

**PSC 2022/23** (375AA, 9CFU)

Principles for Software Composition

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08a - Complete Partial Orders

# Partial orders

# Partially ordered set

## (Poset or just PO)

a set

a binary relation

$$(P, \sqsubseteq) \quad \sqsubseteq \subseteq P \times P$$

reflexive

$$\forall p \in P.$$

$$p \sqsubseteq p$$

antisymmetric

$$\forall p, q \in P.$$

$$p \sqsubseteq q \wedge q \sqsubseteq p \Rightarrow p = q$$

transitive

$$\forall p, q, r \in P.$$

$$p \sqsubseteq q \wedge q \sqsubseteq r \Rightarrow p \sqsubseteq r$$

$q$



$p$

$$p \sqsubseteq q$$

means that  $p$  and  $q$  are **comparable** and that  $p$  is less than (or equal to)  $q$

$$p \sqsubset q$$

means  $p \sqsubseteq q \wedge p \neq q$

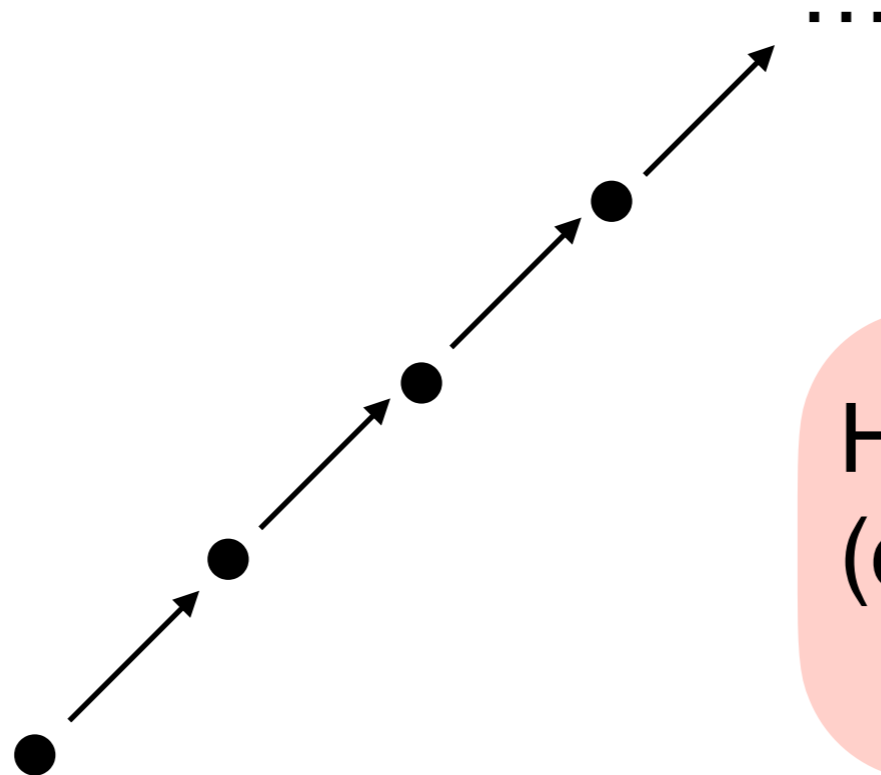
# Total orders

$(P, \sqsubseteq)$  PO

total

$$\forall p, q \in P. \quad p \sqsubseteq q \vee q \sqsubseteq p$$

a PO where every two elements are **comparable**



Hasse diagram notation  
(omit: reflexive arcs,  
transitive arcs)

# Discrete orders

$(P, \sqsubseteq)$  PO

discrete

$$\forall p, q \in P. \quad p \sqsubseteq q \Leftrightarrow p = q$$

any element is **comparable** only to itself

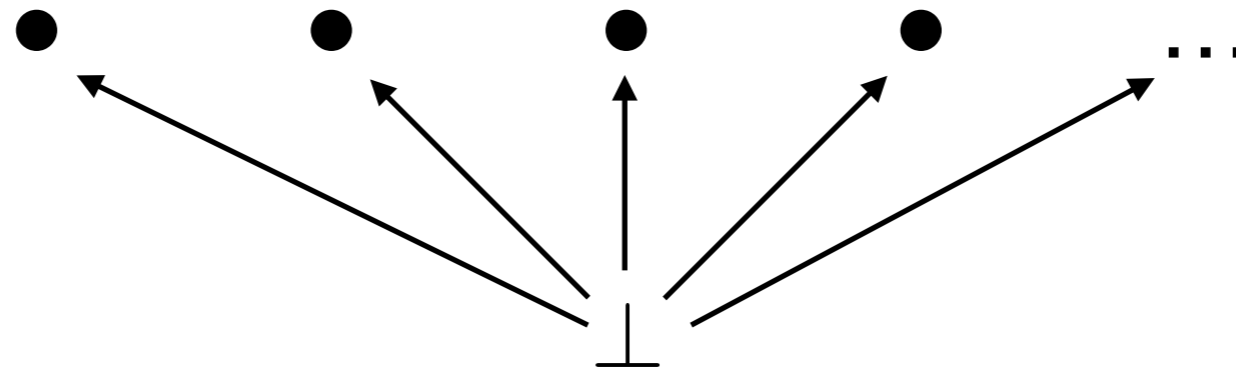


# Flat orders

$(P, \sqsubseteq)$  PO

flat  $\forall p, q \in P. \quad p \sqsubseteq q \Leftrightarrow p = q \vee p = \perp$

any element is **comparable** only to itself  
and with a distinguished (smaller) element  $\perp$





# Exercise

$(\mathbb{N}, \leq)$

PO?

Total?

Discrete?

Flat?



...



3



2



1



0



# Exercise

$(\wp(S), \subseteq)$

PO?



Total?

$$|S| < 2$$

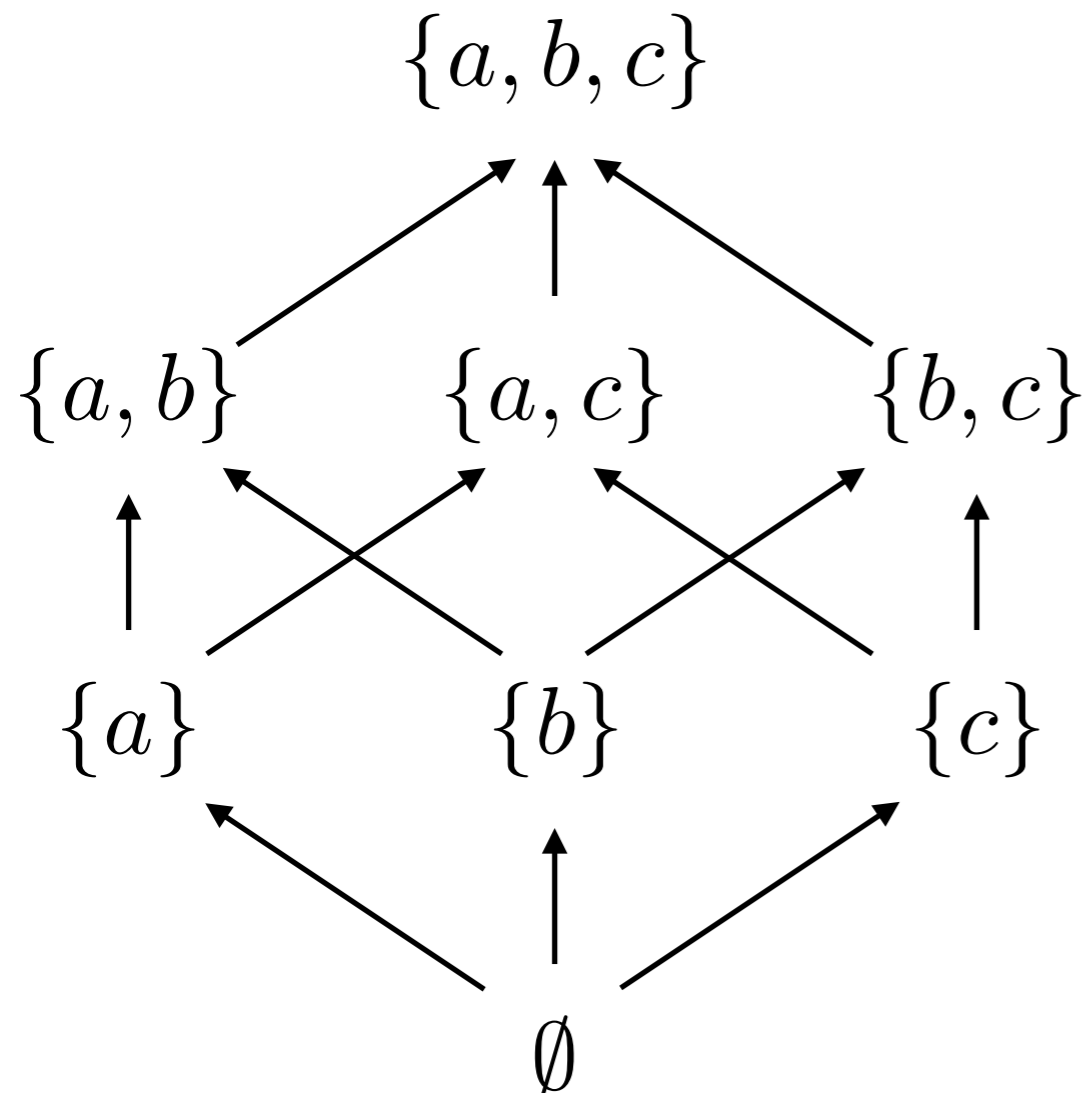
Discrete?

$$S = \emptyset$$

Flat?

$$|S| < 2$$

example:  $S = \{a, b, c\}$



$$\{a, b\} \not\subseteq \{b, c\}$$

$$\{a\} \not\subseteq \{b\}$$





# Exercise

$(\mathbb{N}, =)$

PO?

Total?

Discrete?

Flat?



0

1

2

3

...



# Exercise

$(\mathbb{N} \cup \{\perp\}, \{(\perp, n) \mid n \in \mathbb{N}\}^*)$

PO?



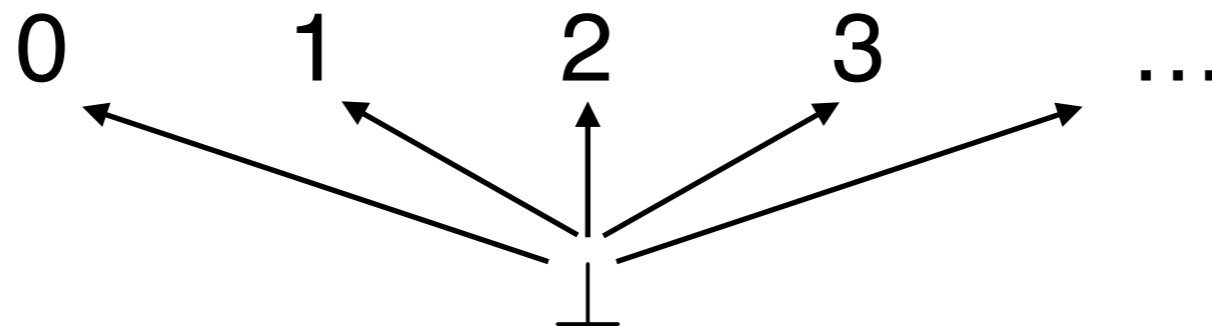
Total?



Discrete?



Flat?





# Exercise

$(\mathbb{N} \cup \{\infty\}, \leq)$

PO?



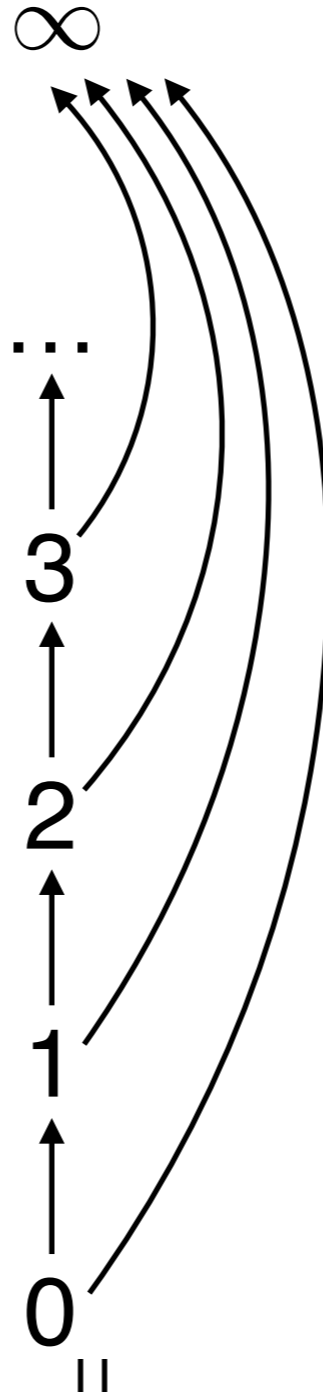
Total?



Discrete?



Flat?





# Exercise

	PO?	Total?	Discrete?	Flat?
$(\mathbb{N}, <)$	✗	✗	✗	✗
$(\mathbb{Z}, \leq)$	✓	✓	✗	✗
$(\mathbb{Z} \cup \{-\infty, \infty\}, \leq)$	✓	✓	✗	✗
$(T_\Sigma, \prec)$	✗	✗	✗	✗
$(\mathbb{N}, \neq)$	✗	✗	✗	✗

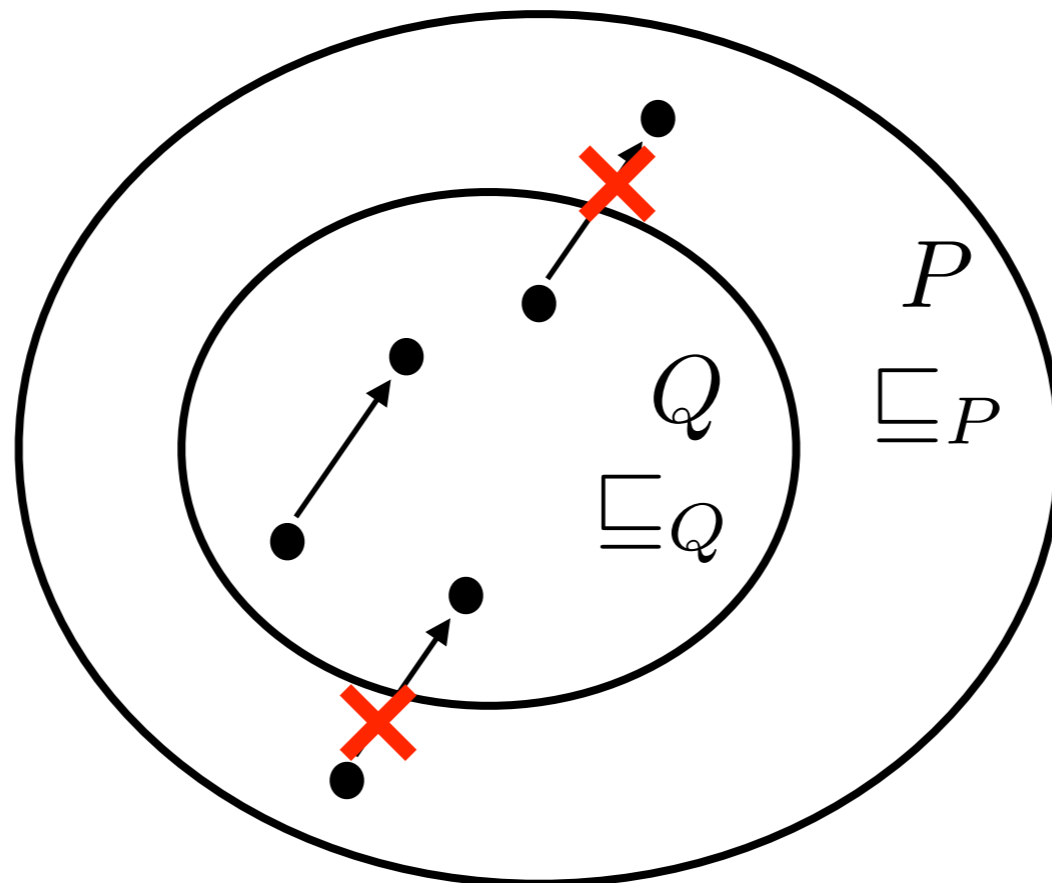
# Subset of a PO

$(P, \sqsubseteq_P)$  PO     $Q \subseteq P$

let  $\sqsubseteq_Q \triangleq \sqsubseteq_P \cap (Q \times Q)$

TH.  $(Q, \sqsubseteq_Q)$  is a PO

TH. if  $(P, \sqsubseteq_P)$  is total, then  $(Q, \sqsubseteq_Q)$  is total



PO  $\sqsubseteq$

w.f.  $\prec$

reflexive

not reflexive (otherwise cycle!)

antisymmetric

antisymmetric (otherwise cycle!)  
 $p \prec q \wedge q \prec p$  is always false

transitive

can be transitive ( $\prec^+$  w.f.)

has infinite descending chains  
(if nonempty)

no infinite descending chain

$\prec^*$  is always a PO

$\sqsubseteq$  can be w.f.

# Element properties (least, minimal, ...)

# Least element

$(P, \sqsubseteq)$  PO  $Q \subseteq P$   $l \in Q$

$l$  is a **least** element of  $Q$  if  $\forall q \in Q. l \sqsubseteq q$

TH. (uniqueness of least element)

$(P, \sqsubseteq)$  PO  $Q \subseteq P$   $l_1, l_2$  least elements of  $Q$  implies  $l_1 = l_2$

$$\left. \begin{array}{l} l_1 \text{ least element of } Q \Rightarrow l_1 \sqsubseteq l_2 \\ l_2 \text{ least element of } Q \Rightarrow l_2 \sqsubseteq l_1 \end{array} \right\} \Rightarrow l_1 = l_2$$

by antisymmetry




# Bottom

$(P, \sqsubseteq)$  PO the least element of  $P$   
(if it exists) is called **bottom** and denoted  $\perp$

sometimes written  $\perp_P$

## Examples

PO	bottom?
$(\mathbb{N} \cup \{\infty\}, \leq)$	0
$(\wp(S), \subseteq)$	$\emptyset$
$(\mathbb{Z}, \leq)$	

# Minimal element

$(P, \sqsubseteq)$  PO     $Q \subseteq P$      $m \in Q$

$m$  is a **minimal** element of  $Q$  if  $\forall q \in Q. q \sqsubseteq m \Rightarrow q = m$

(no smaller element can be found in  $Q$ )

**least**  $\forall q \in Q. l \sqsubseteq q$

**minimal**  $\forall q \in Q. q \sqsubseteq m \Rightarrow q = m$

unique

not necessarily unique

minimal

not necessarily least  
can be least

# Reverse order

TH.  $(P, \sqsubseteq)$  PO implies  $(P, \sqsupseteq)$  PO

note:

$$\sqsupseteq = \sqsubseteq^{-1}$$

proof. it is immediate to check that  $\sqsupseteq$

is reflexive

is antisymmetric

is transitive

$(P, \sqsubseteq)$  PO  $Q \subseteq P$

**greatest** element: least element of  $Q$  w.r.t.  $(P, \sqsupseteq)$

**top** element:  $\top$  greatest element of  $P$  (if it exists)

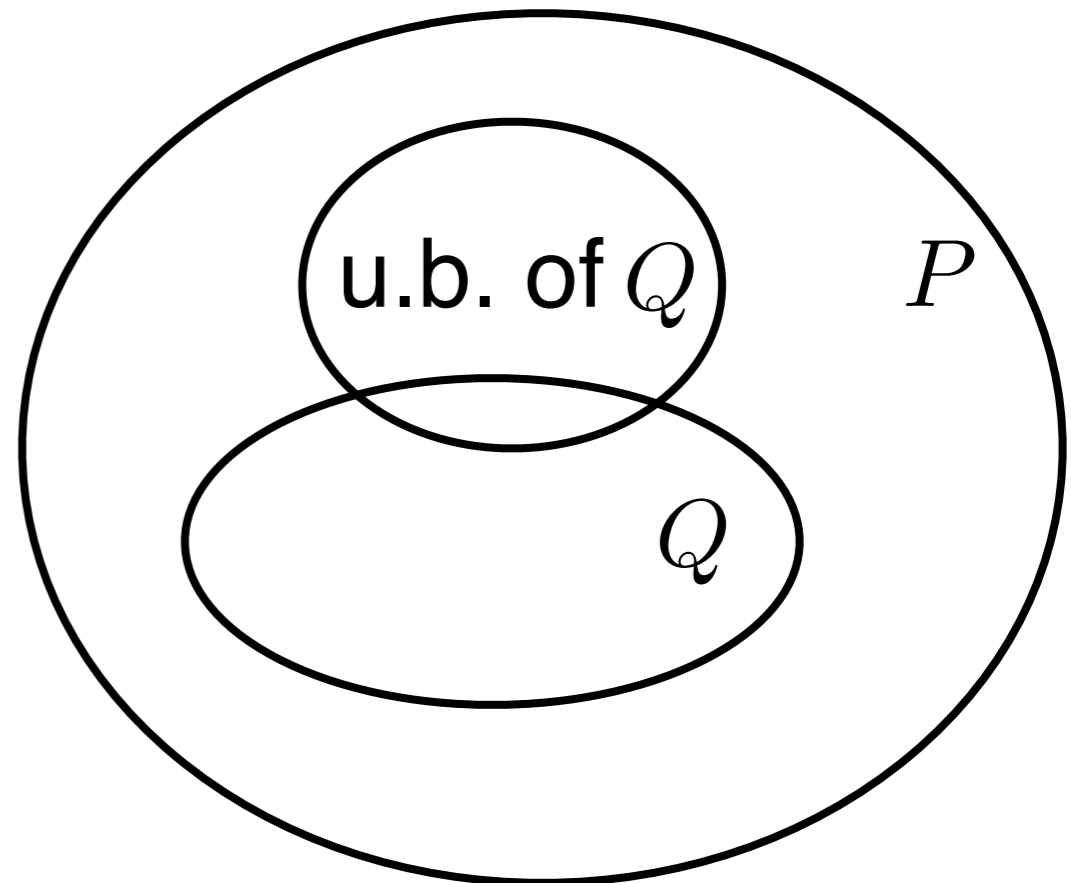
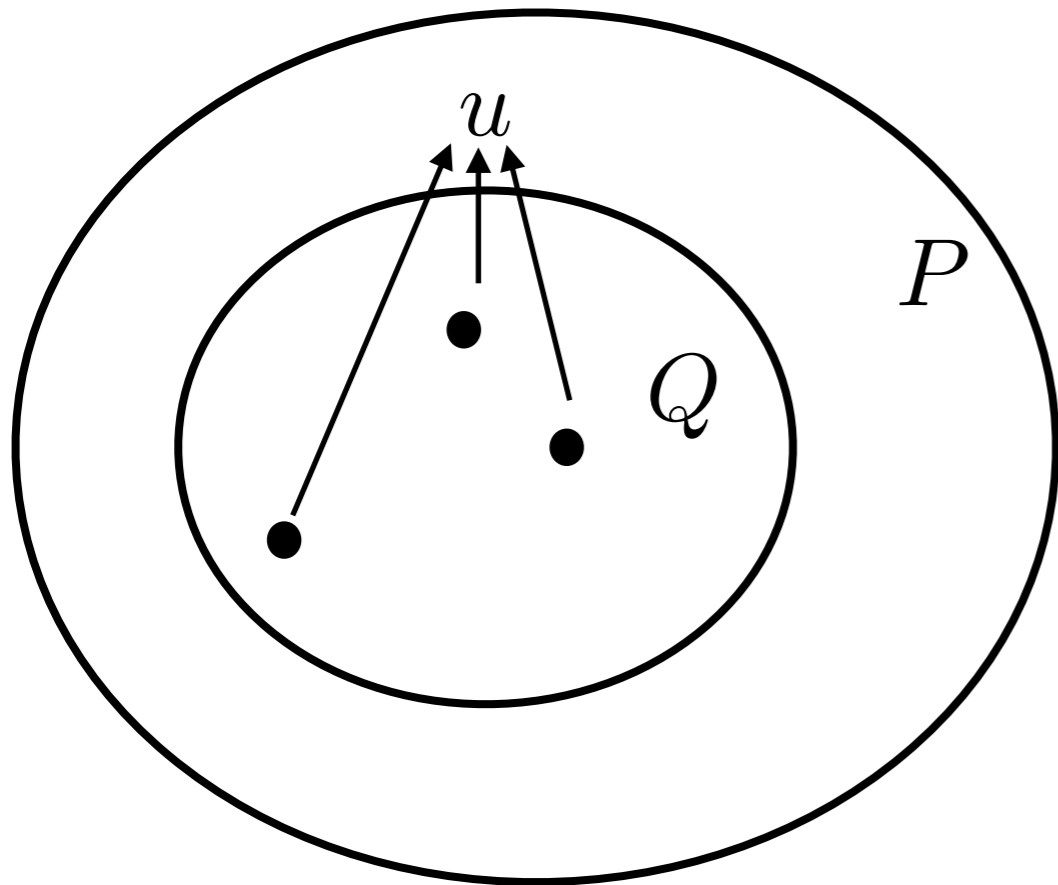
**maximal** element: minimal element of  $Q$  w.r.t.  $(P, \sqsupseteq)$

# Upper bound

$(P, \sqsubseteq)$  PO     $Q \subseteq P$      $u \in P$

$u$  is an **upper bound** of  $Q$  if  $\forall q \in Q. q \sqsubseteq u$   
(all the elements of  $Q$  are smaller than  $u$ )

$Q$  may have many upper bounds



# Least upper bound

$(P, \sqsubseteq)$  PO  $Q \subseteq P$   $p \in P$

$p$  is the **least upper bound (lub)** of  $Q$  if

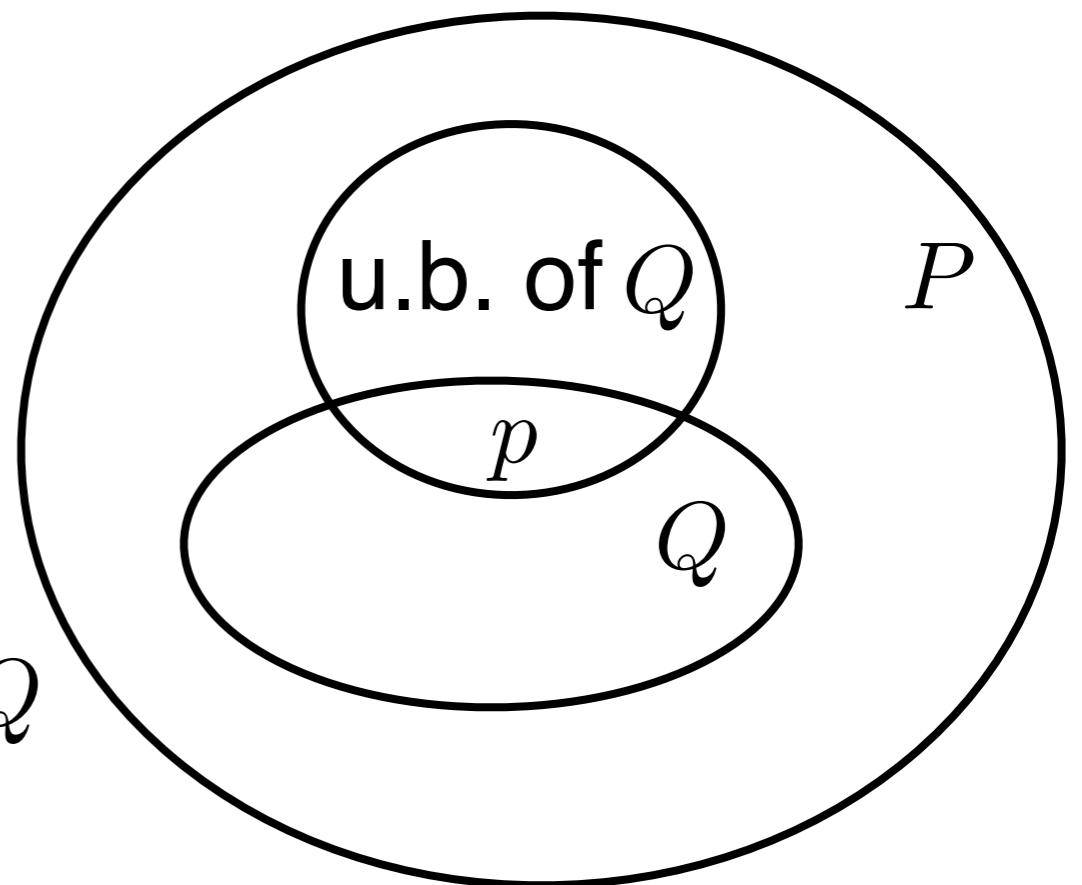
1. it is an upper bound of  $Q$   $\forall q \in Q. q \sqsubseteq p$
2. it is smaller than any other upper bound of  $Q$

$$\forall u \in P. (\forall q \in Q. q \sqsubseteq u) \Rightarrow p \sqsubseteq u$$

we write  $p = \text{lub } Q$

intuitively, it is the least element that represents all of  $Q$

$p$  not necessarily an element of  $Q$

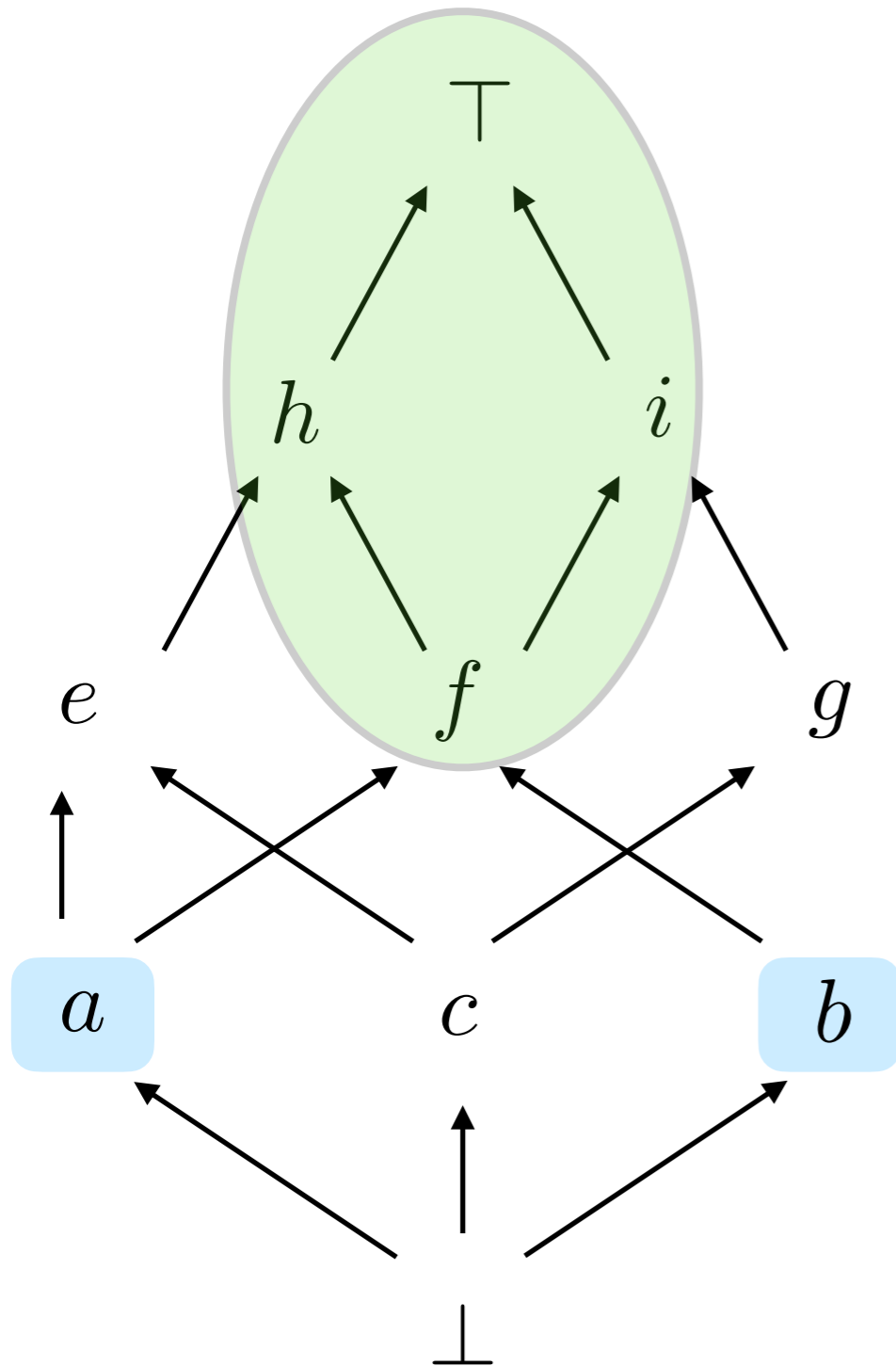




# Exercise

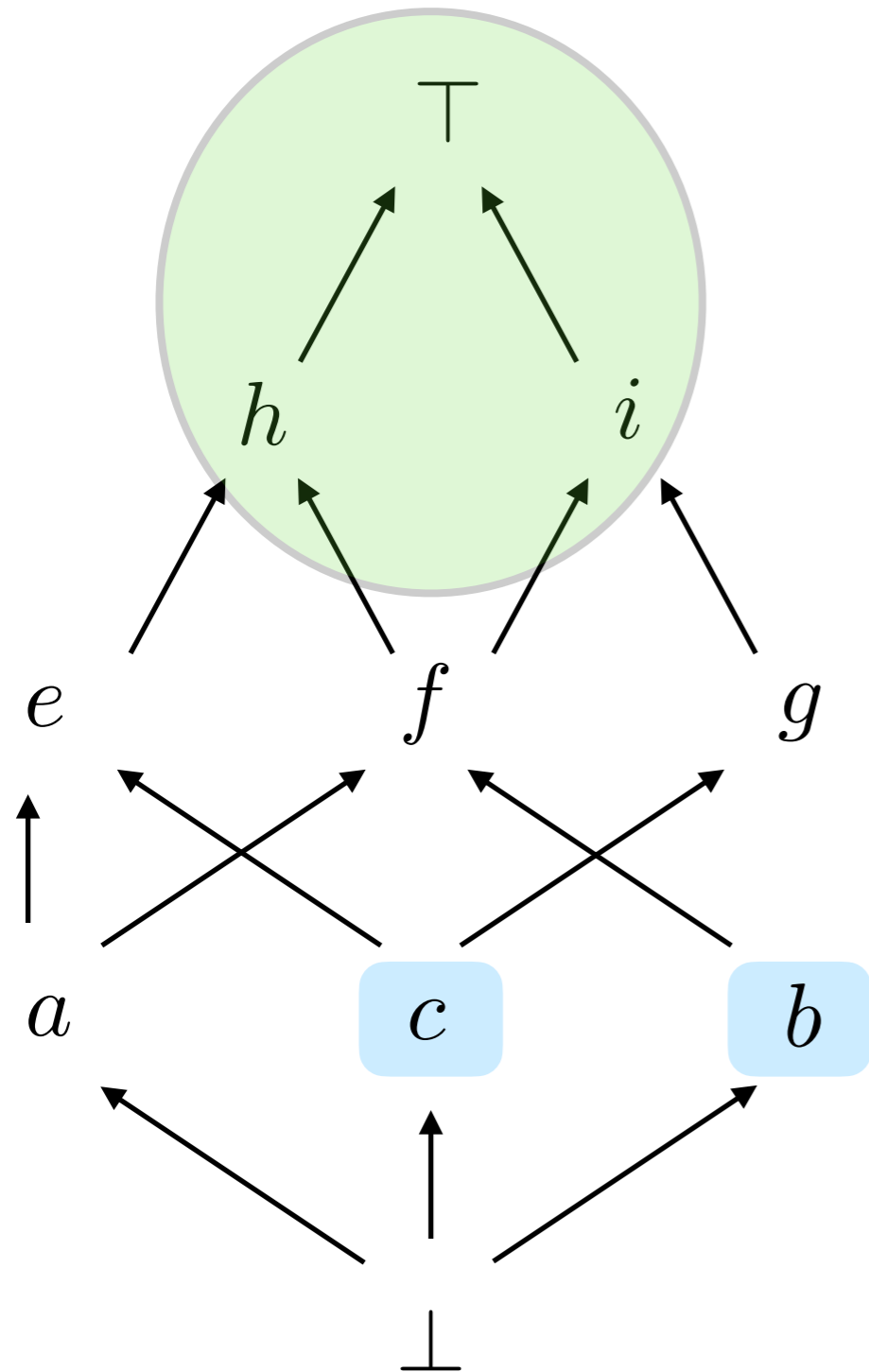
Upper bounds of  $\{a, b\}$  ?  $\{f, h, i, \top\}$

lub?  $f$





# Exercise



Upper bounds of  $\{b, c\}$  ?  $\{h, i, \top\}$

lub? no lub!



# Exercise

$(\mathbb{N}, \leq)$

$Q \subseteq \mathbb{N}$

lub?

if  $Q$  finite  $\text{lub } Q = \max Q$   
otherwise no lub



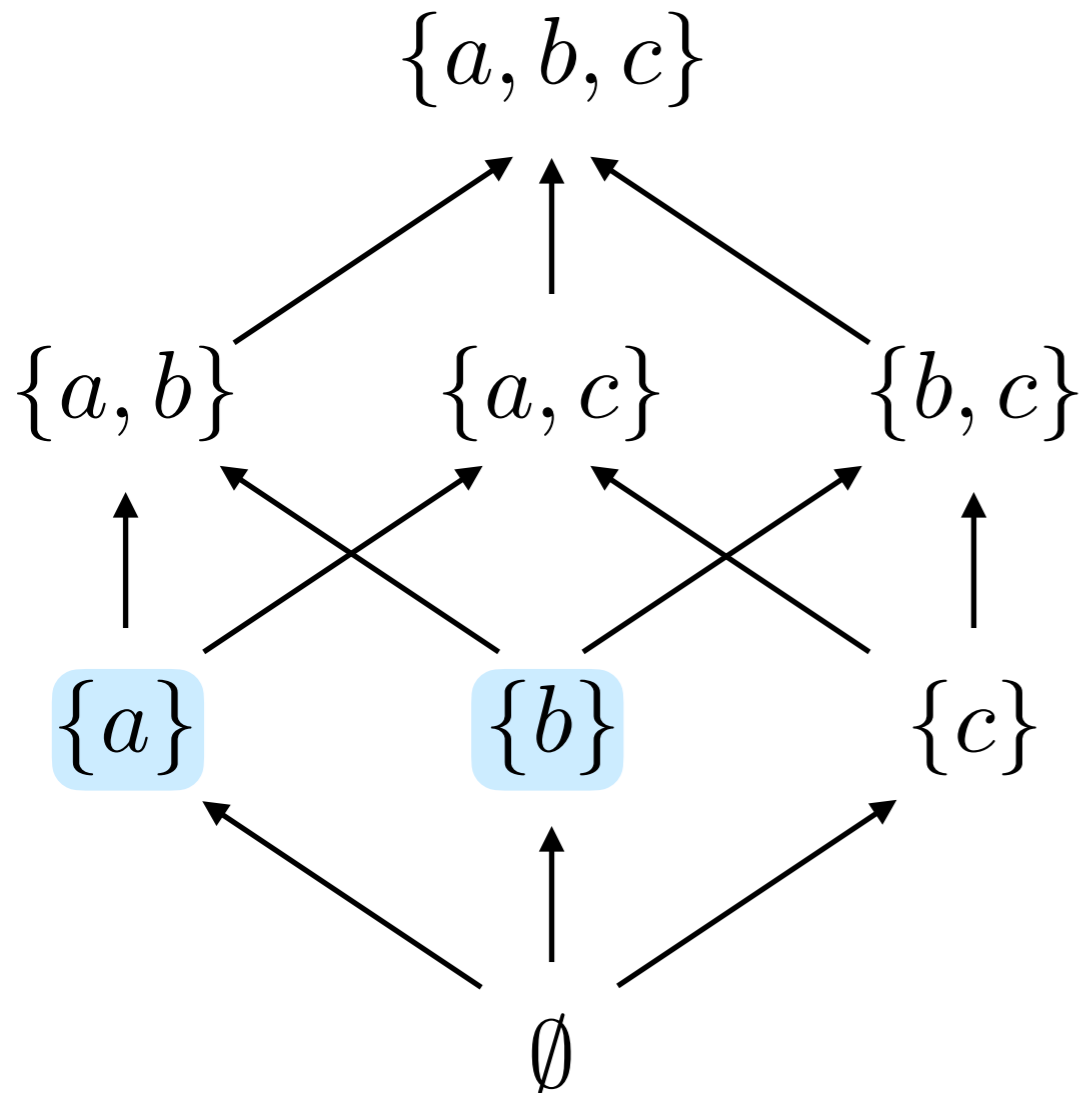




# Exercise

$(\wp(S), \subseteq)$   $Q \subseteq \wp(S)$  lub?

$$\text{lub } Q = \bigcup_{T \in Q} T$$



$$\text{lub } \{\{a\}, \{b\}\} = \{a, b\}$$

# Complete partial orders (CPO)

# Completeness: the idea

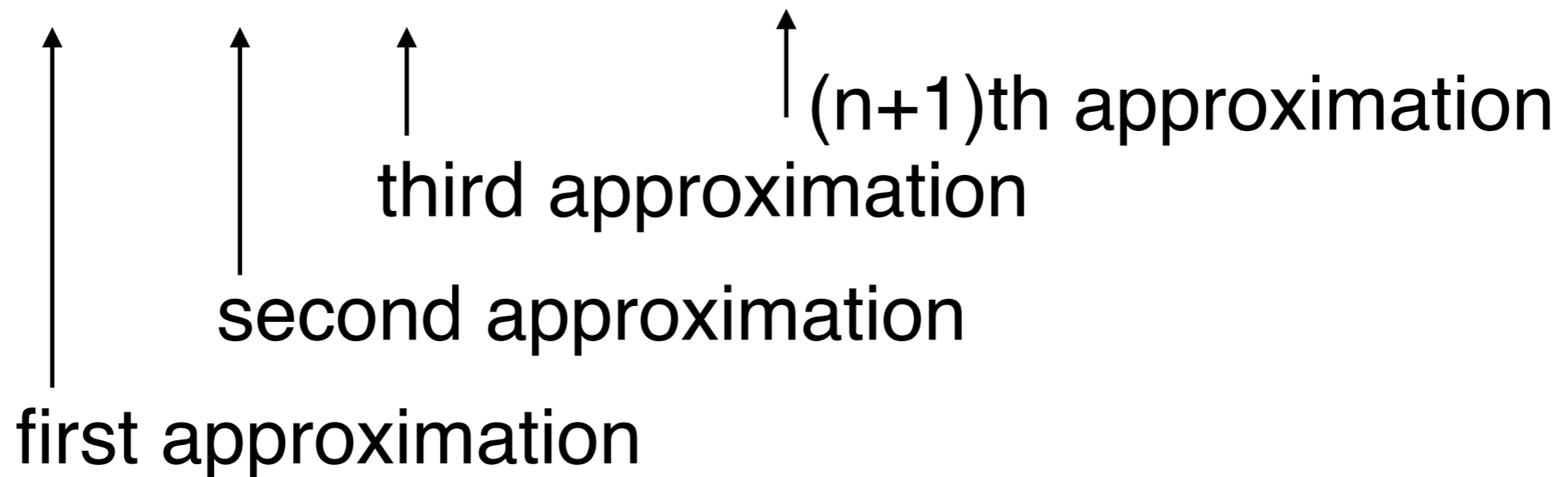
$D$  a domain

$\sqsubseteq$  a way to compare element

$x \sqsubseteq y$   $x$  is a (less precise) approximation of  $y$   
 $x$  and  $y$  are consistent,  
but  $y$  is more accurate than  $x$

} PO

$x_0 \sqsubseteq x_1 \sqsubseteq x_2 \sqsubseteq \dots \sqsubseteq x_n \sqsubseteq \dots$



does any sequence of approximations tend to some limit?

# Chain

$(P, \sqsubseteq)$  PO  $\{d_i\}_{i \in \mathbb{N}}$  is a **chain** if  $\forall i \in \mathbb{N}. d_i \sqsubseteq d_{i+1}$

$$d_0 \sqsubseteq d_1 \sqsubseteq d_2 \sqsubseteq \cdots \sqsubseteq d_n \sqsubseteq \cdots$$

any chain is an infinite list

**finite** chain: there are only finitely many distinct elements

$$\exists k \in \mathbb{N}. \forall i \geq k. d_i = d_{i+1}$$

or equivalently

$$\exists k \in \mathbb{N}. \forall i \geq k. d_i = d_k$$

# Example

$(\mathbb{N}, \leq)$

$0 \leq 2 \leq 4 \leq \dots \leq 2n \leq \dots$  is an infinite chain

$0 \leq 1 \leq 3 \leq 3 \leq 5 \leq \dots \leq 5 \leq \dots$  is a finite chain

any chain has infinite length

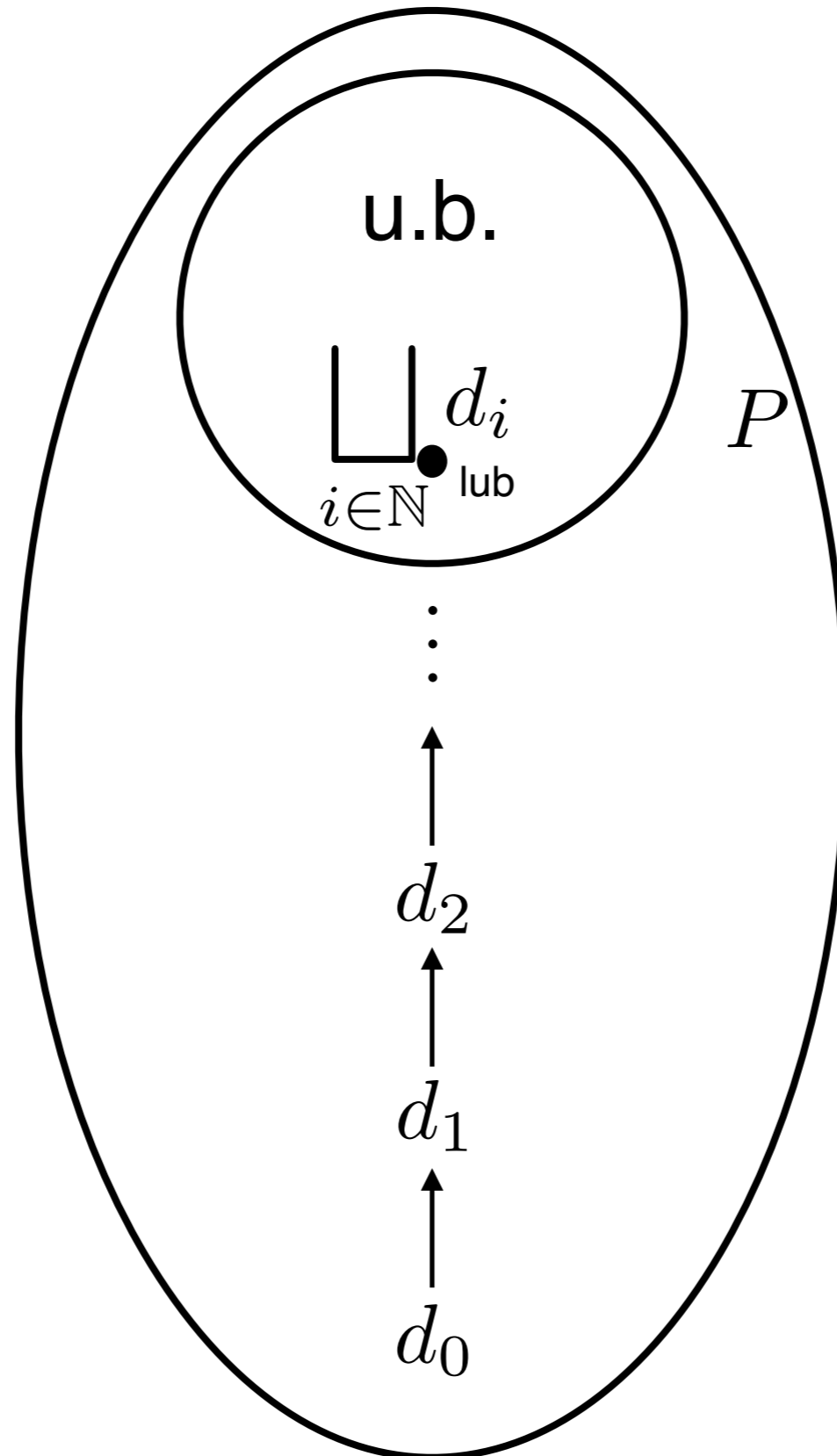
# Limit of a chain

$(P, \sqsubseteq)$  PO  $\{d_i\}_{i \in \mathbb{N}}$  a chain

we denote by  $\bigsqcup_{i \in \mathbb{N}} d_i$  the lub of  $\{d_i\}_{i \in \mathbb{N}}$  if it exists

and call it the **limit** of the chain

# Limit illustrated



# Example

$(\mathbb{N}, \leq)$

$0 \leq 2 \leq 4 \leq \dots \leq 2n \leq \dots$  has no lub  
(empty set of upper bounds)

$0 \leq 1 \leq 3 \leq 3 \leq 5 \leq \dots \leq 5 \leq \dots$  has lub 5  
(which upper bounds?)



# Lemma on finite chains

Lemma (any finite chain has a limit)

$(P, \sqsubseteq)$  PO  $\{d_i\}_{i \in \mathbb{N}}$  a finite chain  $\Rightarrow \bigsqcup_{i \in \mathbb{N}} d_i$  exists

proof.

$\{d_i\}_{i \in \mathbb{N}}$  finite  $\Rightarrow \exists k. \forall i. d_{i+k} = d_k$

the elements of the chain are totally ordered

$d_k$  is the greatest element of the chain

$d_k$  is an upper bound  $\forall i. d_i \sqsubseteq d_k$

$d_k$  is the least upper bound

take  $u$  such that  $\forall i. d_i \sqsubseteq u$  then  $d_k \sqsubseteq u$

# Prefix independence

Lemma (prefix independence)  $(P, \sqsubseteq)$  PO  $\{d_i\}_{i \in \mathbb{N}}$  a chain

$$\text{if } \bigsqcup_{i \in \mathbb{N}} d_i \text{ exists } \Rightarrow \forall k. \bigsqcup_{i \in \mathbb{N}} d_{i+k} = \bigsqcup_{i \in \mathbb{N}} d_i$$

$$\begin{array}{ccc} d_0 \sqsubseteq d_1 \sqsubseteq d_2 \sqsubseteq \cdots \sqsubseteq d_k \sqsubseteq d_{k+1} \sqsubseteq \cdots & \bigsqcup_{i \in \mathbb{N}} d_i & \\ & = & \\ & & \bigsqcup_{i \in \mathbb{N}} d_{i+k} \end{array}$$

# Prefix independence

Lemma (prefix independence)  $(P, \sqsubseteq)$  PO  $\{d_i\}_{i \in \mathbb{N}}$  a chain

$$\text{if } \bigsqcup_{i \in \mathbb{N}} d_i \text{ exists } \Rightarrow \forall k. \bigsqcup_{i \in \mathbb{N}} d_{i+k} = \bigsqcup_{i \in \mathbb{N}} d_i$$

proof.

take a generic  $k$

we prove that  $\{d_i\}_{i \in \mathbb{N}}$  and  $\{d_{i+k}\}_{i \in \mathbb{N}}$  have the same u.b.  
(and thus the same lub)

1. if  $u$  is an u.b. of  $\{d_i\}_{i \in \mathbb{N}}$  then is an u.b. of  $\{d_{i+k}\}_{i \in \mathbb{N}}$

because  $\{d_{i+k}\}_{i \in \mathbb{N}} \subseteq \{d_i\}_{i \in \mathbb{N}}$

2. if  $u$  is an u.b. of  $\{d_{i+k}\}_{i \in \mathbb{N}}$  we need to show  $\forall j. d_j \sqsubseteq u$   
for  $j \geq k$  it is obvious

if  $j < k$  then  $d_j \sqsubseteq d_k \sqsubseteq u$  because  $d_k \in \{d_{i+k}\}_{i \in \mathbb{N}}$

# Complete partial order

$(P, \sqsubseteq)$  PO      $P$  is **complete** if each chain has a limit (lub)

TH. Any finite chain has a limit  
(the last element in the sequence)

If  $P$  has only finite chains it is complete

If  $P$  is finite it is complete

Any discrete order is complete

Any flat order is complete

# Example

$(\mathbb{N}, \leq)$  is not complete  
(it is enough to exhibit a chain with no limit)

$0 \leq 2 \leq 4 \leq \dots \leq 2n \leq \dots$  has no lub  
(empty set of u.b.)

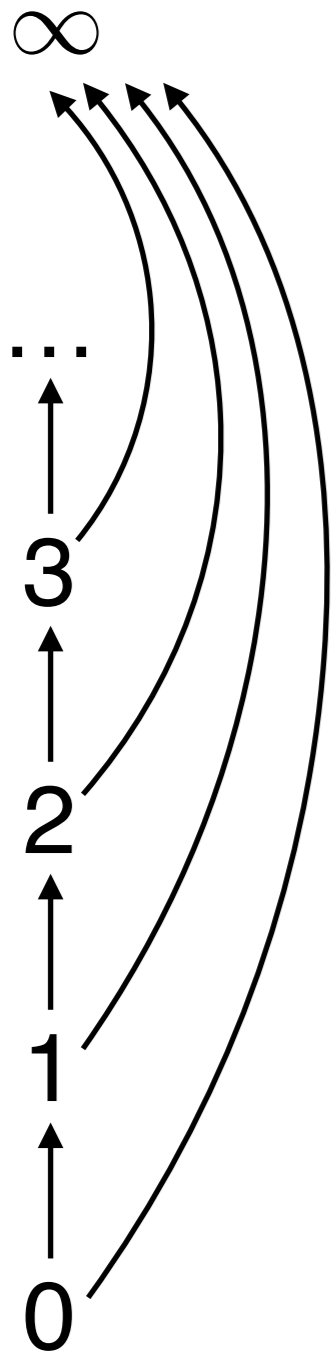


# Exercise

$(\mathbb{N} \cup \{\infty\}, \leq)$

complete? 

any infinite chain has limit  $\infty$   
(set of u.b.  $\{\infty\}$  )





# Exercise

$(\wp(S), \subseteq)$

complete? 

$\{S_i\}_{i \in \mathbb{N}}$

$$\bigsqcup_{i \in \mathbb{N}} S_i = \bigcup_{i \in \mathbb{N}} S_i = \{x \mid \exists k \in \mathbb{N}. x \in S_k\}$$



# Exercise

$(\mathbb{N} \cup \{\infty_1, \infty_2\}, \leq)$  complete? ✘

any infinite chain has no limit  
(set of u.b.  $\{\infty_1, \infty_2\}$  )

