



**PSC 2023/24** (375AA, 9CFU)

Principles for Software Composition

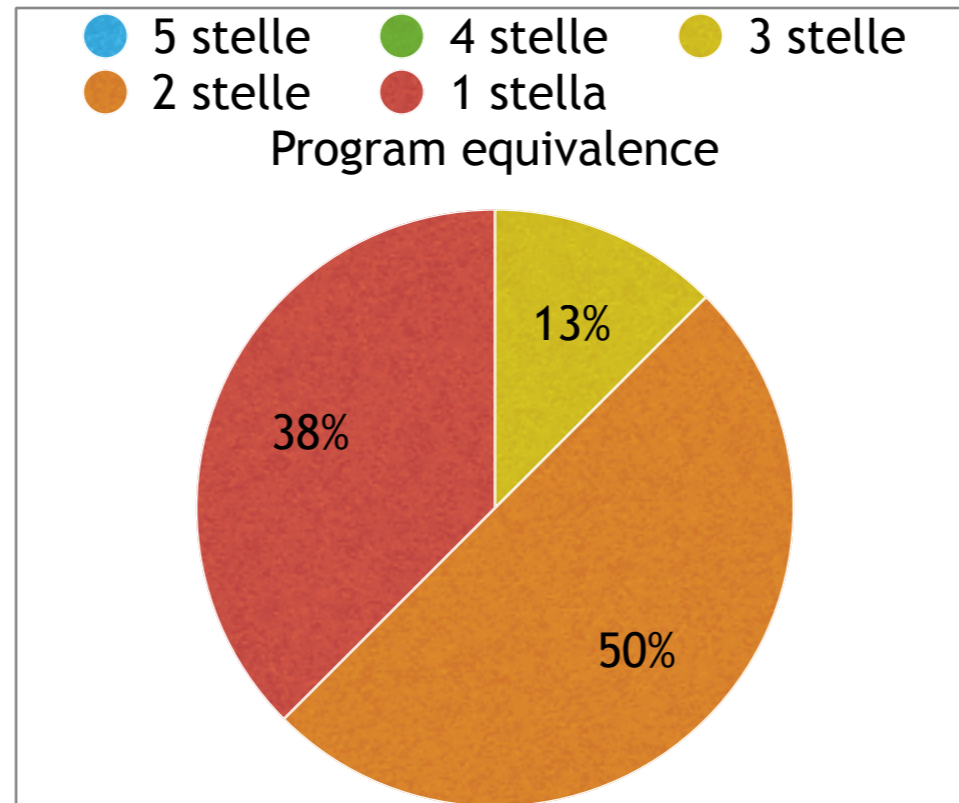
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## 06 - Operational equivalence

# From your forms



(over 8 answers)

# Operational equivalence

# Program equivalence

Symbolic manipulation of programs is possible  
if we can guarantee that their semantics is preserved

Useful to study program transformations  
(interpreters, compilers, code optimisation, refactoring...)

but... when are two commands equivalent?

# Operational equivalence

A first attempt

default initial memory

$$\forall x. \sigma_0(x) = 0$$

$$c_1 \sim_o c_2 \stackrel{\Delta}{=} \forall \sigma. ( \langle c_1, \sigma_0 \rangle \longrightarrow \sigma \Leftrightarrow \langle c_2, \sigma_0 \rangle \longrightarrow \sigma )$$

What's wrong with this definition?

$$x := y + 1 \sim_o x := 1$$

$$C[\bullet] \stackrel{\Delta}{=} y := 1; [\bullet]$$

$$C[x := y + 1] \stackrel{\Delta}{=} y := 1; x := y + 1$$

$$C[x := 1] \stackrel{\Delta}{=} y := 1; x := 1$$

$$C[x := y + 1] \not\sim_o C[x := 1]$$

# Operational equivalence

any initial memory

$$c_1 \sim_o c_2 \quad \triangleq \quad \forall \sigma, \sigma'. ( \langle c_1, \sigma \rangle \longrightarrow \sigma' \Leftrightarrow \langle c_2, \sigma \rangle \longrightarrow \sigma' )$$

much better:  
it is an equivalence relation  
and also a congruence  
(but we will see this later)

# Concrete equivalence

$$c \stackrel{?}{\sim}_o w$$

$$w \triangleq \text{while } x > 0 \text{ do } x := x - 1$$

$$c \triangleq \text{if } x > 0 \text{ then } x := 0 \text{ else skip}$$

Take a generic  $\sigma$  if  $\sigma(x) \leq 0$ :

$$\langle c, \sigma \rangle \longrightarrow \sigma'$$

$$\swarrow \langle x > 0, \sigma \rangle \longrightarrow \mathbf{ff}, \langle \mathbf{skip}, \sigma \rangle \longrightarrow \sigma'$$

$$\swarrow^* \langle \mathbf{skip}, \sigma \rangle \longrightarrow \sigma'$$

$$\swarrow_{\sigma' = \sigma} \square$$

$$\langle c, \sigma \rangle \longrightarrow \sigma$$

$$\langle w, \sigma \rangle \longrightarrow \sigma'$$

$$\swarrow_{\sigma' = \sigma} \langle x > 0, \sigma \rangle \longrightarrow \mathbf{ff}$$

$$\swarrow^* \square$$

$$\langle w, \sigma \rangle \longrightarrow \sigma$$

# Example (ctd)

$c \stackrel{?}{\sim}_o w$

$w \triangleq \text{while } x > 0 \text{ do } x := x - 1$

$c \triangleq \text{if } x > 0 \text{ then } x := 0 \text{ else skip}$

Take a generic  $\sigma$  if  $\sigma(x) > 0$ :

$\langle c, \sigma \rangle \longrightarrow \sigma'$

$\swarrow \langle x > 0, \sigma \rangle \longrightarrow \mathbf{tt}, \langle x := 0, \sigma \rangle \longrightarrow \sigma'$

$\swarrow^* \langle x := 0, \sigma \rangle \longrightarrow \sigma'$

$\swarrow^* \sigma' = \sigma[0/x] \quad \square$

$\langle c, \sigma \rangle \longrightarrow \sigma[0/x]$

$\forall n > 0. P(n) \quad P(n) \triangleq \forall \sigma. \sigma(x) = n \Rightarrow \langle w, \sigma \rangle \longrightarrow \sigma[0/x]$



# Example (ctd)

$c \stackrel{?}{\sim}_o w$

$w \triangleq \text{while } x > 0 \text{ do } x := x - 1$

$c \triangleq \text{if } x > 0 \text{ then } x := 0 \text{ else skip}$

$\forall n > 0. P(n) \quad P(n) \triangleq \forall \sigma. \sigma(x) = n \Rightarrow \langle w, \sigma \rangle \longrightarrow \sigma[0/x]$

$P(1) \triangleq \forall \sigma. \sigma(x) = 1 \Rightarrow \langle w, \sigma \rangle \longrightarrow \sigma[0/x]$

take a generic  $\sigma$  and suppose  $\sigma(x) = 1$

$\langle w, \sigma \rangle \longrightarrow \sigma'$

$\swarrow \langle x > 0, \sigma \rangle \longrightarrow \mathbf{tt}, \langle x := x - 1, \sigma \rangle \longrightarrow \sigma'', \langle w, \sigma'' \rangle \longrightarrow \sigma'$

$\swarrow^* \langle x := x - 1, \sigma \rangle \longrightarrow \sigma'', \langle w, \sigma'' \rangle \longrightarrow \sigma'$

$\swarrow_{\sigma'' = \sigma[0/x]}^* \langle w, \sigma[0/x] \rangle \longrightarrow \sigma'$

$\swarrow_{\sigma' = \sigma[0/x]} \langle x > 0, \sigma[0/x] \rangle \longrightarrow \mathbf{ff}$

$\swarrow^* \square$

$\langle w, \sigma \rangle \longrightarrow \sigma[0/x]$

# Example (ctd)

$c \stackrel{?}{\sim}_o w$

$w \triangleq \text{while } x > 0 \text{ do } x := x - 1$

$c \triangleq \text{if } x > 0 \text{ then } x := 0 \text{ else skip}$

$\forall n > 0. P(n) \quad P(n) \triangleq \forall \sigma. \sigma(x) = n \Rightarrow \langle w, \sigma \rangle \longrightarrow \sigma[0/x]$

$\forall n > 0. P(n) \Rightarrow P(n+1) \quad \text{take a generic } n > 0$

we assume  $P(n) \triangleq \forall \sigma. \sigma(x) = n \Rightarrow \langle w, \sigma \rangle \longrightarrow \sigma[0/x]$

we prove  $P(n+1) \triangleq \forall \sigma. \sigma(x) = n+1 \Rightarrow \langle w, \sigma \rangle \longrightarrow \sigma[0/x]$

take a generic  $\sigma$  and suppose  $\sigma(x) = n+1$

$\langle w, \sigma \rangle \longrightarrow \sigma'$

$\swarrow \langle x > 0, \sigma \rangle \longrightarrow \mathbf{tt}, \langle x := x - 1, \sigma \rangle \longrightarrow \sigma'', \langle w, \sigma'' \rangle \longrightarrow \sigma'$

$\swarrow^* \langle x := x - 1, \sigma \rangle \longrightarrow \sigma'', \langle w, \sigma'' \rangle \longrightarrow \sigma'$

$\swarrow^*_{\sigma'' = \sigma[n/x]} \langle w, \sigma[n/x] \rangle \longrightarrow \sigma'$

by inductive hypothesis  $P(n)$  we know  $\langle w, \sigma[n/x] \rangle \longrightarrow \sigma[0/x]$

by determinacy it must be  $\sigma' = \sigma[0/x] \quad \langle w, \sigma \rangle \longrightarrow \sigma[0/x]$

# Parametric equivalence

$$w \stackrel{?}{\sim}_o ww \qquad w \triangleq \text{while } b \text{ do } c$$

$$ww \triangleq \text{while } b \text{ do } (\overbrace{\text{while } b \text{ do } c}^w)$$

Take a generic  $\sigma$  if  $\langle b, \sigma \rangle \longrightarrow \text{ff}$ :

$$\begin{array}{l} \langle w, \sigma \rangle \longrightarrow \sigma' \\ \swarrow_{\sigma' = \sigma} \langle b, \sigma \rangle \longrightarrow \text{ff} \\ \swarrow^* \square \end{array} \qquad \langle w, \sigma \rangle \longrightarrow \sigma$$
  

$$\begin{array}{l} \langle ww, \sigma \rangle \longrightarrow \sigma' \\ \swarrow_{\sigma' = \sigma} \langle b, \sigma \rangle \longrightarrow \text{ff} \\ \swarrow^* \square \end{array} \qquad \langle ww, \sigma \rangle \longrightarrow \sigma$$

# Example (ctd)

$$w \stackrel{?}{\sim}_o ww \qquad w \triangleq \text{while } b \text{ do } c$$

$$ww \triangleq \text{while } b \text{ do } (\overbrace{\text{while } b \text{ do } c}^w)$$

Take a generic  $\sigma$       if  $\langle b, \sigma \rangle \longrightarrow \mathbf{tt}$ :      if  $\langle w, \sigma \rangle \longrightarrow \sigma'$ :

$$\langle ww, \sigma \rangle \longrightarrow \sigma'_1$$

$$\swarrow \langle b, \sigma \rangle \longrightarrow \mathbf{tt}, \langle w, \sigma \rangle \longrightarrow \sigma'', \langle ww, \sigma'' \rangle \longrightarrow \sigma'_1$$

$$\swarrow^* \langle w, \sigma \rangle \longrightarrow \sigma'', \langle ww, \sigma'' \rangle \longrightarrow \sigma'_1$$

$$\swarrow_{\sigma'' = \sigma'}^* \langle ww, \sigma' \rangle \longrightarrow \sigma'_1$$

$$\swarrow_{\sigma'_1 = \sigma'} \langle b, \sigma' \rangle \longrightarrow \mathbf{ff}$$

$$\swarrow^* \square \qquad \langle ww, \sigma \rangle \longrightarrow \sigma'$$

see home exercises

$$\forall b, c, \sigma, \sigma'. \langle \text{while } b \text{ do } c, \sigma \rangle \longrightarrow \sigma' \Rightarrow \langle b, \sigma' \rangle \longrightarrow \mathbf{ff}$$

# Example (ctd)

$$w \stackrel{?}{\sim}_o ww \qquad w \triangleq \text{while } b \text{ do } c$$

$$ww \triangleq \text{while } b \text{ do } (\overbrace{\text{while } b \text{ do } c}^w)$$

Take a generic  $\sigma$       if  $\langle b, \sigma \rangle \longrightarrow \mathbf{tt}$ :      if  $\langle w, \sigma \rangle \not\rightarrow$ :

$$\langle ww, \sigma \rangle \longrightarrow \sigma'_1$$

$$\swarrow \langle b, \sigma \rangle \longrightarrow \mathbf{tt}, \langle w, \sigma \rangle \longrightarrow \sigma'', \langle ww, \sigma'' \rangle \longrightarrow \sigma'_1$$


$$\swarrow^* \langle w, \sigma \rangle \longrightarrow \sigma'', \langle ww, \sigma'' \rangle \longrightarrow \sigma'_1$$


will diverge, because by hypothesis  $\langle w, \sigma \rangle \not\rightarrow$


$$\langle ww, \sigma \rangle \not\rightarrow$$

# A note on divergence

any two (always) divergent commands are equivalent

$$w_1 \overset{?}{\sim}_o w_2$$


$$w_1 \overset{?}{\sim}_o w_3$$


$$w_1 \overset{?}{\sim}_o w_4$$


$w_1 \triangleq$  **while true do skip**

$w_2 \triangleq$   $x := 0$ ; **while**  $x \geq 0$  **do**  $x := x + 1$

$w_3 \triangleq$   $y := 0$ ; **while**  $y \leq 0$  **do**  $y := y - 100$

$w_4 \triangleq$   $y := x + 1$ ; **while**  $x \neq y$  **do** ( $x := x - 1$ ;  $y := y + 1$ )

# A difficult problem

$w_1 \triangleq$  while true do skip

$p \triangleq$   $x := 3;$   $w_1 \stackrel{?}{\sim}_o p$   
 $s := 1;$   
**while**  $x \neq s$  **do** (  
     $x := x + 2;$   
     $s := 1;$   
     $i := 2;$   
    **while**  $2 \times i \leq x$  **do** (  
        **if**  $x \% i = 0$  **then**  $s := s + i$   
        **else skip;**  
         $i := i + 1$   
    )  
)

# A difficult problem

$w_1 \triangleq$  while true do skip

$w_1 \stackrel{?}{\sim}_o p$

always terminate

when  $x \geq 2$ , at the end,  $s$  takes the sum of all proper divisors of  $x$

```
 $s := 1;$   
 $i := 2;$   
while  $2 \times i \leq x$  do (  
    if  $x \% i = 0$  then  $s := s + i$   
    else skip;  
     $i := i + 1$   
)
```



# A difficult problem

$w_1 \triangleq$  while true do skip

$w_1 \stackrel{?}{\sim}_o p$

$$s := \sum_{i \in \text{div}(x)} i$$

where  $\text{div}(x) \triangleq \{1\} \cup \{d \mid 1 < d < x, x \% d = 0\}$

$$i := 1 + (x/2)$$

# A difficult problem

$w_1 \triangleq$  while true do skip

$p \triangleq$   $x := 3;$  an odd number  $w_1 \stackrel{?}{\sim}_o p$   
 $s := 1;$  sum of proper divisors of  $x$   
**while**  $x \neq s$  **do** ( exit when  $x$  is perfect  
 $x := x + 2;$  next odd number

$s := \sum_{i \in \text{div}(x)} i$  sum of proper divisors of  $x$

where  $\text{div}(x) \triangleq \{1\} \cup \{d \mid 1 < d < x, x \% d = 0\}$

$i := 1 + (x/2)$

)

a number is perfect

if it equals the sum of its proper divisors

$$6 = 1 + 2 + 3$$

# A difficult problem

$w_1 \triangleq$  while true do skip

$p \triangleq$

terminate when the first **odd perfect number** is found

$w_1 \overset{?}{\sim}_o p$



as of 2022

it is not known if any odd perfect number exists

Ochem, Pascal; Rao, Michaël.

Odd perfect numbers are greater than  $10^{1500}$ .

Mathematics of Computation n.81(279): 1869–1877.

(2012) doi:10.1090/S0025-5718-2012-02563-4

a number is **perfect**

if it equals the sum of its proper divisors

# For the curious minds

Euclid proved that  $2^{p-1}(2^p - 1)$  is an even perfect number whenever  $2^p - 1$  is prime

Euler proved that any even perfect number has the form  $2^{p-1}(2^p - 1)$  with  $2^p - 1$  prime

Prime numbers of the form  $2^p - 1$  are now called [Mersenne primes](#)

Even perfect numbers and Mersenne primes are thus in 1-to-1 correspondence

For  $2^p - 1$  to be prime it is necessary that  $p$  is prime  
(1st badge exercise)

If  $p$  is prime,  $2^p - 1$  is not necessarily prime  
(1st badge exercise)

As of 2022, there are only 51 Mersenne primes known

It is not known if there are infinitely many Mersenne primes

# For the curious minds

| prime number $p$ | Mersenne prime $2^p - 1$ | even perfect number $2^{p-1}(2^p - 1)$ |
|------------------|--------------------------|----------------------------------------|
| 2                | 3                        | 6                                      |
| 3                | 7                        | 28                                     |
| 5                | 31                       | 496                                    |
| 7                | 127                      | 8.128                                  |
| 13               | 8.191                    | 33.550.336                             |
| 17               | 131.071                  | 8.589.869.056                          |
| ...              | ...                      | ...                                    |
| 82.589.933       | 24.862.048 digits!       | 49.724.095 digits!                     |

# A tricky problem

$w_1 \triangleq$  while true do skip

$p_1 \triangleq$   $x := 2;$

$s := 1;$

while  $x \neq s$  do (

$x := x + 2;$

$s := 1;$

$i := 2;$

while  $2 \times i \leq x$  do (

if  $x \% i = 0$  then  $s := s + i$

else skip;

$i := i + 1$

)

)

$w_1 \stackrel{?}{\sim}_o p_1$



$x := 6 \stackrel{?}{\sim}_o p_1$



$x := 6;$

$s := 6; \stackrel{?}{\sim}_o p_1$

$i := 4;$



# A variant

$w_1 \triangleq$  while true do skip

$p_2 \triangleq$   $s := 1;$

**while**  $x \neq s$  **do** (

$x := x + 2;$

$s := 1;$

$i := 2;$

**while**  $2 \times i \leq x$  **do** (

**if**  $x \% i = 0$  **then**  $s := s + i$

**else skip;**

$i := i + 1$

)

)

$w_1 \stackrel{?}{\sim}_o p_2$



$x := 6;$

$s := 6;$   $\stackrel{?}{\sim}_o p_2$

$i := 4;$



# A last instance

$p_2 \stackrel{?}{\sim}_o p_3$



```
 $p_2 \triangleq$   $s := 1;$   
  while  $x \neq s$  do (  
     $x := x + 2;$   
     $s := 1;$   
     $i := 2;$   
    while  $2 \times i \leq x$  do (  
      if  $x \% i = 0$  then  $s := s + i$   
        else skip;  
       $i := i + 1$   
    )  
  )
```

```
 $p_3 \triangleq$   $s := 1;$   
  while  $x \neq s$  do (  
     $x := x + 1;$   
     $s := 1;$   
     $i := 2;$   
    while  $2 \times i \leq x$  do (  
      if  $x \% i = 0$  then  $s := s + i$   
        else skip;  
       $i := i + 1$   
    )  
  )
```