The rsync algorithm

https://rsync.samba.org/tech_report/tech_report.html

An easy problem

 I have two files A and B. I want to make B equals to A

- What is the cost?
 - CPU
 - Data moved (reads, writes)

The problem of rsync

- A is stored in computer alpha and B in computer beta
- The network link can be slow (at least it is much slower than CPU)

• How can I save bandwidth?

A naïve approach

- Beta compute a hash of the file B and send it to alpha
- Alpha compute the hash of A and send back to beta either the hash (if the two hash are the same) or the content of A if they differ
- Beta check if the message is the hash or has to update B
- What is the cost?
- What is the hash function?

Cryptographic hash

- 1. Deterministic
- 2. Quick to compute
- 3. Infeasible to generate a message from the hash
- 4. A small change in the message should drastically change the hash
- 5. It is infeasible to find collisions

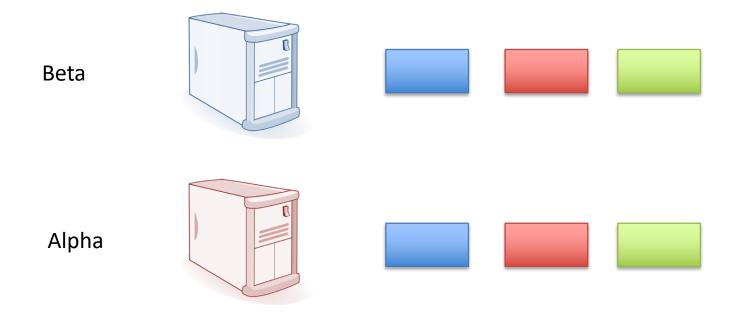
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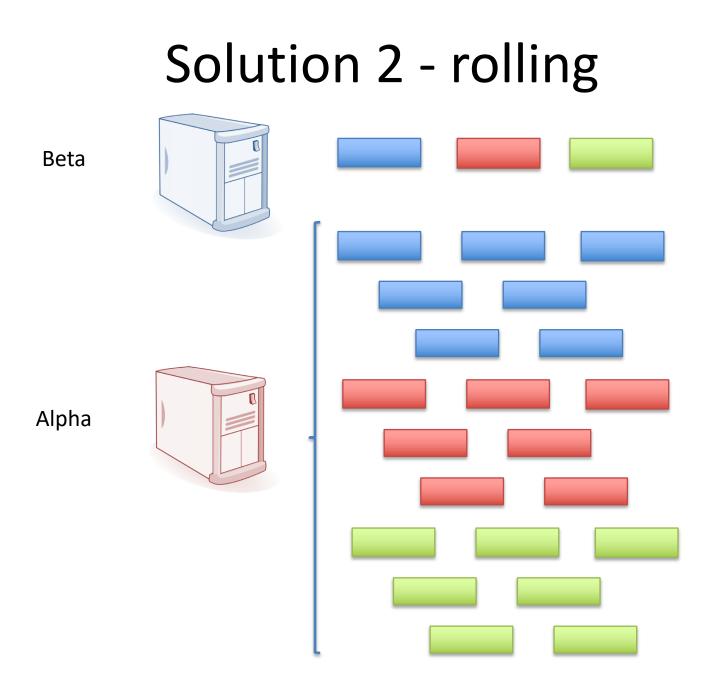
Can I do better?

 Can I save bandwidth when A and B are similar?

Solution 1 - bucketing



- Weakness?
- Can I do better?



Can I do better?

• Intense use of cpu in **alpha**

Solution 3 – rolling hashing

A two hashing strategy

$$Document = X_{1,}X_{2} \dots X_{n}$$

$$a(k,l) = \left(\sum_{i=k}^{l} X_i\right) \mod M$$
$$b(k,l) = \left(\sum_{i=k}^{l} (l-i+1)X_i\right) \mod M$$
$$s(k,l) = a(k,l) + 2^{16} b(k,l)$$

Solution 3 – rolling hashing

A convenient way to derive next hash

$$a(k+1, l+1) = (a(k, l) + X_{l+1} - X_k) \mod M$$

$$b(k + 1, l + 1) = (b(k, l) - (l - k + 1)X_k + a(k + 1, l + 1)) \mod M$$

- Is it M=2¹⁶ a good idea?
- Collisions?

Update an example (1)

- Sequence: ABCDE
- Window size: 4
- Get rid of the modulo for simplicity
- a(1,4) = A + B + C + D
- a(2, 5) = a(1,4) A + E == A + B + C + D - A + E == B + C + D + E

Update an example (2)

- Sequence: ABCDE, window size = 4
- b(1,4) = 4A + 3B + 2C + 1D

=

- b(2,5) = b(1,4) 4A + a(2,5) =
 - = 4A + 3B + 2C + 1D 4A + a(2,5) =
 - = 3B + 2C + 1D + a(2,5) =
 - = 3B + 2C + 1D + B + C + D + E =
 - 4B + 3C + 2D + E

Can I do better?

- Collision probability high enough to ensure equality of blocks
- One scan of the file A in alpha for each block of B in beta

Solution 4 - rsync

- Use two hash functions
- The rolling hashing for each possible offset
- A stronger 128bit hash in case a collision is detected
 - Rsync uses MD4

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- Use two hash functions
- The rolling hashing for each possible offset
- A stronger 128bit hash in case a collision is detected
 - Rsync uses MD4
- How to generate collisions in MD4

 https://eprint.iacr.org/2005/151.pdf

Checksum searching

- Beta send several checksums
- For each test alpha performs a search on these checksums

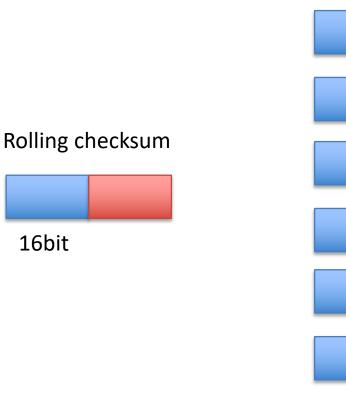
• Is linear scanning an option?

Checksum searching: possible solutions

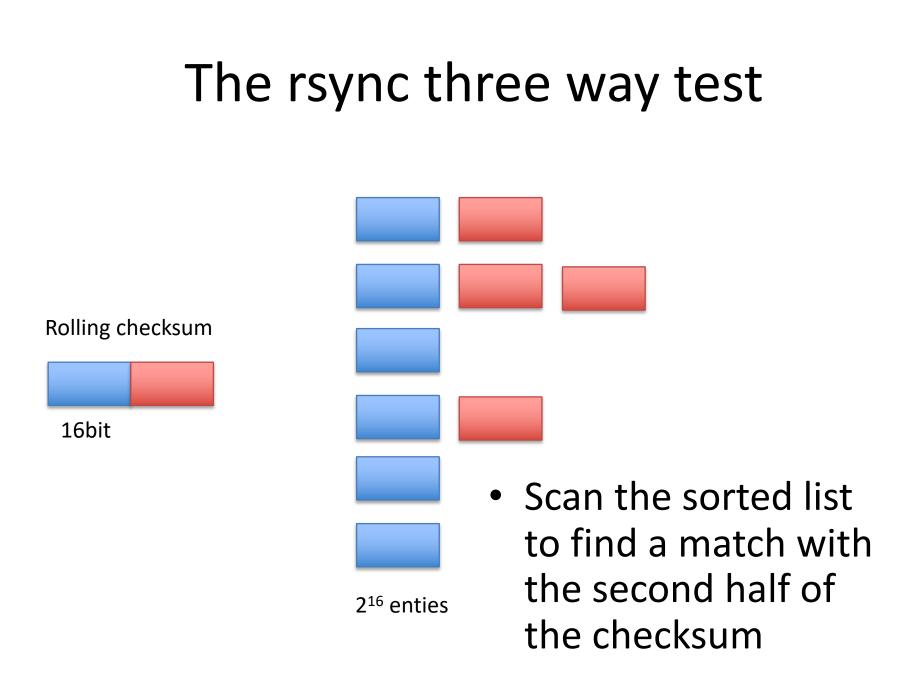
- Binary search
 - Preprocessing requires sorting O(n lg n)
 - Searching requires O (lg n0
- Bloom filters
 - Constant time insert and query, but can have false positives
- Perfect hashing
 - Preprocessing space/time tradeoff
 - Constant time searching

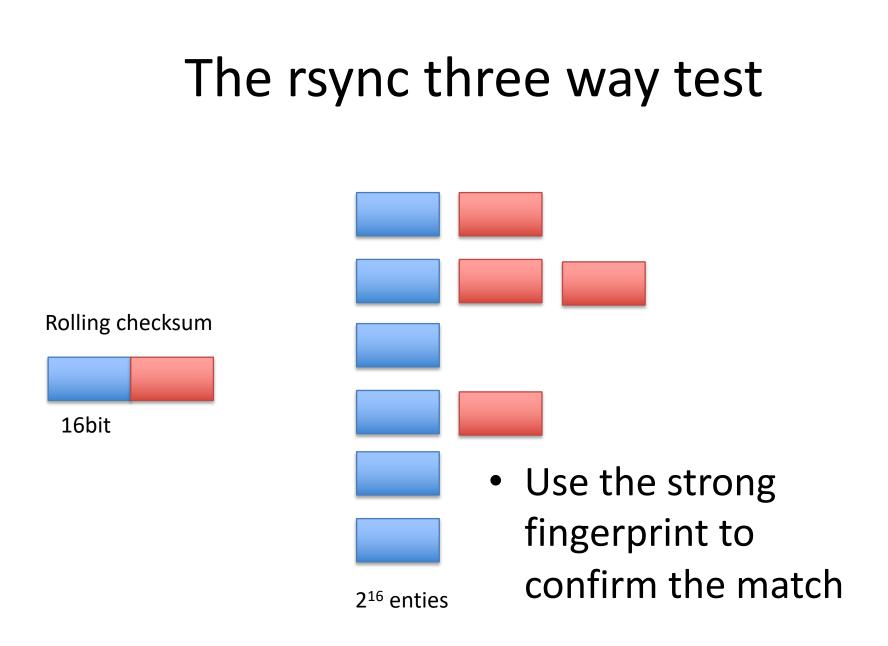
The rsync three way test

2¹⁶ enties



- Search for a match in the table
 - If nul the block is not found





The rsync three way test

- What happens if two blocks in B have the same fingerprint?
- How the list of blocks can be organized?
- Is it possible to copy a corrupted file?

Things you may want to try and discuss next week

- Test binary search or perfect hashing
- Test the impact of the length of the block
- Small vs huge files