Models of computation (MOD) 2014/15

Exam – July 1, 2015

[Ex. 1] Add to IMP the command

fix a of c

that keeps executing c until the value of the expression a is not changed by the last execution of c (for example fix x of x := x + 1 diverges for any σ , while fix x of (x := 1; y := y + 1) terminates after having incremented y by 2 if initially $\sigma(x) \neq 1$ or by 1 otherwise).

- 1. Define the operational semantics of the new command.
- 2. Define the denotational semantics of the new command, paying attention to the introduction of lifting whenever necessary. *Hint:* defines the semantics as the fixpoint of a suitable function $\Gamma_{a.c.}$
- 3. Extend the proof of completeness of the denotational semantics to take into account the new command.

[Ex. 2] Let $\mathcal{D} = (D, \sqsubseteq)$ be a CPO. Let us consider the pair $(\wp(D), \preceq)$, where

$$S \preceq S' \quad \Leftrightarrow \quad \forall s \in S. \; \exists s' \in S'. \; s \sqsubseteq s'$$

Is $(\wp(D), \preceq)$ a CPO_⊥?

[Ex. 3] Let us consider the following definition

$$f(x) \stackrel{\text{def}}{=} \mathbf{if} \ x = 0 \mathbf{then} \ 1 \mathbf{ else} \ (1 + f(x - 1))$$

- 1. Define a (recursive) HOFL term t that corresponds to the above definition and prove that it has type $int \rightarrow int$.
- 2. Prove that for any $n \ge 0$ it holds $(t \ n) \rightarrow n+1$ using the eager operational semantics.
- [Ex. 4] Let us consider the CTMC

$$s_{0} \xrightarrow{\lambda_{2}} s_{1}$$

$$\lambda_{1} \left(\begin{array}{c} \lambda_{1} \\ \lambda_{1} \\ s_{2} \end{array} \right) \xrightarrow{\lambda_{1}} s_{3}$$

- 1. What is the probability to sojourn in s_0 for some time r?
- 2. Assume $\lambda_2 > 2\lambda_1$: are there some bisimilar states?