

$$G = (V, E)$$

$k$ -Coloring:  $\chi: V \rightarrow \{1, 2, \dots, k\}$

s.t.  $\forall (u, v) \in E : \chi(u) \neq \chi(v)$

# Decision

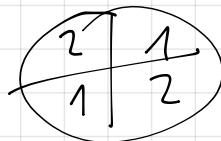
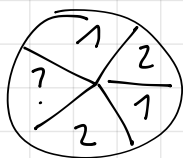
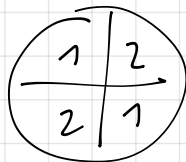
Graph  $G = (V, E)$ , an integer  $k$

Q: Has  $G$  a  $k$ -coloring?

## PLANAR GRAPHS

$k \geq 4$  : answer = YES (because of 4-color theorem)

$k = 2$  : Y . N



$COLOR(u, c):$        $c \in \{1, 2\}$

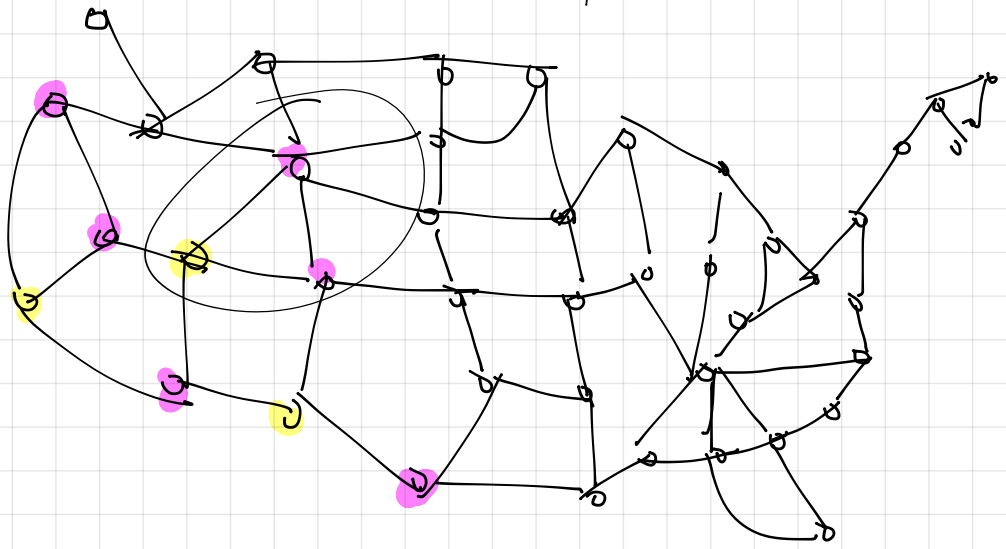
$X(u) := c;$        $\bar{c} = \{1, 2\} - c$

for  $v \in adj(u)$  do

if  $X(v)$  is undefined then  
 $COLOR(v, \bar{c})$

elif if  $X(v) = c$  then it's not 2-col

NO



$K=3$

$\in$  NPC

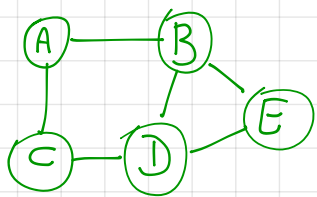
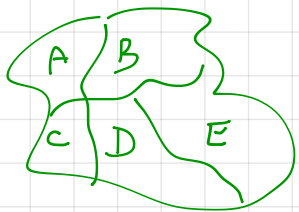
3-COL

proof

3-SAT  $\leq$  3-COL (too long here)

What we see here is 3-COL  $\leq$  3-SAT

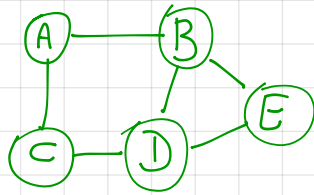
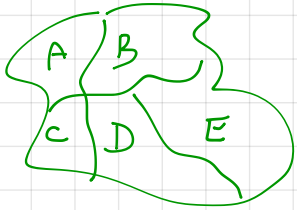
Reduction: show how expressive are Boolean formulas



pleneu jept

3-COL  $\leq$  SAT  $\leq$  3-SAT

└──────────┘ we know this



Boolean variables:

$x_A^1$   $x_A^2$   $x_A^3$  : if country A has color  $\{1, 2, 3\}$   
 $\vdots$   
 $x_E^1$   $x_E^2$   $x_E^3$

Formulas

$$\varphi(x, y, z) = (x \wedge \bar{y} \wedge \bar{z}) \vee (\bar{x} \wedge y \wedge \bar{z}) \vee (\bar{x} \wedge \bar{y} \wedge z)$$

$\varphi(x, y, z)$  is true iff exactly one of  $x, y, z$  is true  
 We need to transform  $\varphi$  into CNF:  $\uparrow$

$$F \equiv \hat{\varphi}(x_A^1, x_A^2, x_A^3) \wedge \dots \wedge \hat{\varphi}(x_E^1, x_E^2, x_E^3) \wedge \boxed{\dots}$$

$$\gamma(x_0, y_0, z_0, x_1, y_1, z_1) \equiv (x_0 \wedge \bar{x}_1) \vee (y_0 \wedge \bar{y}_1) \vee (z_0 \wedge \bar{z}_1)$$

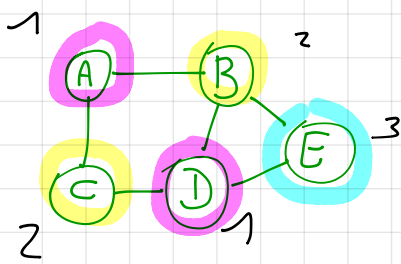
$$\gamma(x_A^1, x_A^2, x_A^3, x_B^1, x_B^2, x_B^3) \rightarrow \neg \text{edge}(A, B)$$

is true iff

A and B have different colors

$\hookrightarrow \neg$





$$\begin{array}{l}
 X_A^1 = T \quad X_A^2 = X_A^3 = F \\
 X_B^2 = T \quad \dots \dots \dots \\
 X_C^2 = T \quad \dots \dots \dots
 \end{array}$$

$$F \equiv \underbrace{\hat{\varphi}(x_A^1, x_A^2, x_A^3)} \wedge \dots \wedge \underbrace{\hat{\varphi}(x_E^1, x_E^2, x_E^3)} \wedge \boxed{\dots}$$

$$\begin{aligned}
 & \wedge \hat{\tau}(A, B) \wedge \hat{\tau}(A, C) \wedge \hat{\tau}(C, D) \\
 & \wedge \hat{\tau}(B, D) \wedge \hat{\tau}(B, E) \wedge \hat{\tau}(D, E)
 \end{aligned}$$

where  $\hat{\tau}(A, B)$  is the CNF of  $\tau(A, B)$