calcolare C richiede

$$\Theta(n^2) \text{ same}$$

$$\Theta(n^3) \text{ mol}$$

Divide et Impera

NO

$$T(n) = 8 T\left(\frac{n}{2}\right) + \Theta(n^2)$$

$$T(1) = \Theta(1)$$

$$A \begin{array}{|c|c|} \hline A_{11} & A_{12} \\ \hline A_{21} & A_{22} \\ \hline \end{array} \times B \begin{array}{|c|c|} \hline B_{11} & B_{12} \\ \hline B_{21} & B_{22} \\ \hline \end{array} = C \begin{array}{|c|c|} \hline C_{11} & C_{12} \\ \hline C_{21} & C_{22} \\ \hline \end{array}$$

$$C_{11} = A_{11} \cdot B_{11} \oplus A_{12} \cdot B_{21}$$

$$C_{21} = A_{21} \cdot B_{11} \oplus A_{22} \cdot B_{21}$$

$$C_{12} = A_{11} \cdot B_{12} \oplus A_{12} \cdot B_{22}$$

$$C_{22} = A_{21} \cdot B_{12} \oplus A_{22} \cdot B_{22}$$

$$a = 8 \quad b = 2$$

$$n^{\log_2 8} = n^3$$

$$f(n) = O(n^{3-\epsilon}) \quad \epsilon \leq 1$$

$$T(n) = \Theta(n^3)$$

$$X_1 = (A_{11} + A_{22}) \cdot (B_{11} + B_{22})$$

$$X_2 = (A_{21} + A_{22}) \cdot B_{11}$$

$$X_3 = A_{11} \cdot (B_{12} - B_{22})$$

$$X_4 = A_{22} \cdot (B_{21} - B_{11})$$

$$X_5 = (A_{11} + A_{12}) \cdot B_{22}$$

$$X_6 = (A_{11} - A_{21}) \cdot (B_{11} + B_{12})$$

$$X_7 = (A_{12} - A_{22}) \cdot (B_{21} + B_{22})$$

Prodotti di sottomatrici
Size $n/2$

$$C_{11} = X_1 + X_2 - X_5 - X_7$$

$$C_{12} = X_3 + X_5$$

$$C_{21} = X_2 + X_4$$

$$C_{22} = X_1 + X_3 - X_2 - X_6$$

$$\Omega(n^2)$$

$$T(n) = 7 T\left(\frac{n}{2}\right) + \underbrace{18 n^2}_{\Theta(n^2)}$$

$$n^{\log_7 7} = n^{\log_2 7} = n^{2.81}$$

$$n^2 \ll n^{2.81}$$

STRASSEN

$$T(n) = \Theta(n^{2.81})$$

$$f(n) = O\left(n^{2.81 - \epsilon}\right)$$

$\epsilon \leq 0.08$

VINOGRAD

$$\Theta(n^{2.376})$$

limite inferiore $\Omega(n^2)$ Bini, Lotti, Romani, Caprioni $\Theta(n^{2.75})$



