Principles for software composition 2021/22 05 - HOFL

[Ex. 1] Determine the type of the HOFL term

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ x. \ ((\lambda y. \ \mathbf{if} \ y \ \mathbf{then} \ 0 \ \mathbf{else} \ 0) \ x).$$

Then compute its (lazy) canonical form.

[Ex. 2] Determine the type of the HOFL term

$$map \stackrel{\text{def}}{=} \lambda f. \ \lambda x. \ ((f \ \mathbf{fst}(x)), (f \ \mathbf{snd}(x)))$$

Then, compute the (lazy) canonical forms of the terms

$$t_1 \stackrel{\text{def}}{=} map \ (\lambda z. \ 2 \times z) \ (1,2)$$
 $t_2 \stackrel{\text{def}}{=} \mathbf{fst} \ (map \ (\lambda z. \ 2 \times z) \ (1,2))$

[Ex. 3] Let (D, \sqsubseteq_D) be a CPO and $f: D \to D$ be a continuous function. Prove that the set of fixpoints of f is itself a CPO (ordered by \sqsubseteq_D).

[Ex. 4] (Test for convergence) We would like to modify the denotational semantics of HOFL assigning to the construct

if
$$t$$
 then t_0 else t_1

- the semantics of t_1 if the semantics of t is $\perp_{\mathbb{Z}_{\perp}}$, and
- the semantics of t_0 otherwise.

Is it possible? If not, why?

[Ex. 5] (Strict conditional) Modify the operational semantics of HOFL by taking the following rules for conditionals:

$$\frac{t \to 0 \quad t_0 \to c_0 \quad t_1 \to c_1}{\text{if } t \text{ then } t_0 \text{ else } t_1 \ \to \ c_0} \qquad \qquad \frac{t \to n \quad n \neq 0 \quad t_0 \to c_0 \quad t_1 \to c_1}{\text{if } t \text{ then } t_0 \text{ else } t_1 \ \to \ c_1}.$$

Without changing the denotational semantics, prove that:

- 1. for any term t and canonical form c, we have $t \to c \implies \forall \rho$. $[t] \rho = [c] \rho$;
- 2. in general $t \Downarrow \not\Rightarrow t \downarrow$ (exhibit a counterexample).

[Ex. 6] Determine the type of the HOFL term

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ f. \ (\lambda x.1, \ \mathbf{fst}(f) \ 0)$$

Then, compute the (lazy) denotational semantics of t.