

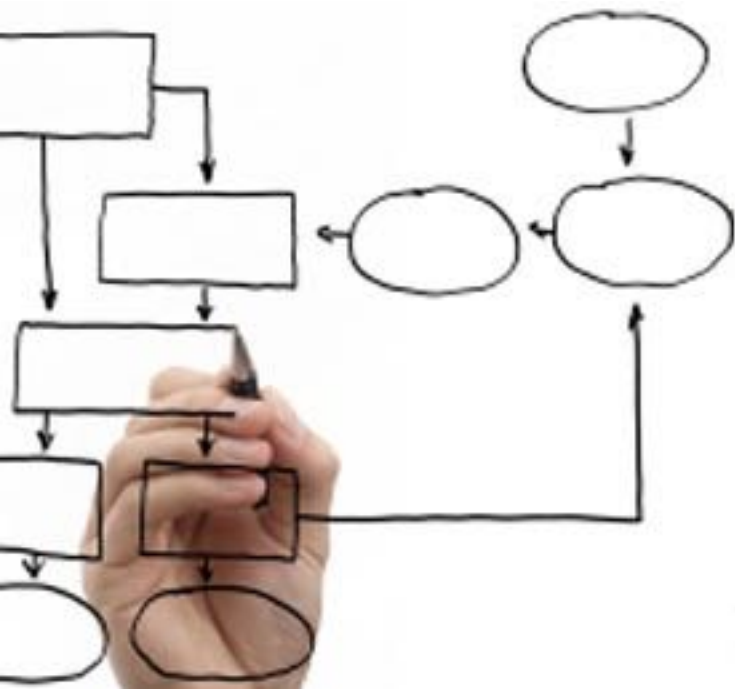
# Business Processes Modelling

## MPB (6 cfu, 295AA)

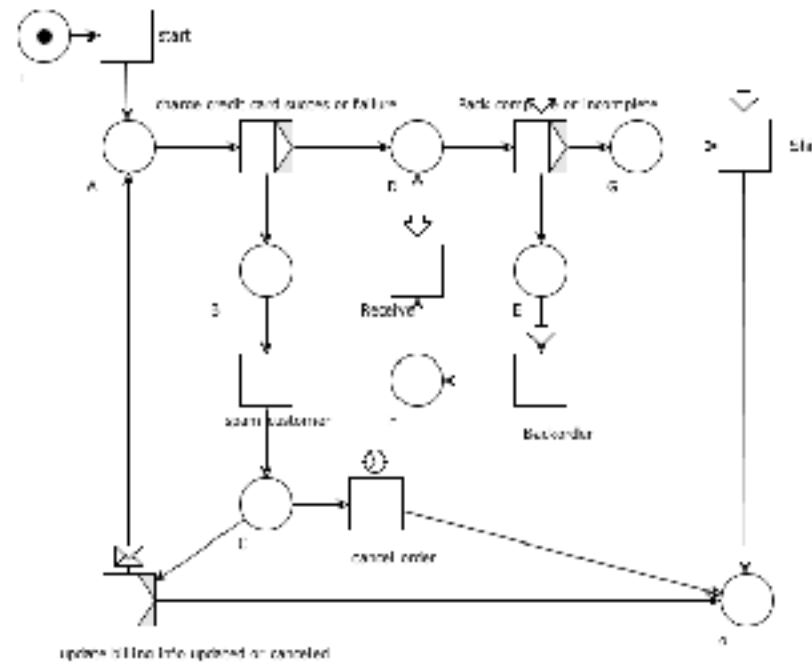
Roberto Bruni

<http://www.di.unipi.it/~bruni>

13 - Workflow nets



# Object



We study some special kind of Petri nets,  
that are suitable models of workflows

There are many, many  
variants of Petri nets

# Condition / Event Systems

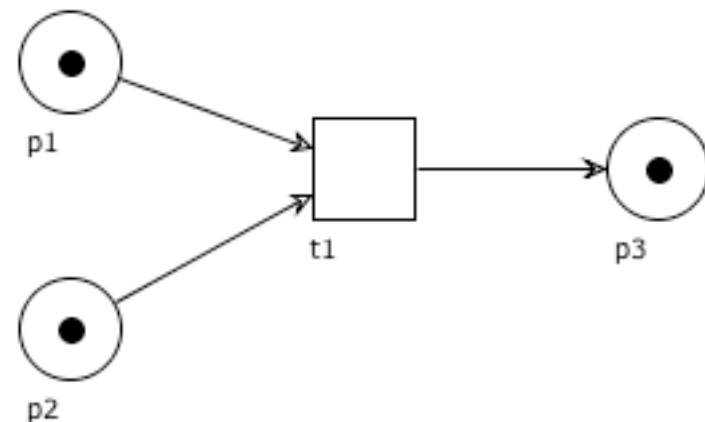
A **C/E system** is a Petri net whose places have all capacity equal to 1  
(i.e., each place can contain one token at most)

Markings are just subsets of  $P$  (not multisets)

Firing rule is more restrictive:

$t$  is enabled at  $M$  if  $\bullet t \subseteq M$  **and**  $t \bullet \cap M = \emptyset$

Is  $t_1$  enabled?



# Condition / Event Systems

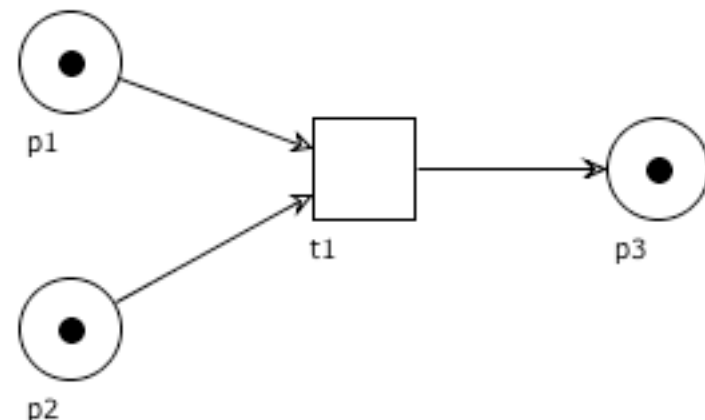
A **C/E system** is a Petri net whose places have all capacity equal to 1  
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Firing rule is more restrictive:

$t$  is enabled at  $M$  if  $\bullet t \subseteq M$  **and**  $t \bullet \cap M = \emptyset$

Is  $t_1$  enabled?



**No,**  
a token  
is already  
in  $p_3$

# Place / Transition Petri nets

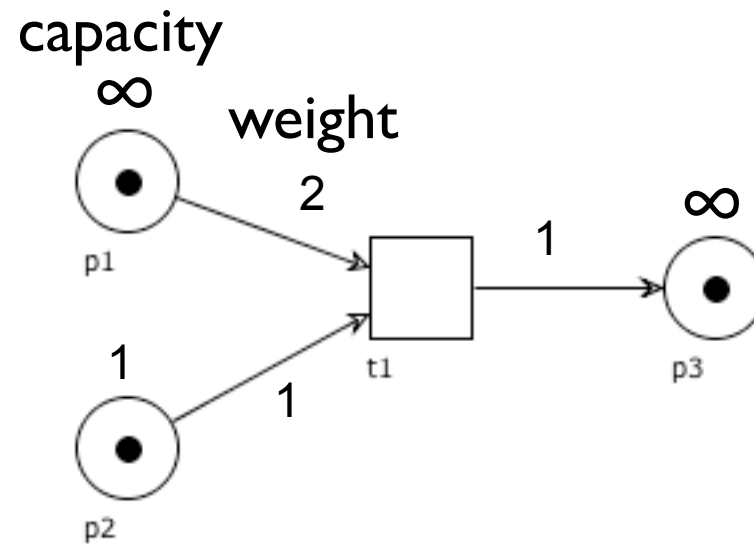
A **P/T net** is a Petri net  $(P, T, F)$  together with a weight function  $w : F \rightarrow \text{Nat}$

Firings consume and produce tokens according to the weight function

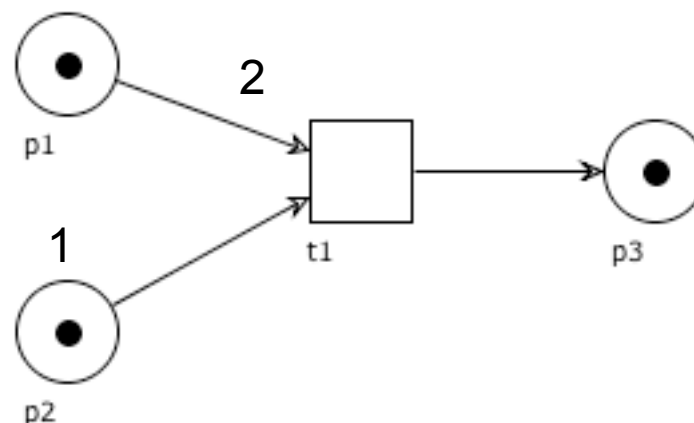
Sometimes a place capacity function  $c : P \rightarrow \text{Nat} \cup \{\infty\}$  is also considered

Firings cannot lead to markings where the capacity of a place is exceeded

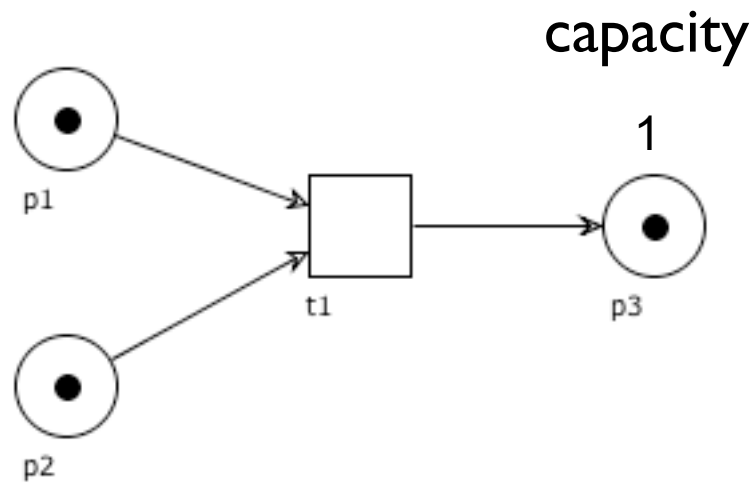
# P/T net: examples



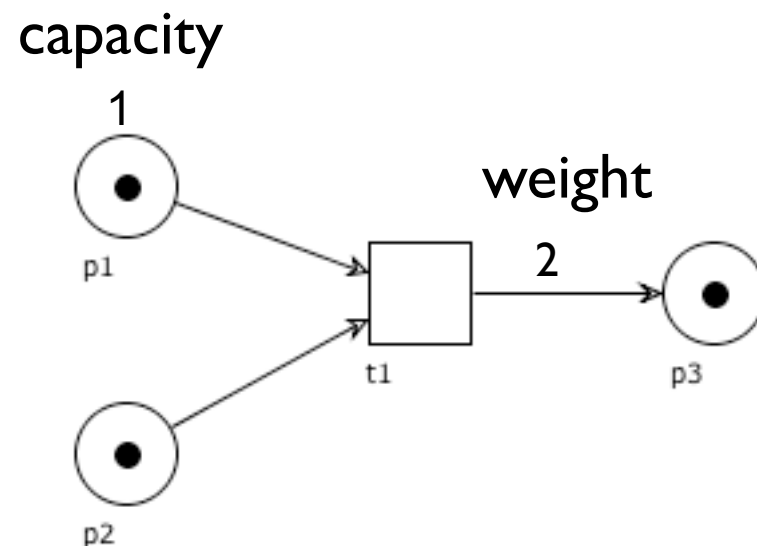
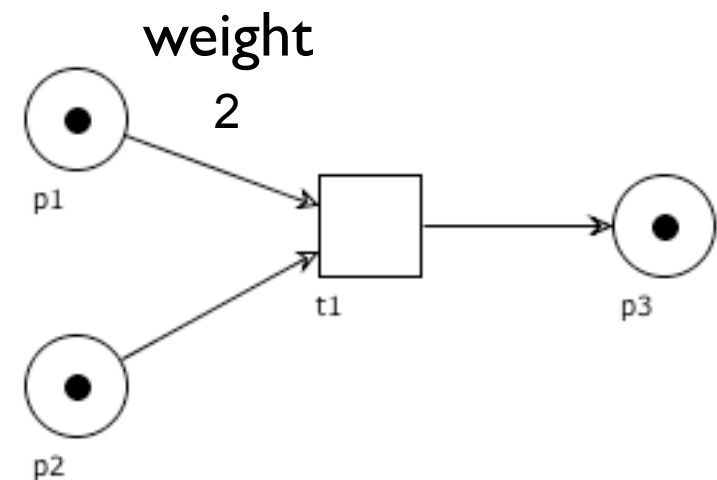
Capacity  $\infty$  is omitted from places  
Weight 1 is omitted from arcs



# P/T net: examples

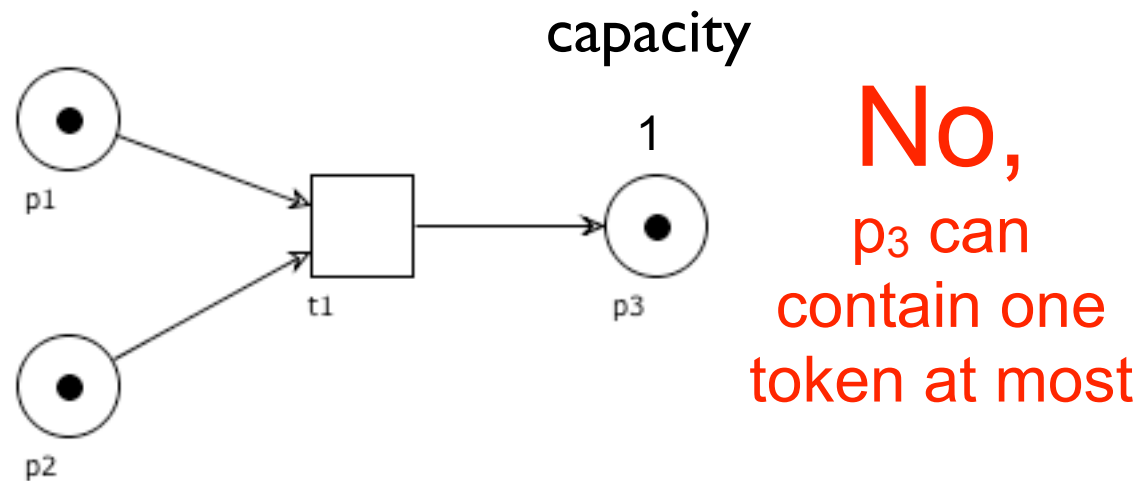


Is  $t_1$  enabled?



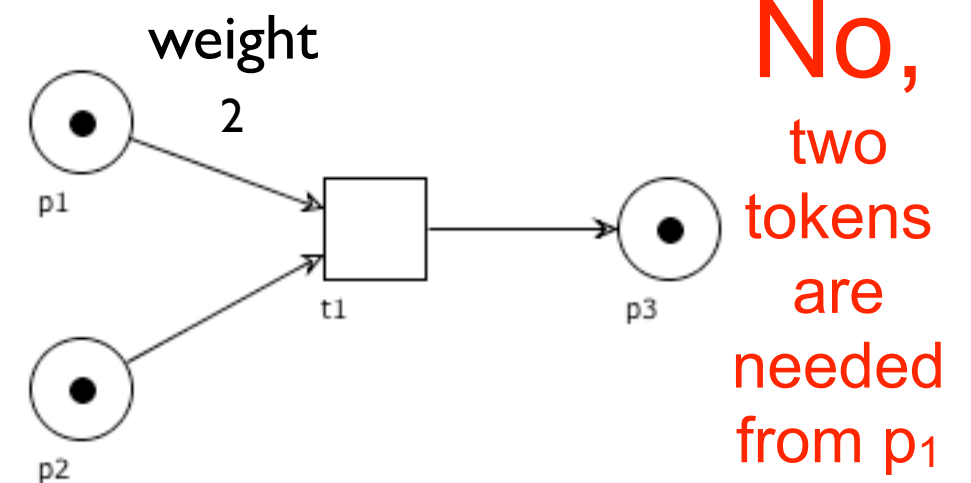


# P/T net: examples

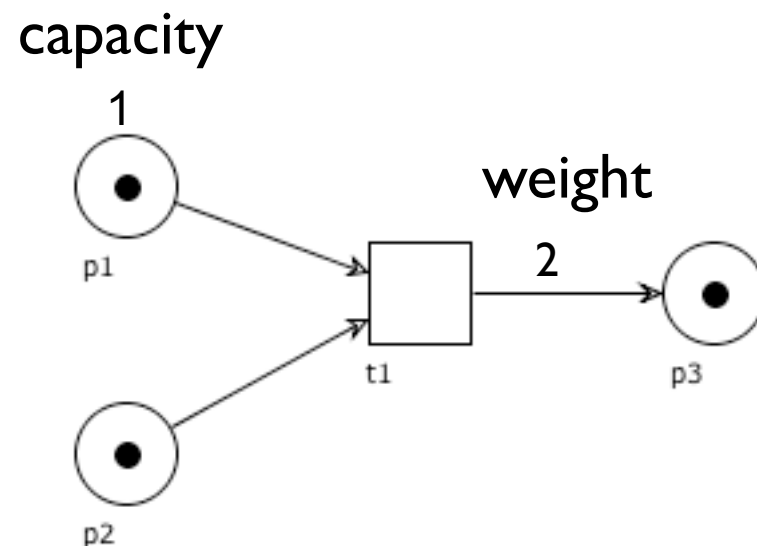


**No,**  
 $p_3$  can  
contain one  
token at most

Is  $t_1$  enabled?



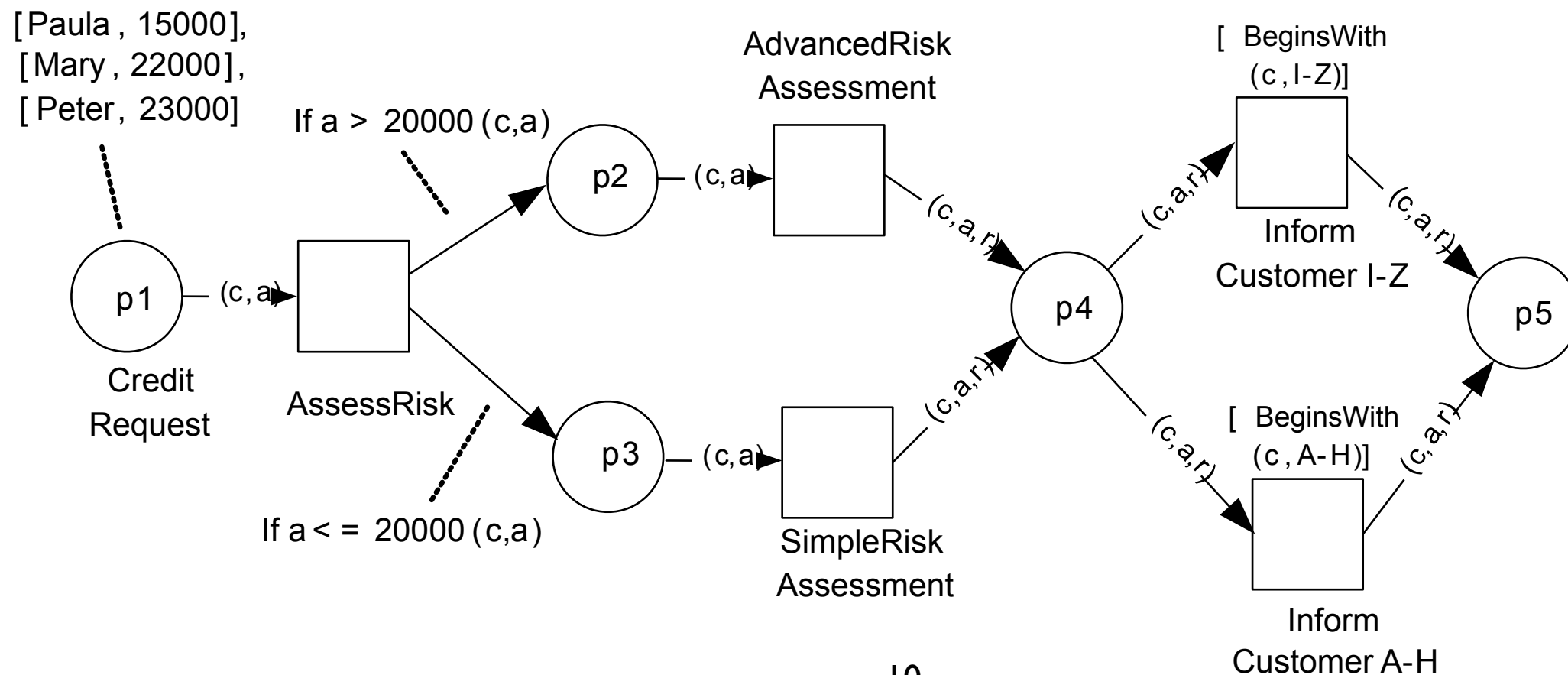
**No,**  
two  
tokens  
are  
needed  
from  $p_1$



**Yes,**  
the firing  
leads to  $3p_3$

# Coloured nets (also called High-Level)

A **coloured net** is a Petri net whose tokens can carry data and whose transitions can check data (see exact definition in Weske's book)



# Workflow nets

# Workflow nets features

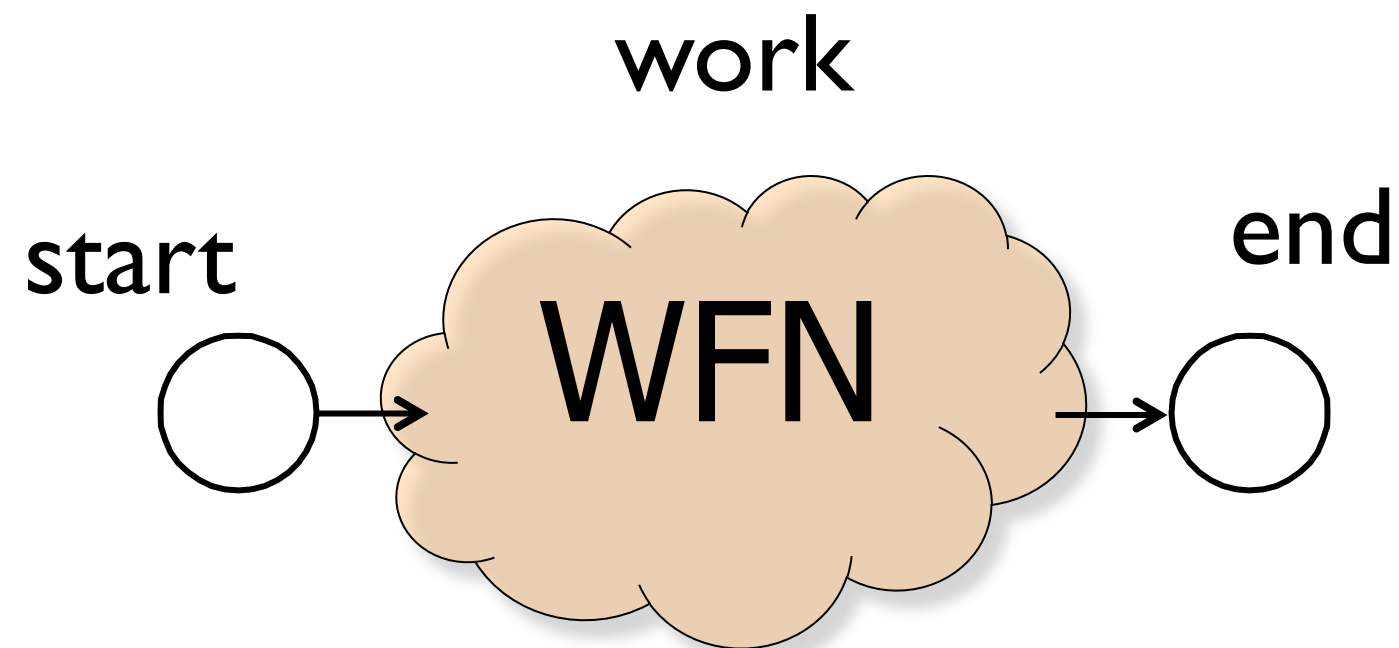
Tailored to the representation of business processes

Formal (unambiguous) semantics

Structural restrictions

Decorated graphical representation

# Workflow net: idea



# Workflow net

## Definition:

A Petri net  $(P, T, F)$  is called **workflow net** if:

1. there is a distinguished *initial place*  $i \in P$  with  $\bullet i = \emptyset$
2. there is a distinguished *final place*  $o \in P$  with  $o \bullet = \emptyset$
3. every other place and transition belongs to a path from  $i$  to  $o$

# Workflow net: Rationale

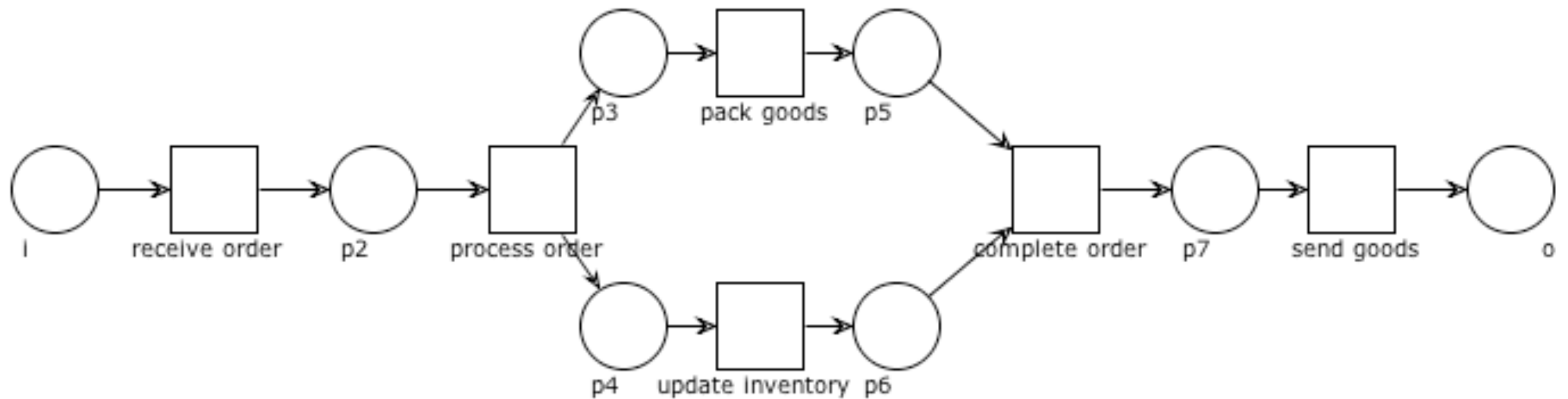
1. a token in  $i$  represents a process instance not yet started
2. a token in  $o$  represents a finished case
3. each place and each transition can participate in a case

## Definition:

A Petri net  $(P, T, F)$  is called **workflow net** if:

1. there is a distinguished *initial place*  $i \in P$  with  $\bullet i = \emptyset$
2. there is a distinguished *final place*  $o \in P$  with  $o \bullet = \emptyset$
3. every other place and transition belongs to a path from  $i$  to  $o$

# WF net: Example





# Basic properties

**Lemma:** In a workflow net there is a **unique** node with no incoming arc

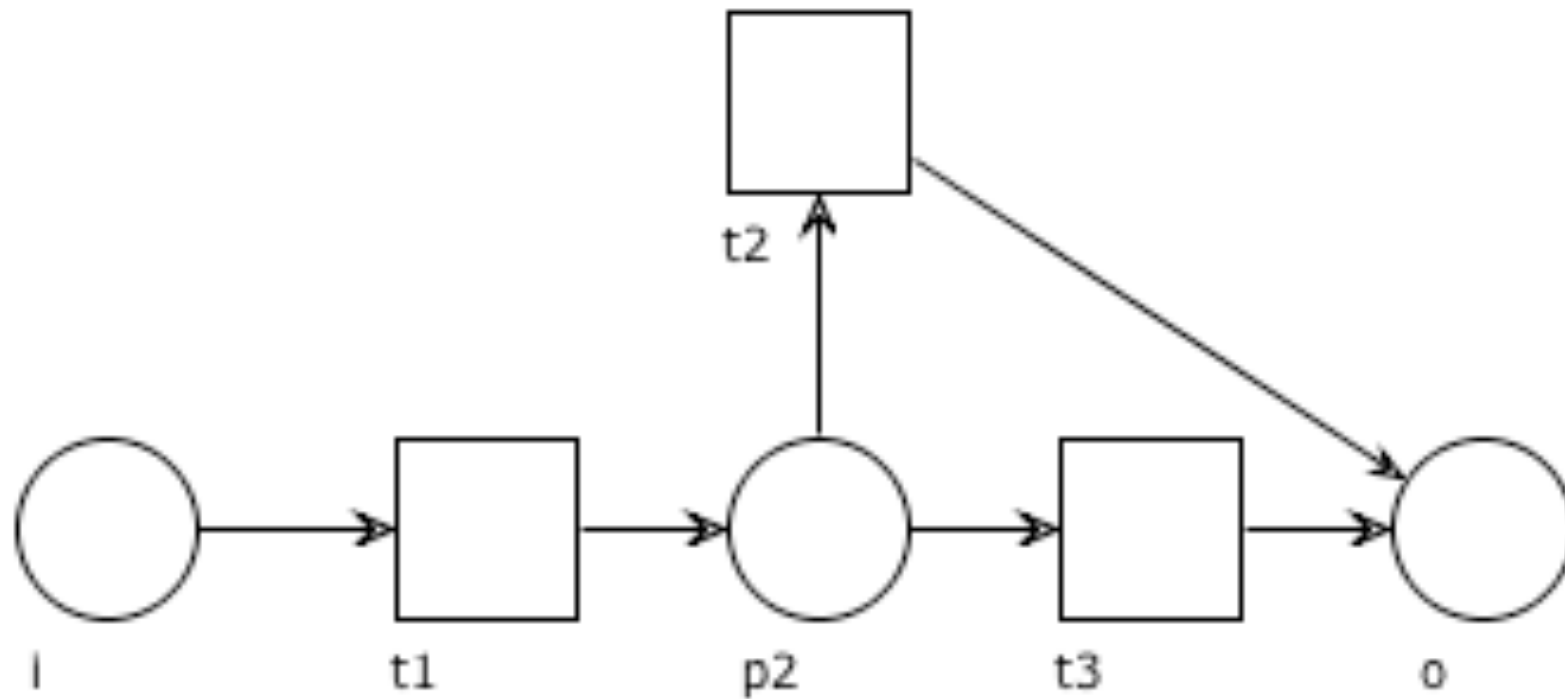
**Lemma:** In a workflow net there is a **unique** node with no outgoing arc

**Exercise:** Guess which nodes are those

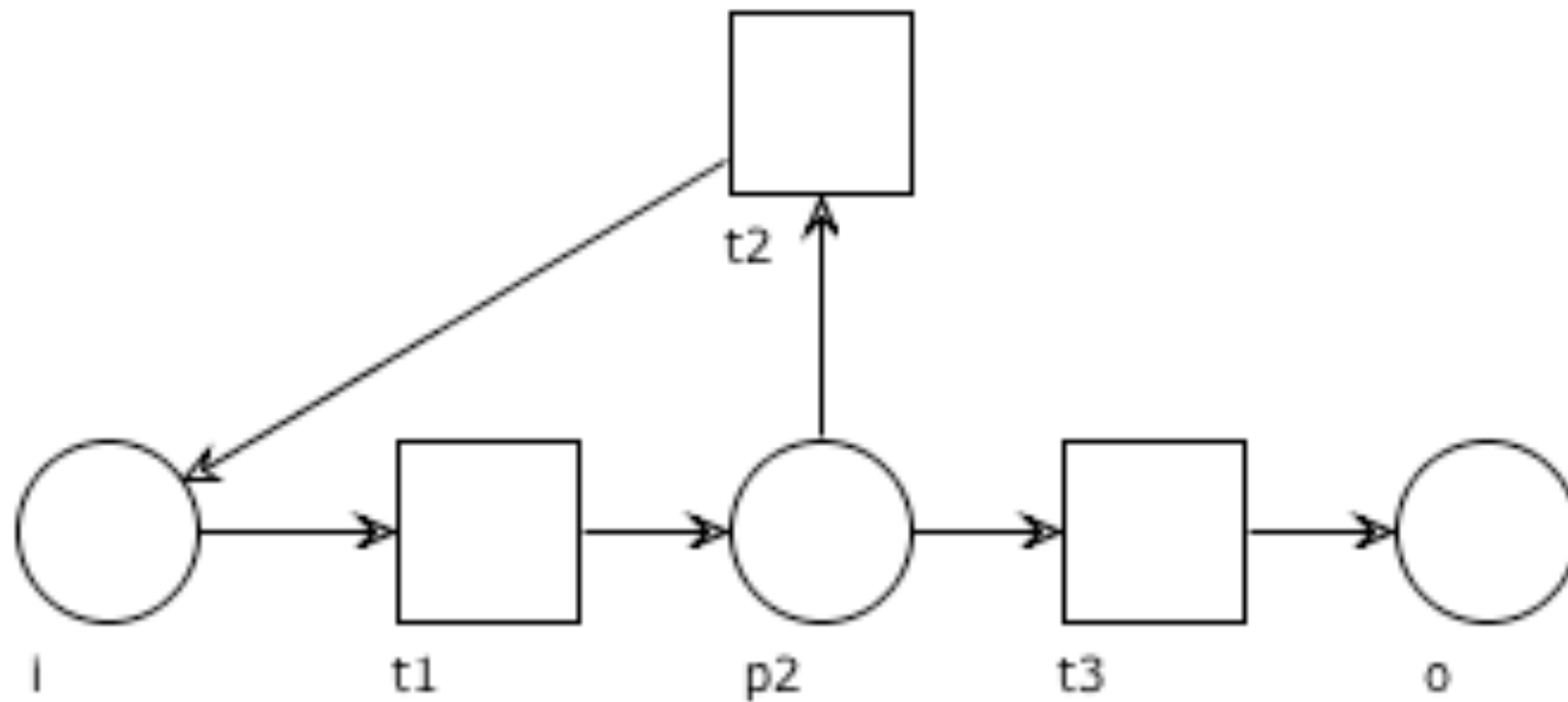
**Exercise:** Prove the above lemmas

(hint: suppose the nodes are not unique, reach a contradiction)

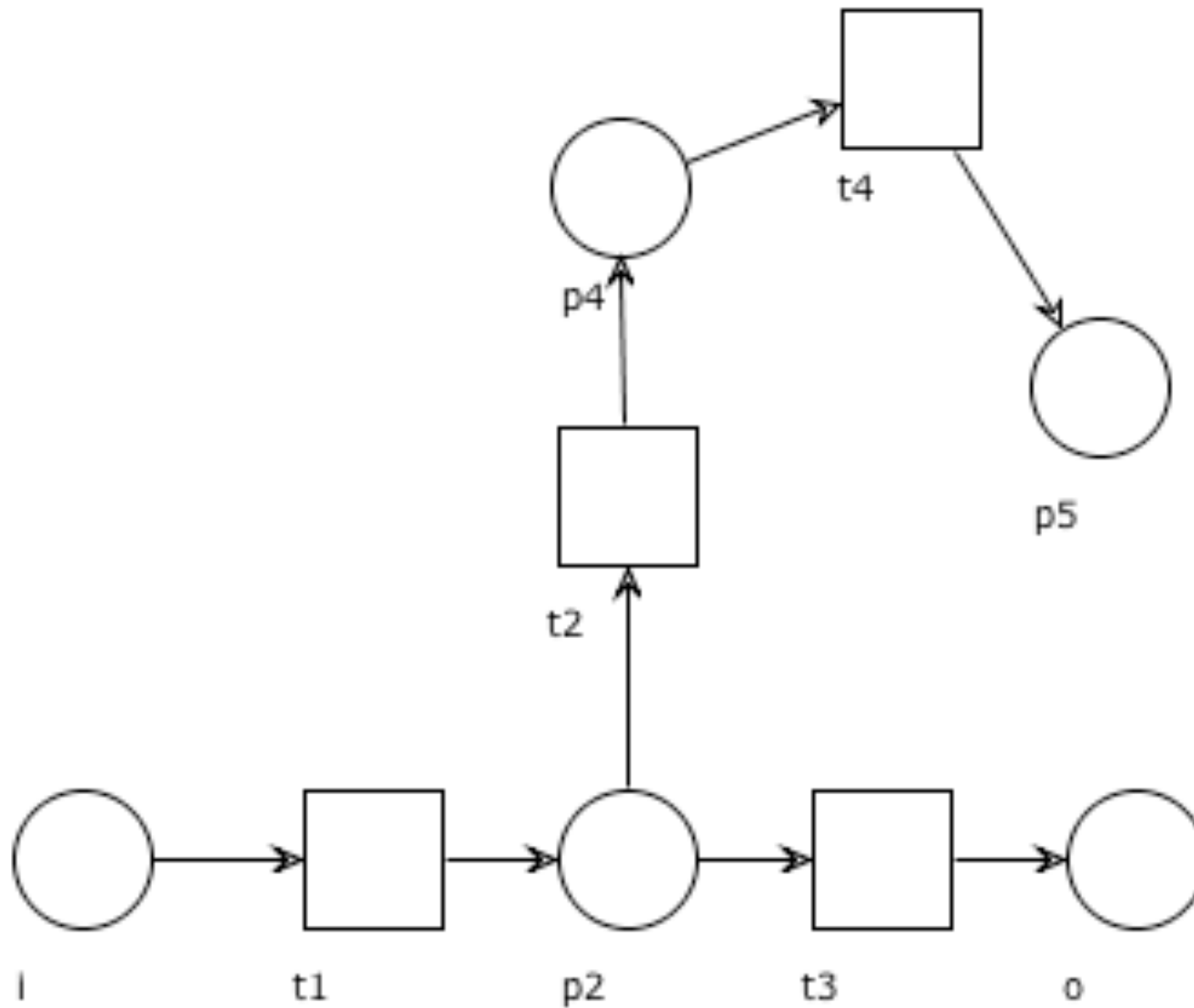
# Question time: WF net?



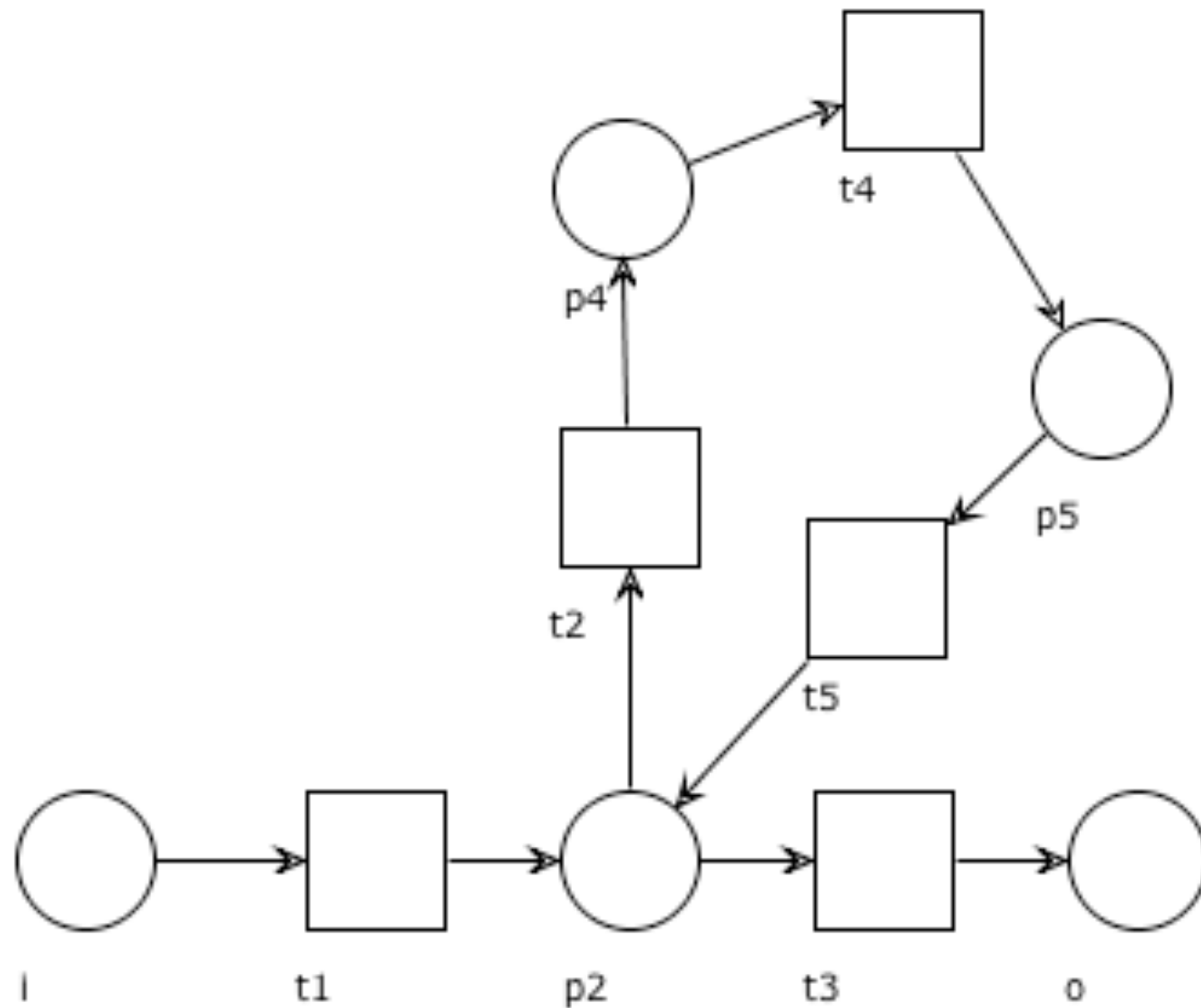
# Question time: WF net?



# Question time: WF net?

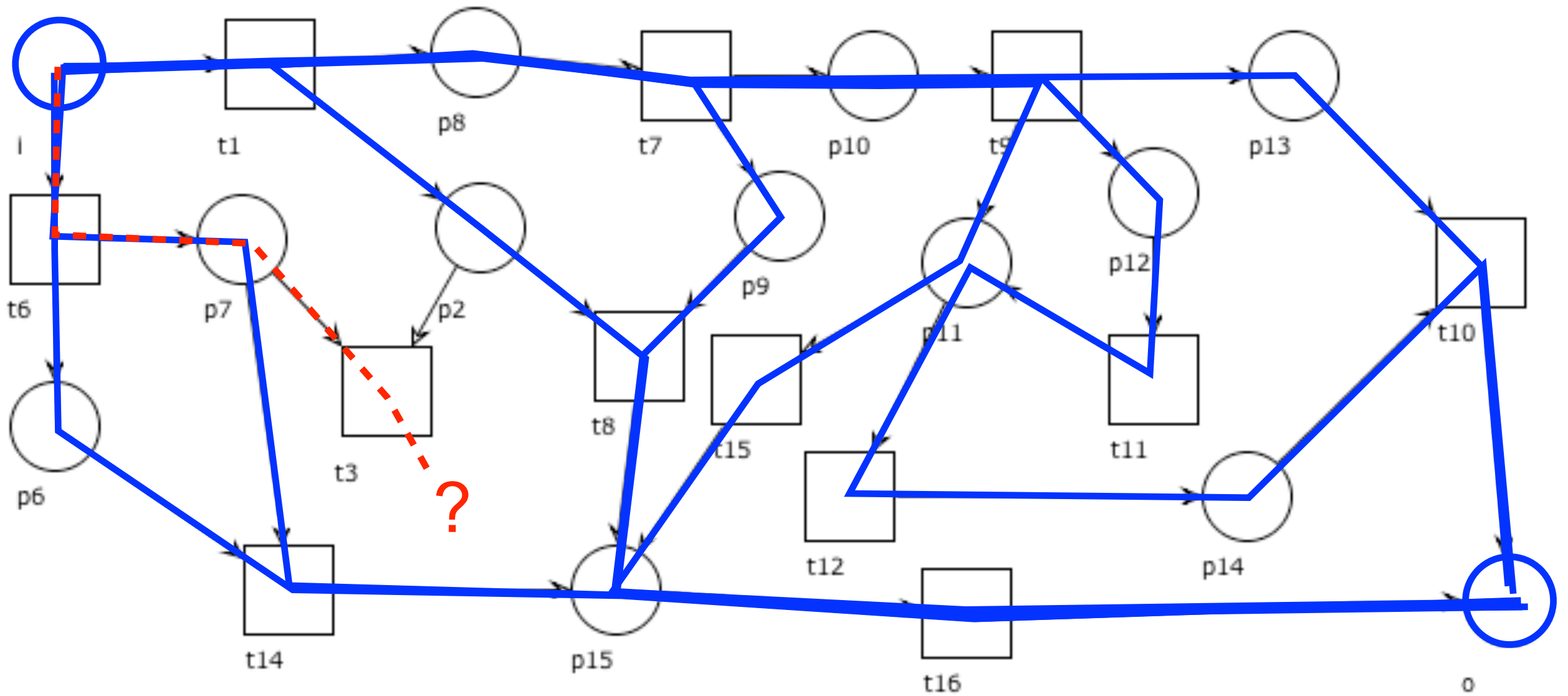


# Question time: WF net?

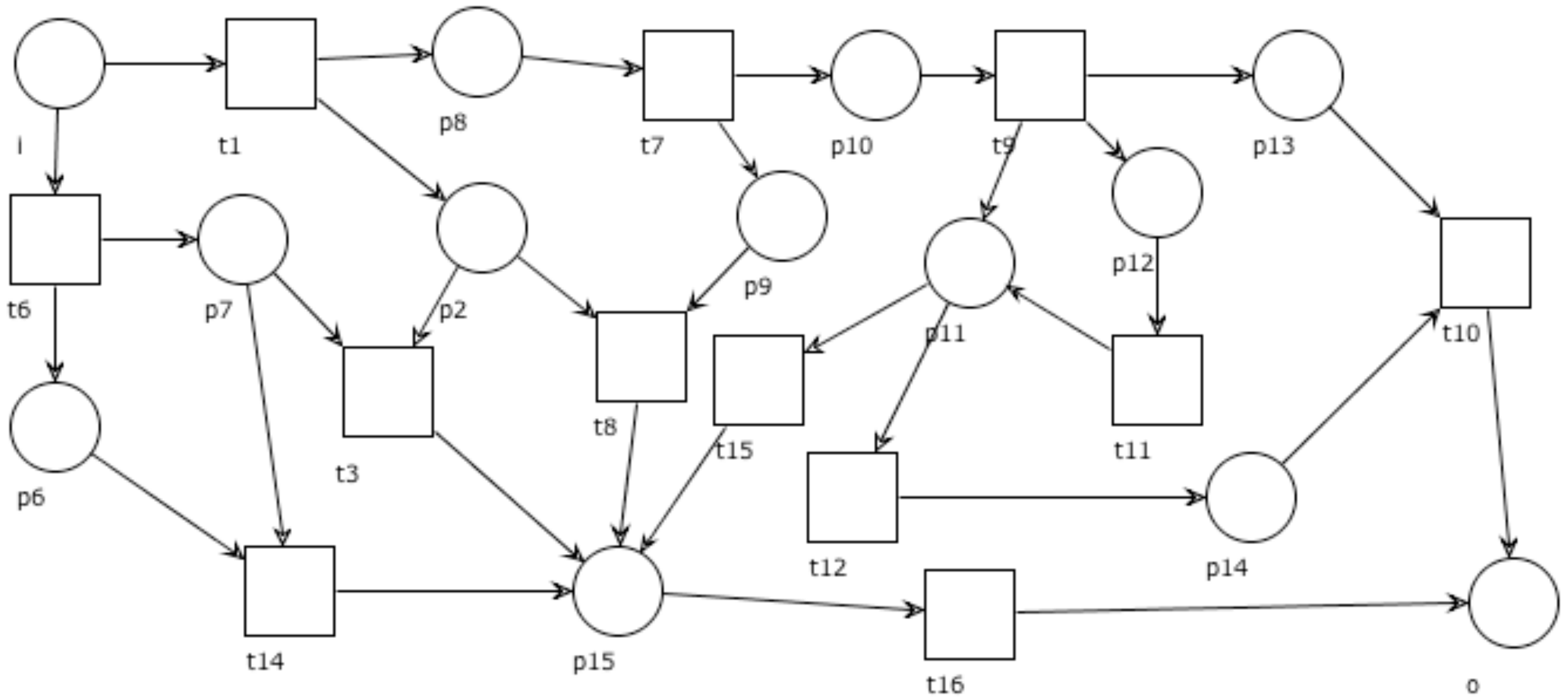




# Question time: WF net?

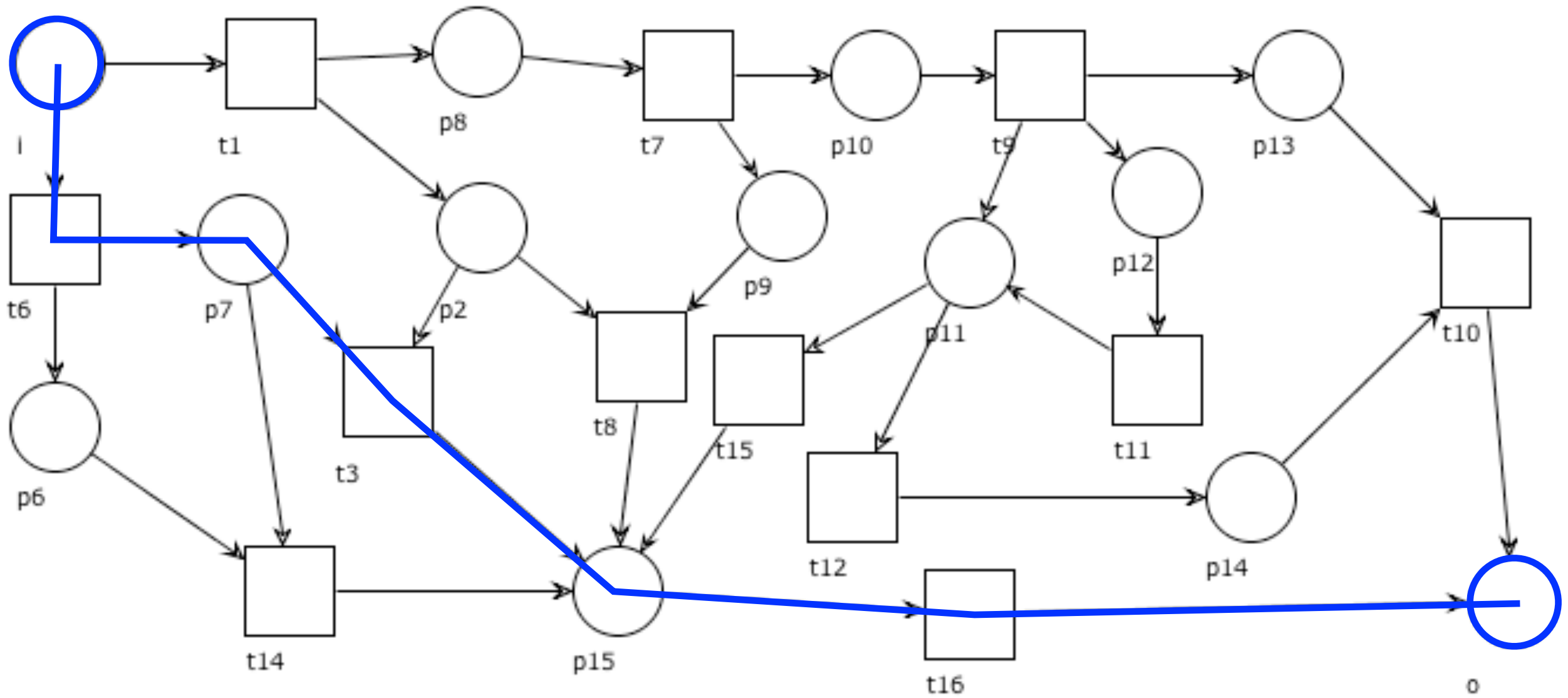


# Question time: WF net?



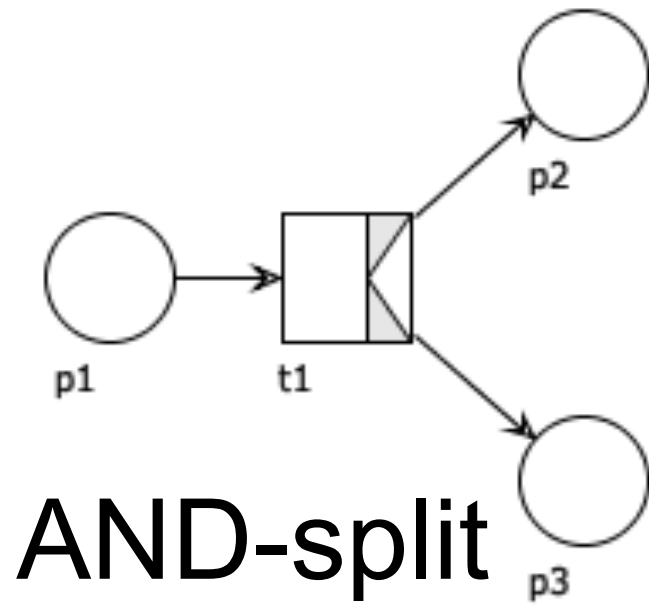


# Question time: WF net?

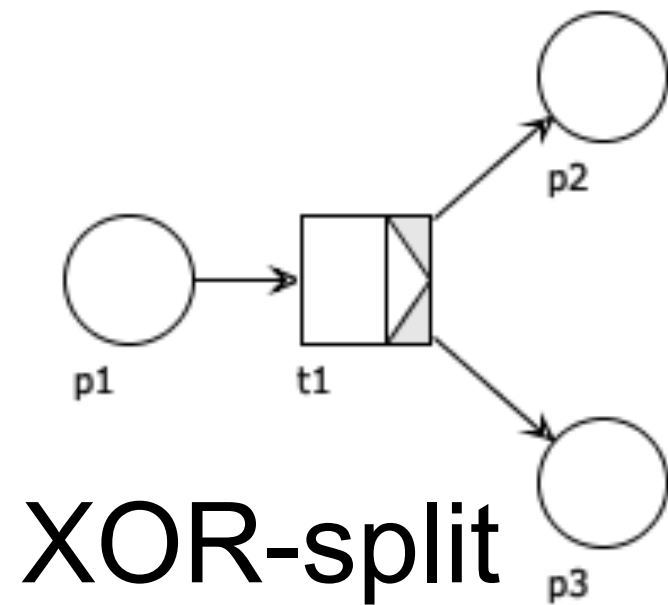
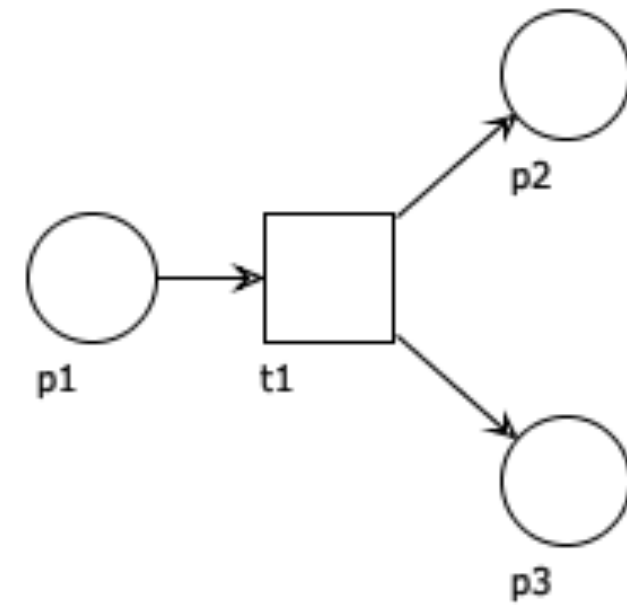


Yes

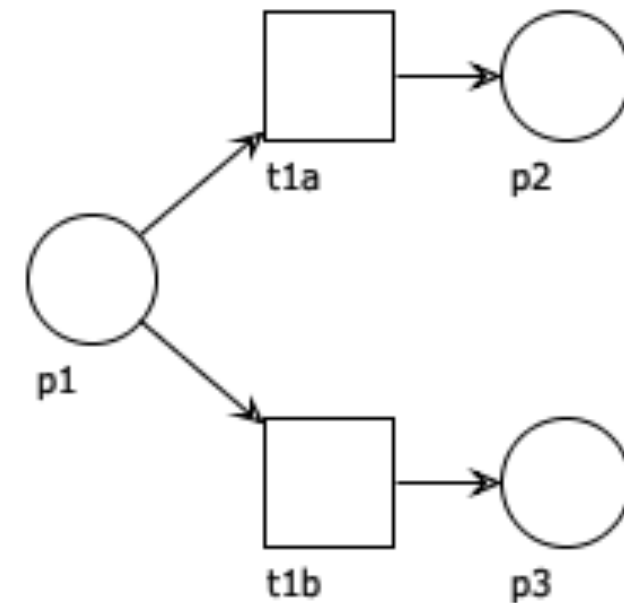
# Syntax sugar: split



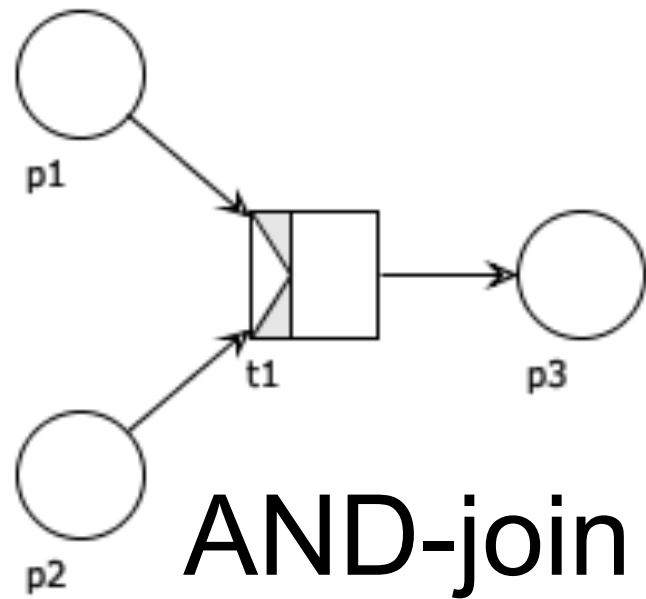
stands for



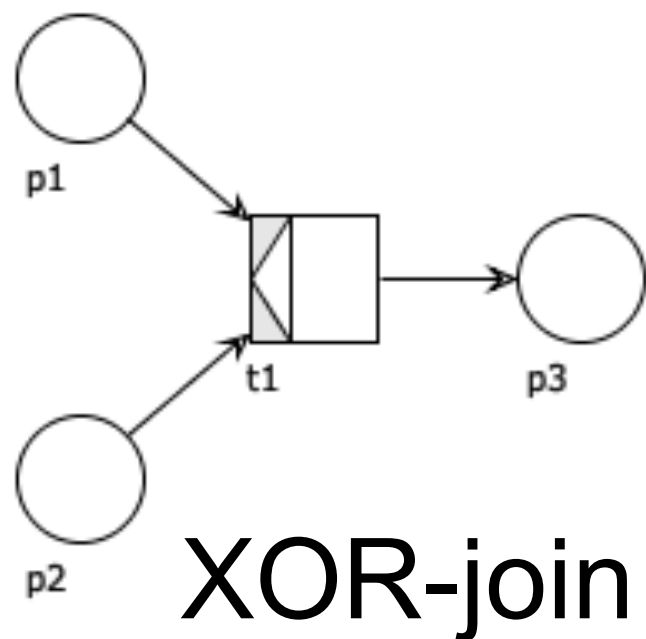
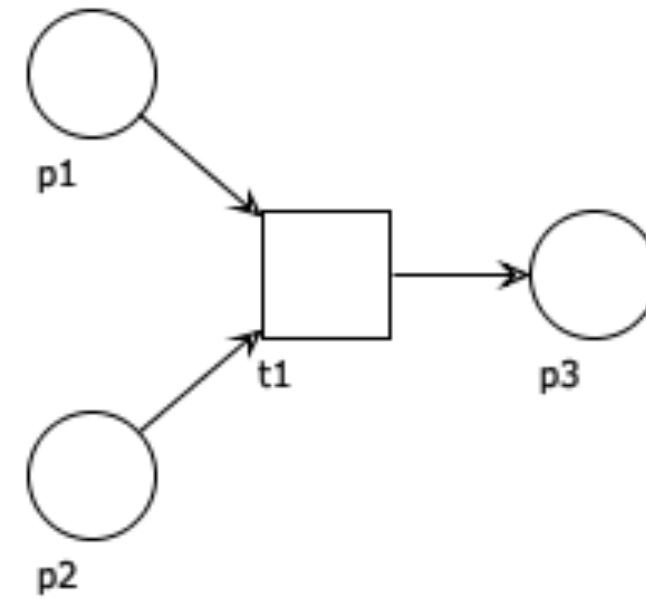
stands for



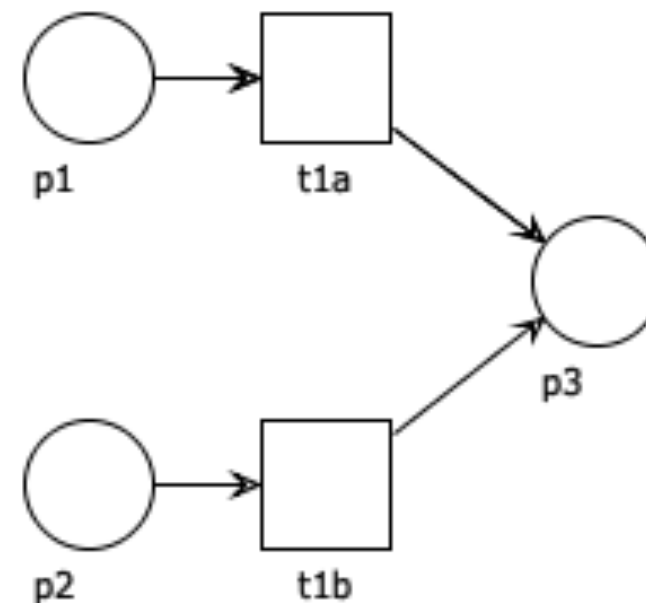
# Syntax sugar: join



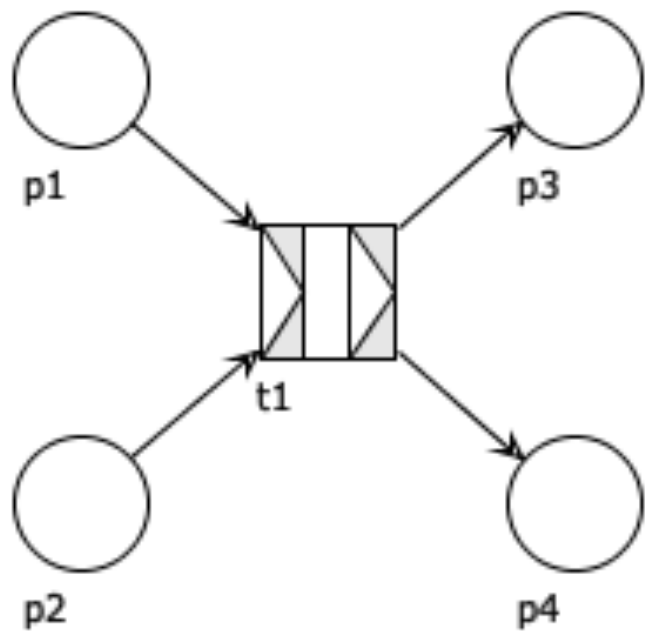
stands for



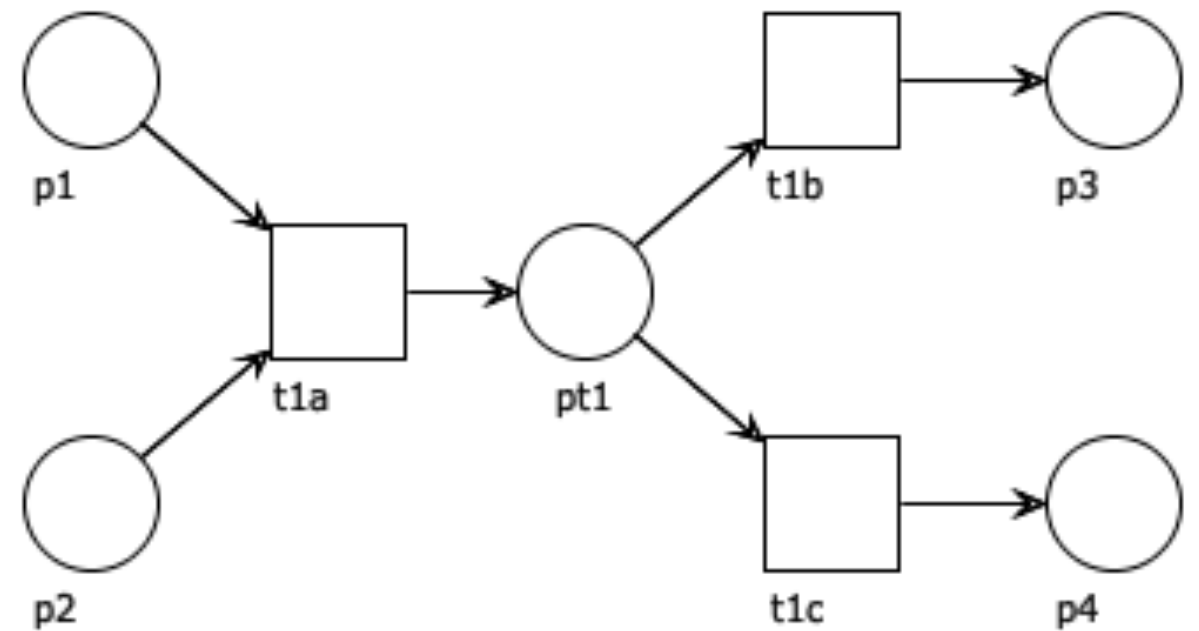
stands for



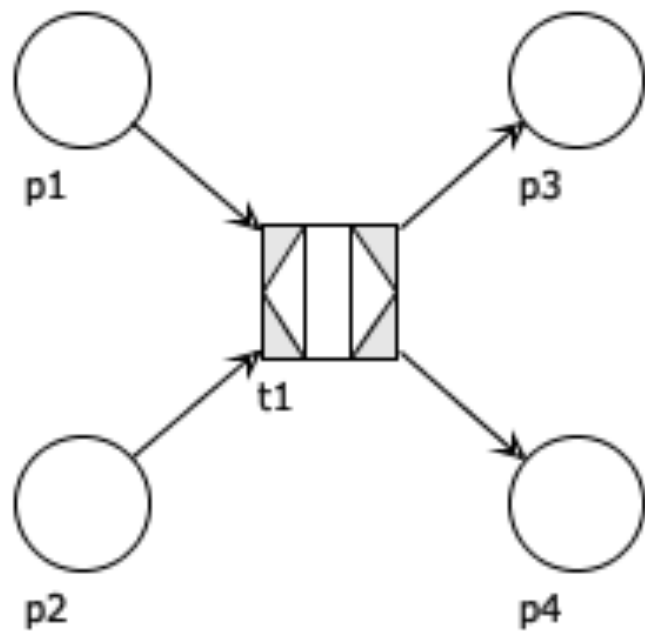
# Syntax sugar: any combination is also possible



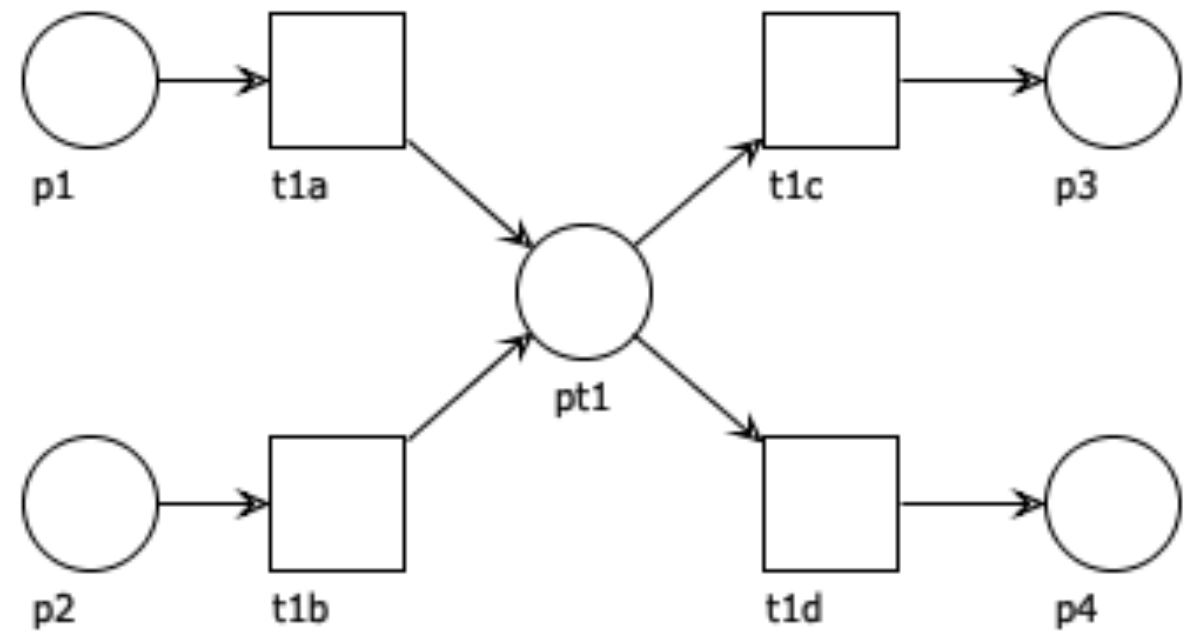
stands for



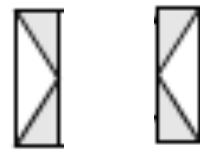
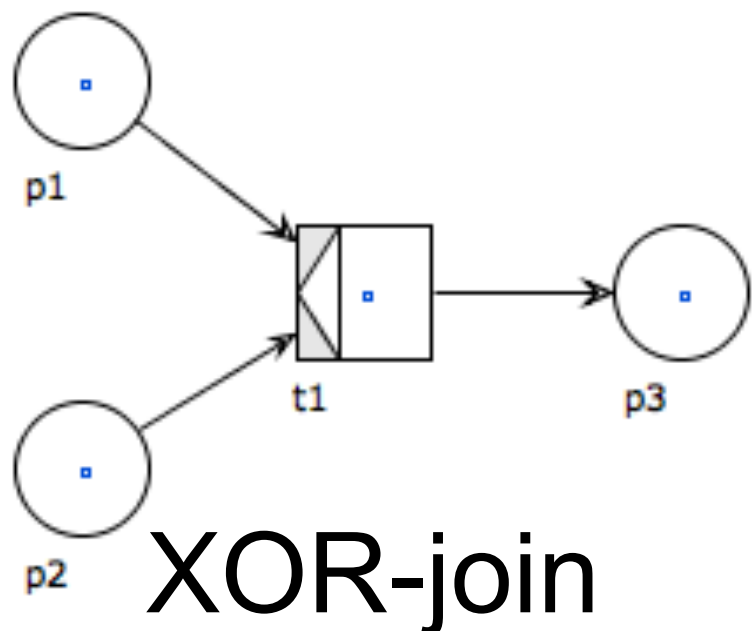
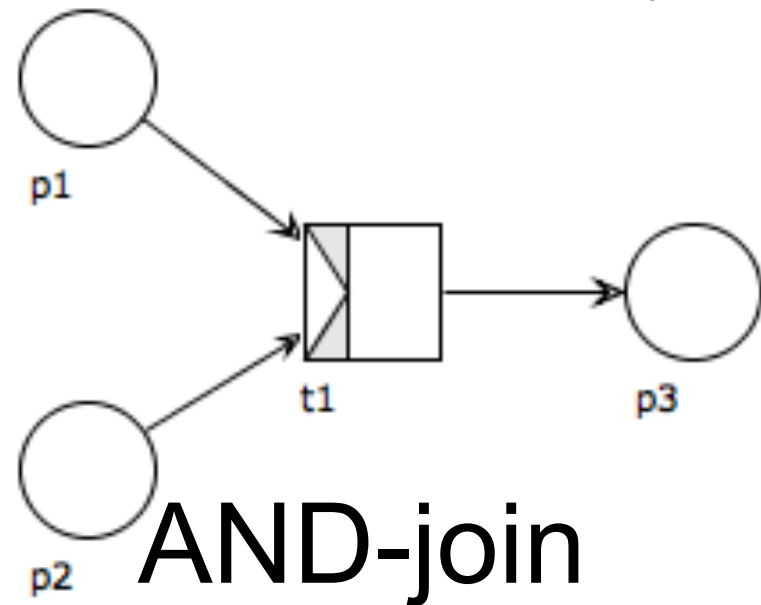
# Syntax sugar: any combination is also possible



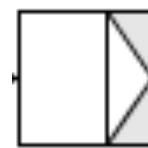
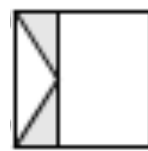
stands for



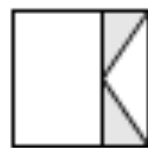
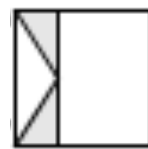
# Syntax sugar: a personal note



Chosen decorations  
are too similar!



Different meanings  
if differently placed!

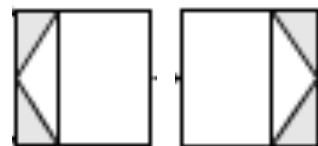
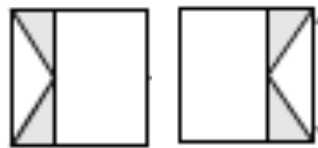


Unnecessary for AND  
(redundant)!

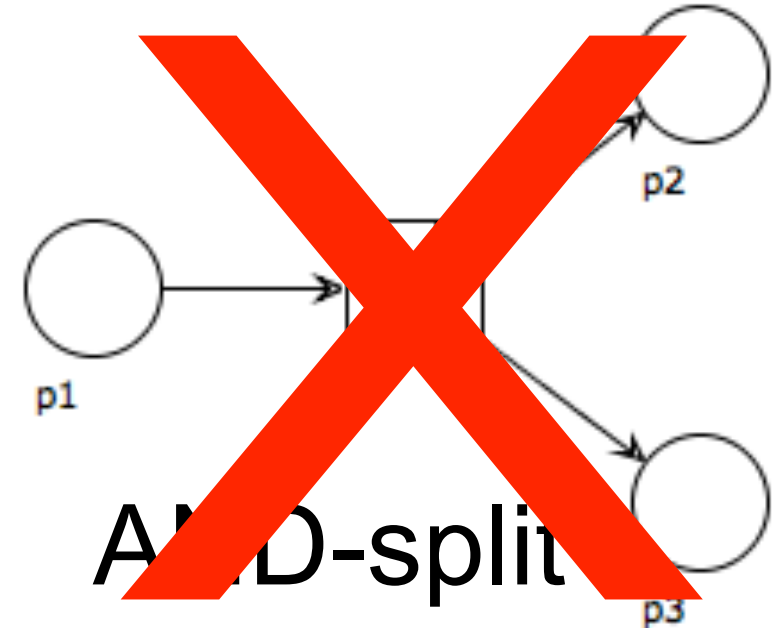
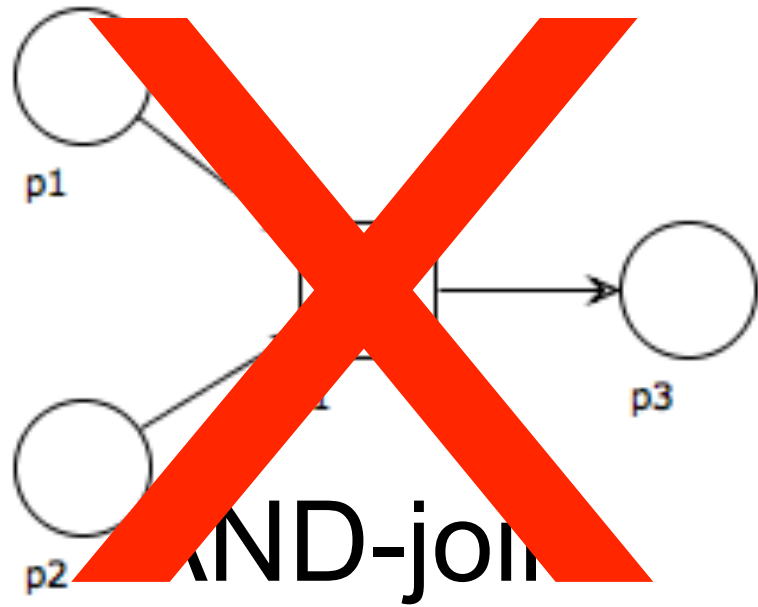
# Syntax sugar: a personal note

Why there?

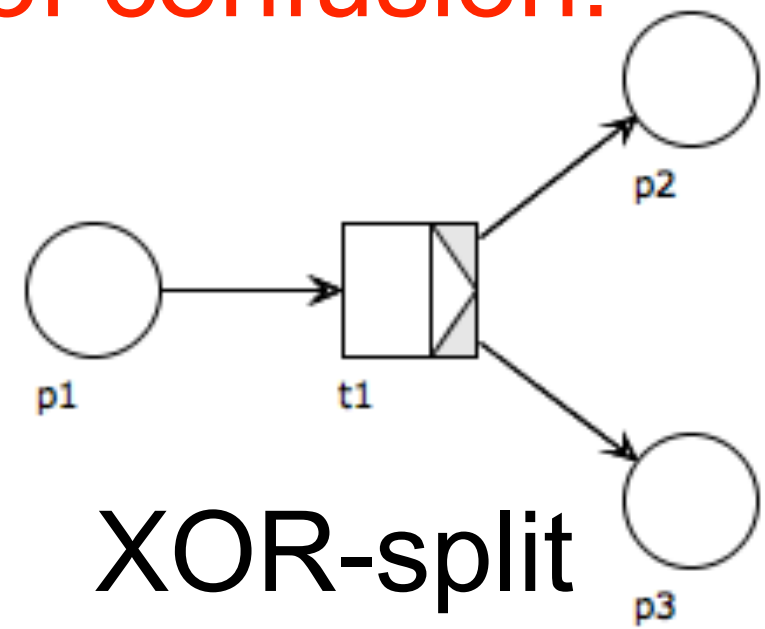
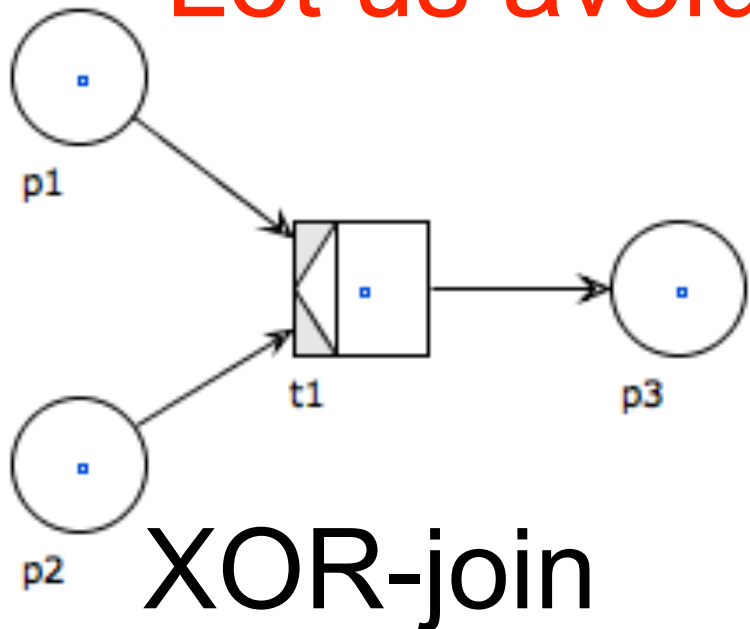
Because of gateways



# Syntax sugar: a personal note



Let us avoid any source of confusion!

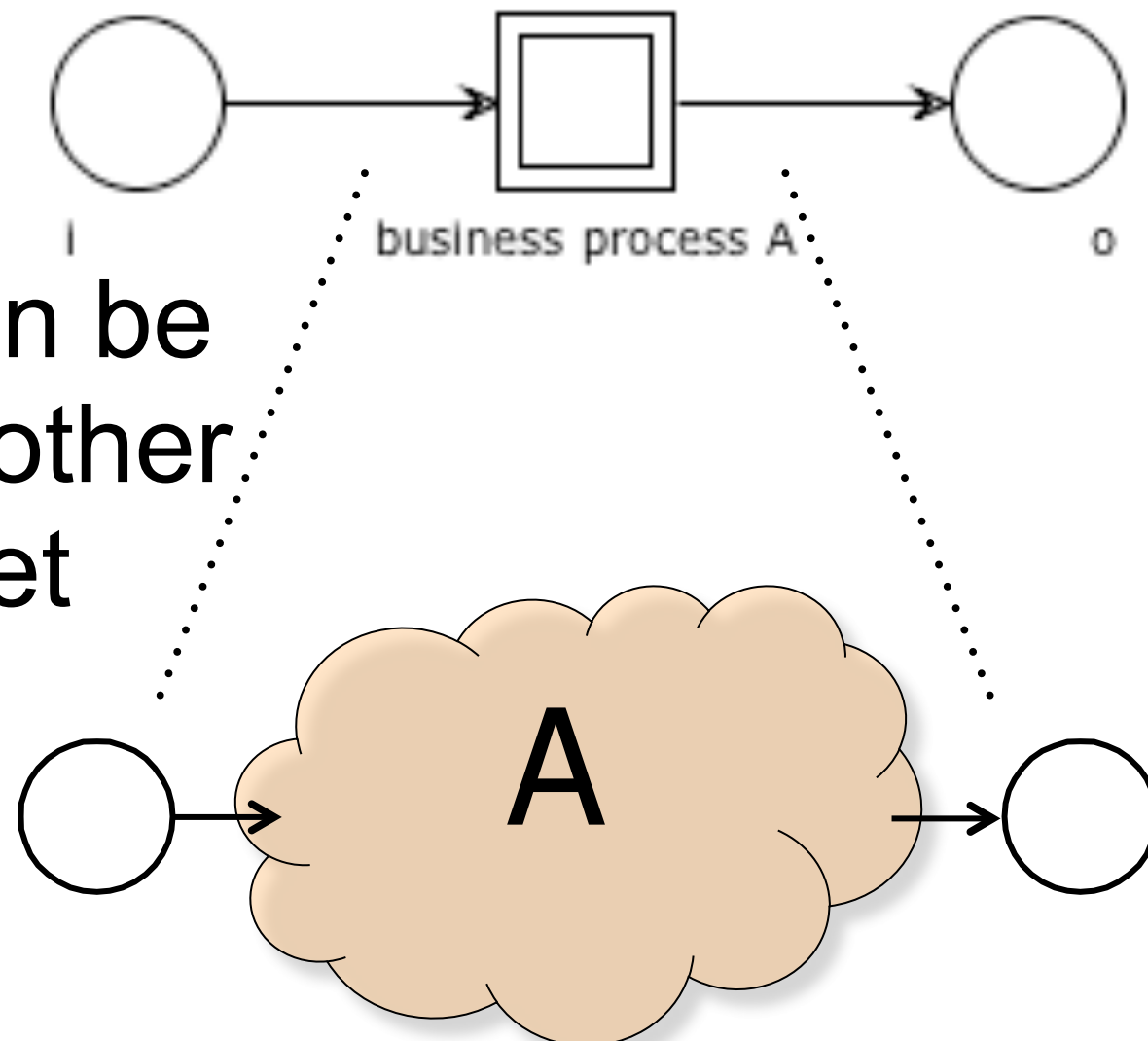




# Hierarchical structuring

Uniqueness of entry / exit point facilitate the hierarchical structuring of WF nets

a transition can be realized by another workflow net

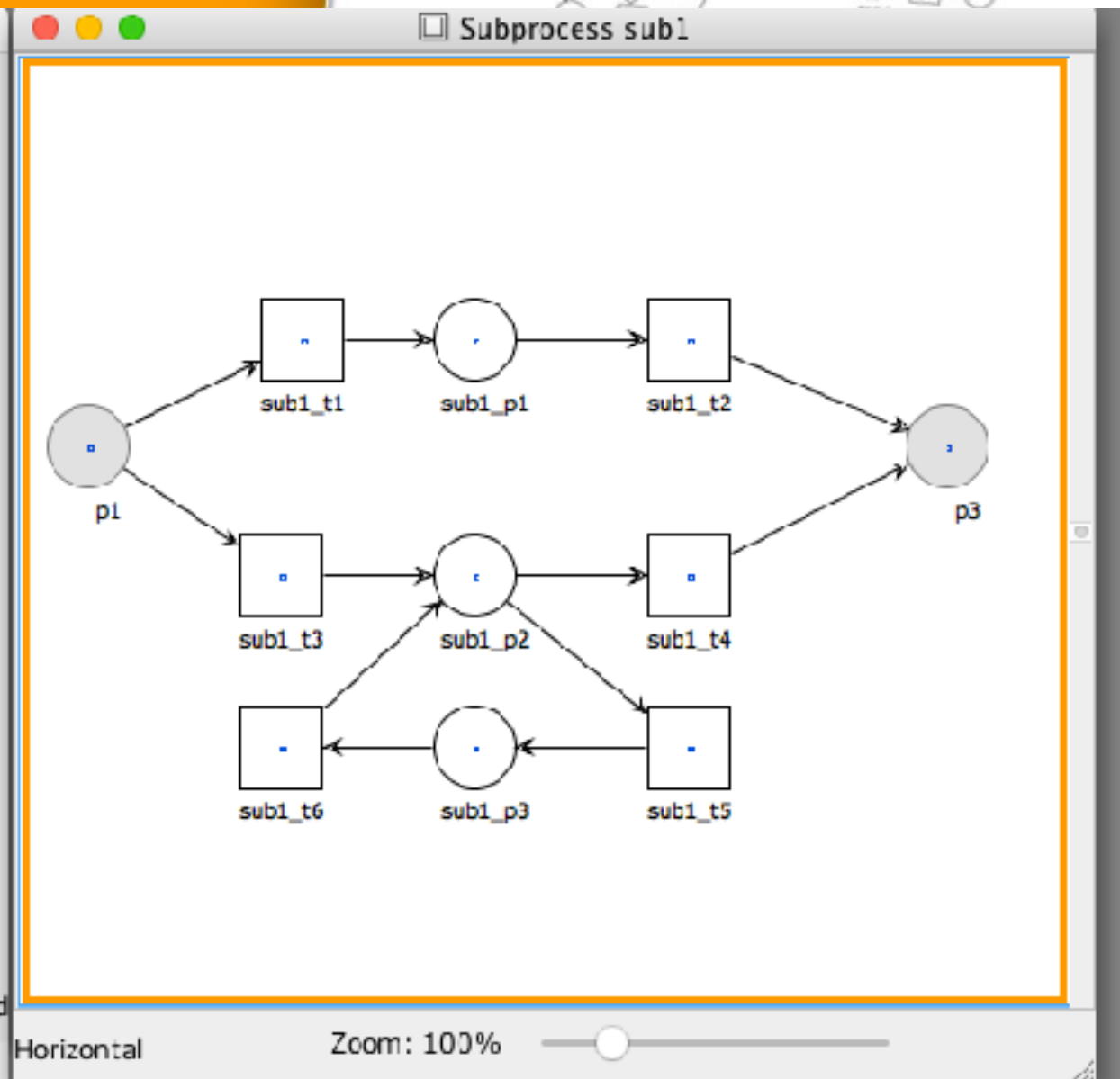
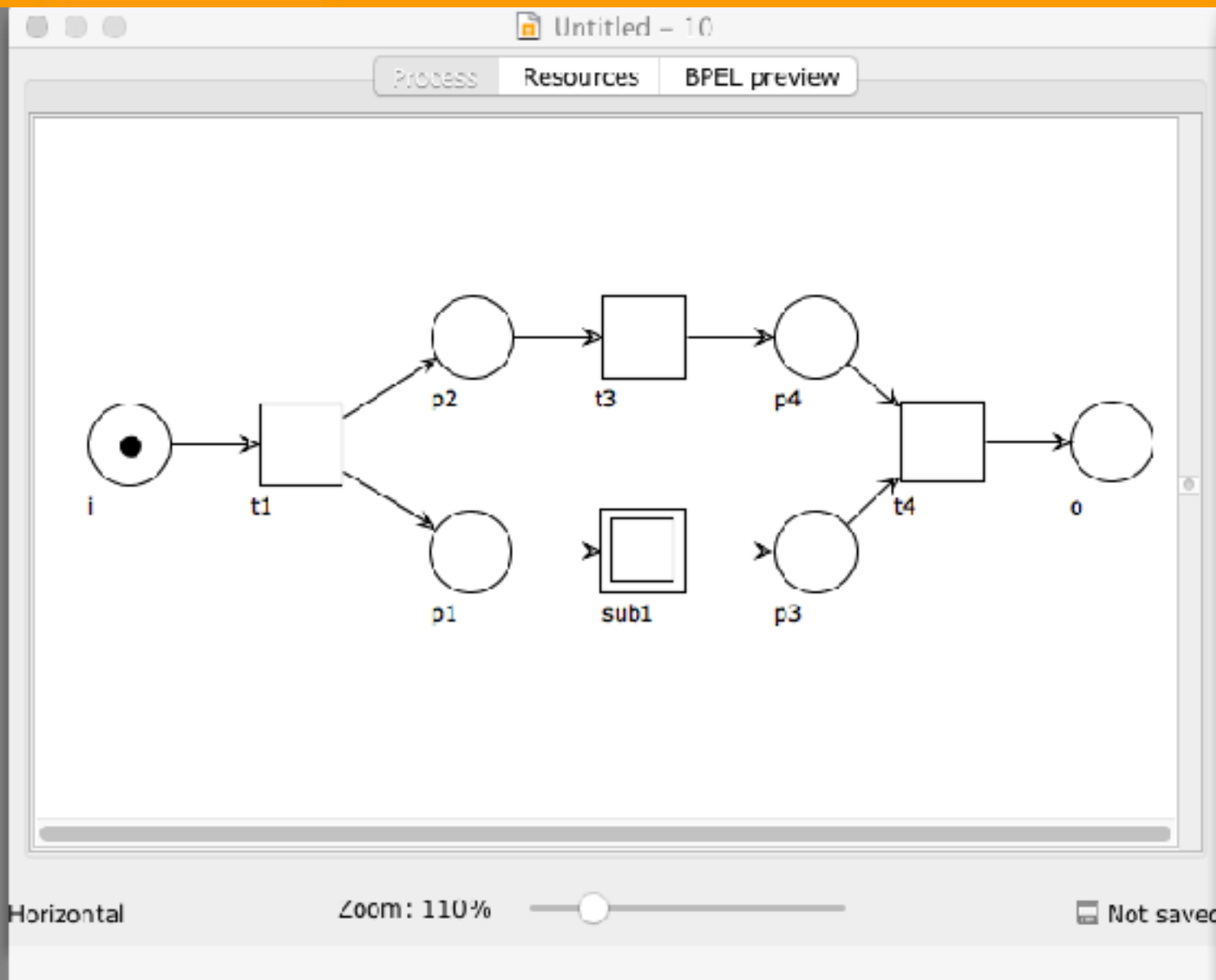
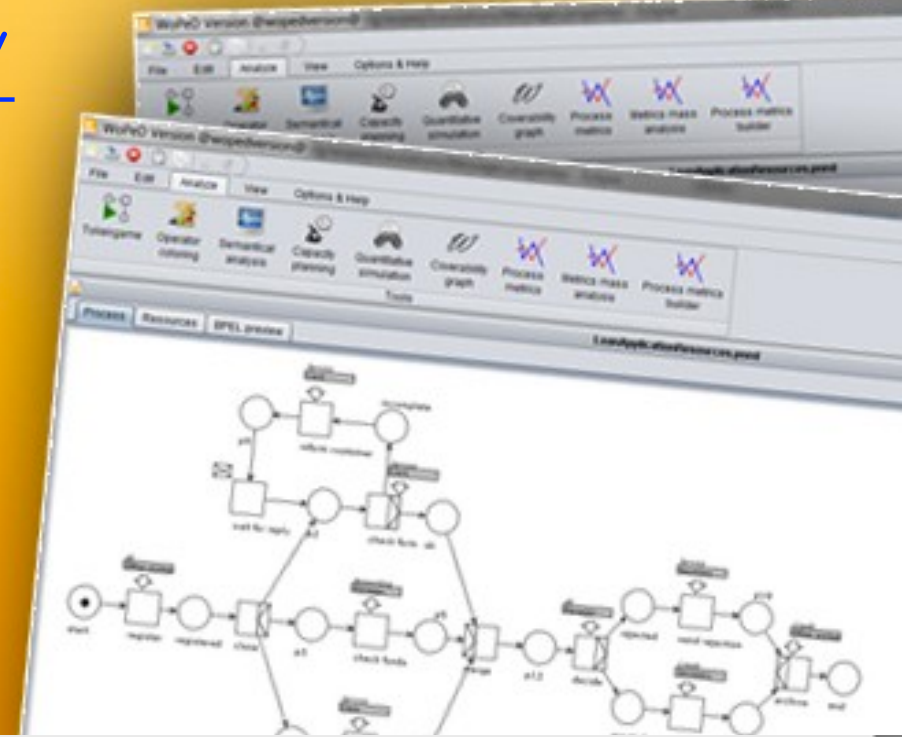


<http://woped.dhbw-karlsruhe.de/woped/>

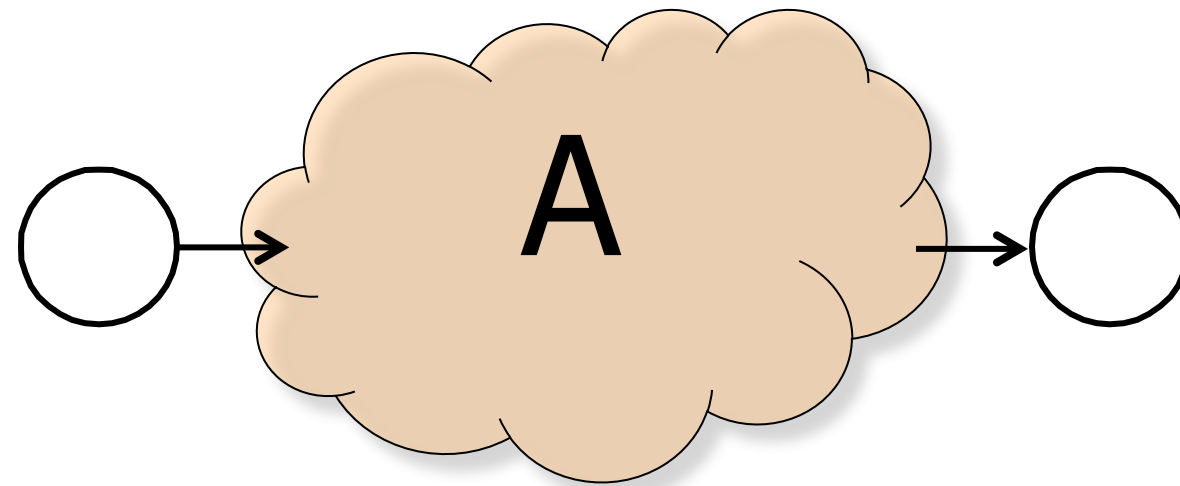
# WoPeD (3.7.1)

## Workflow Petri Net Designer

*Download WoPeD at sourceforge!*



# Language of a workflow net



The language of a workflow net is the set of firing sequences that go from  $i$  to  $o$

$$L(N) = \{ \sigma \mid i \xrightarrow{\sigma} o \}$$

$L(N)$  defines the admissible traces of the workflow

# Typical control flow aspects

Sequencing

Parallelism (AND-split + AND-join)

Selection (XOR-split + XOR-join)

Iteration (XOR-join + XOR-split)

Capacity constraints:

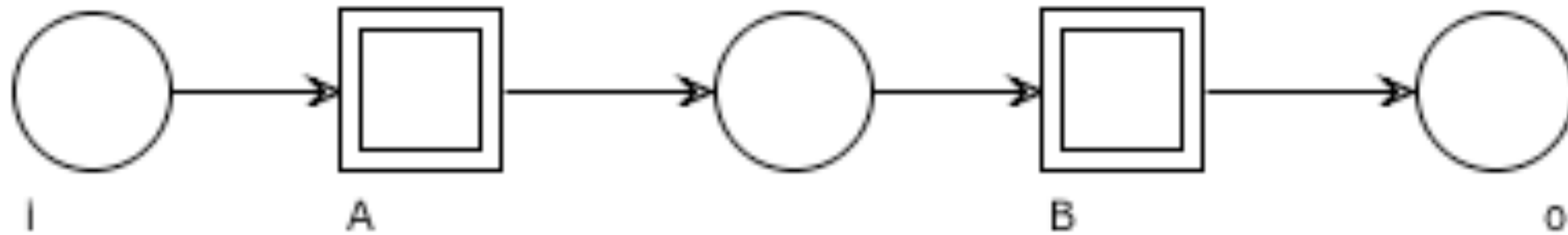
Feedback loop

Mutual exclusion

Alternating

# Sequencing

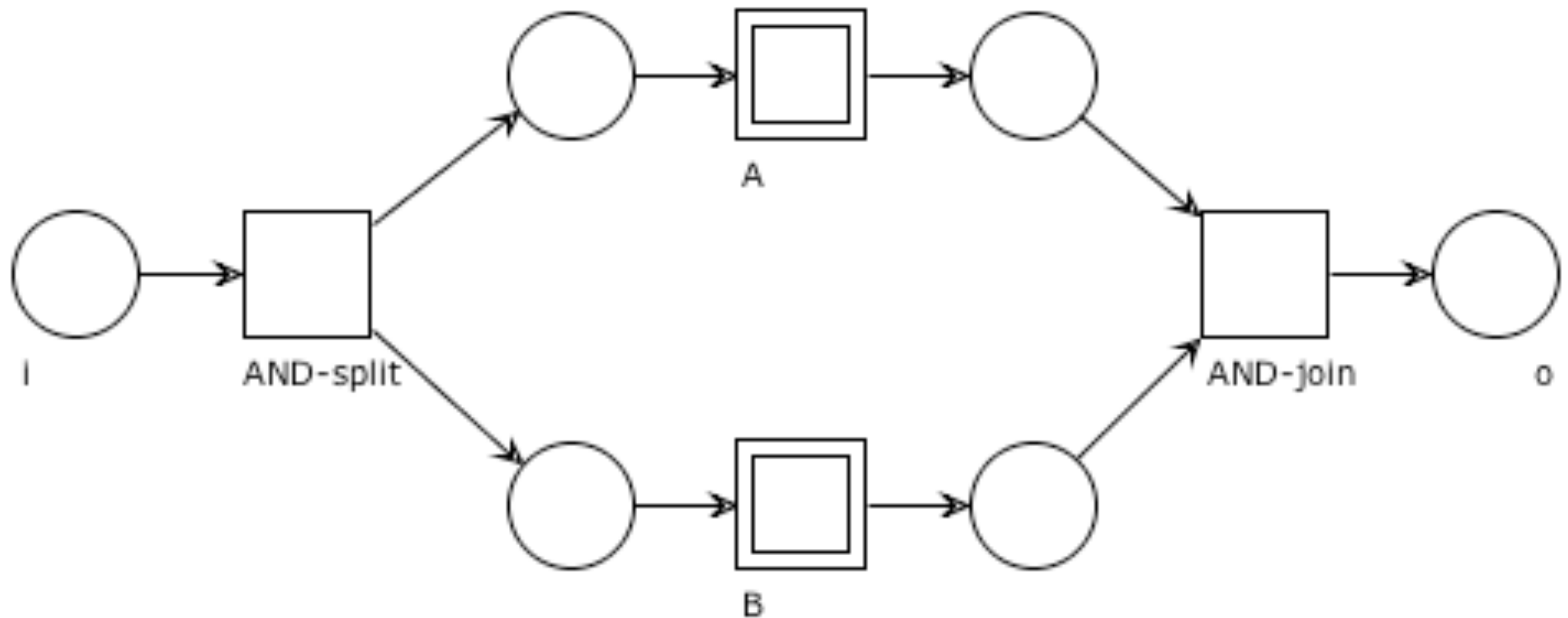
B is executed after A



# Parallelism

(AND-split + AND-join)

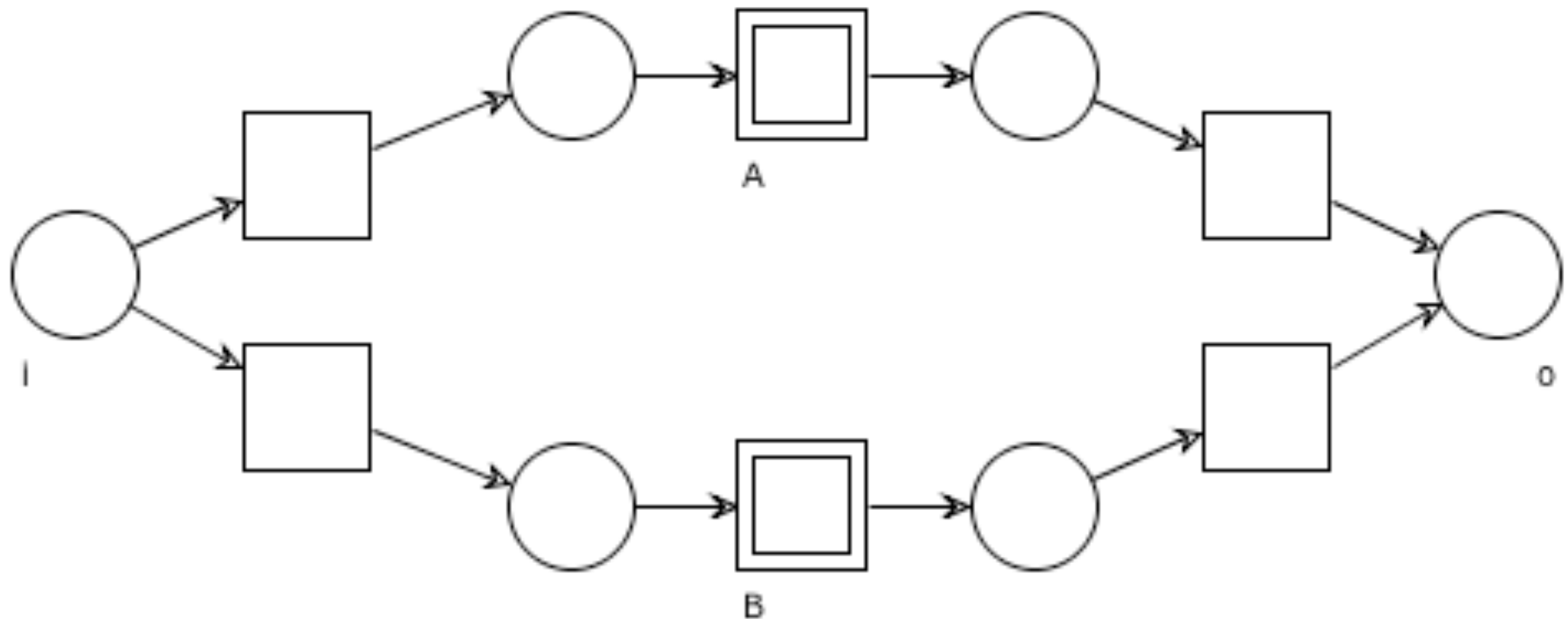
A and B are both executed in no particular order



# Explicit choice

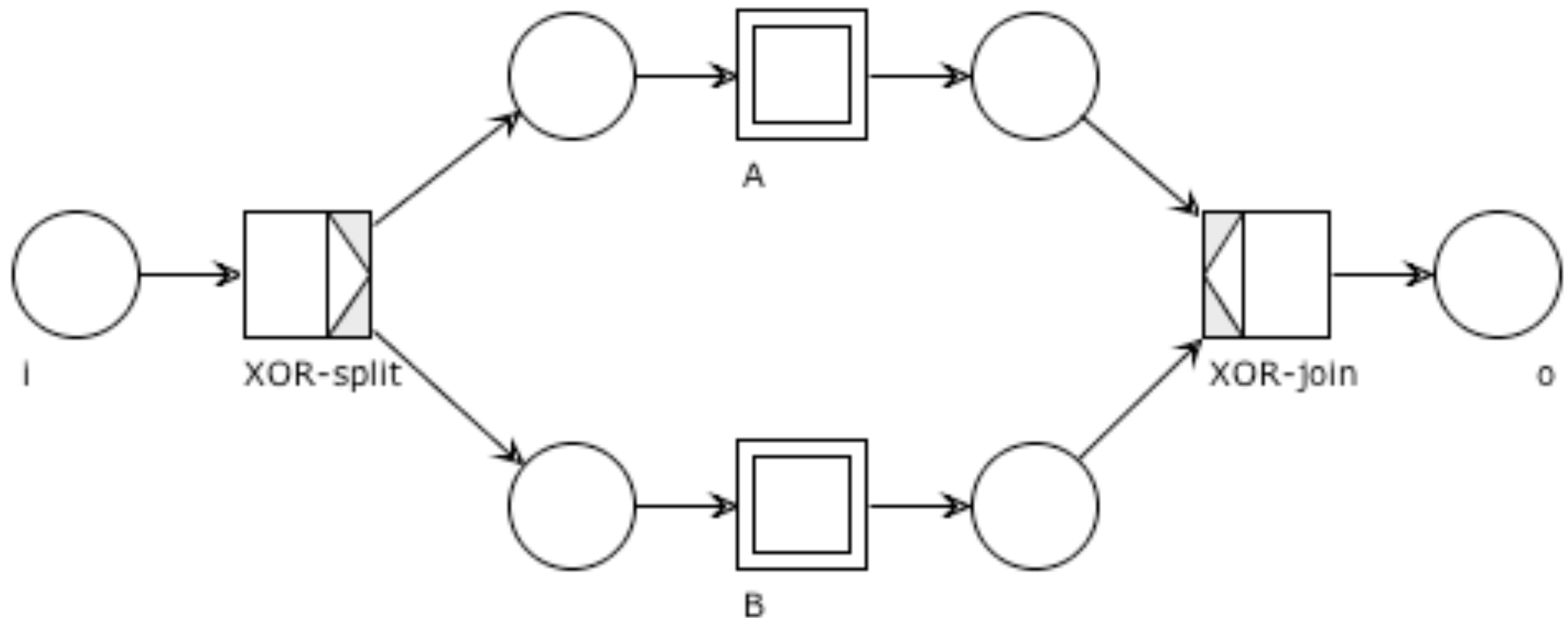
(XOR-split + XOR-join)

Either A or B is executed (choice is **explicit**)



# Explicit choice ("sugared" version)

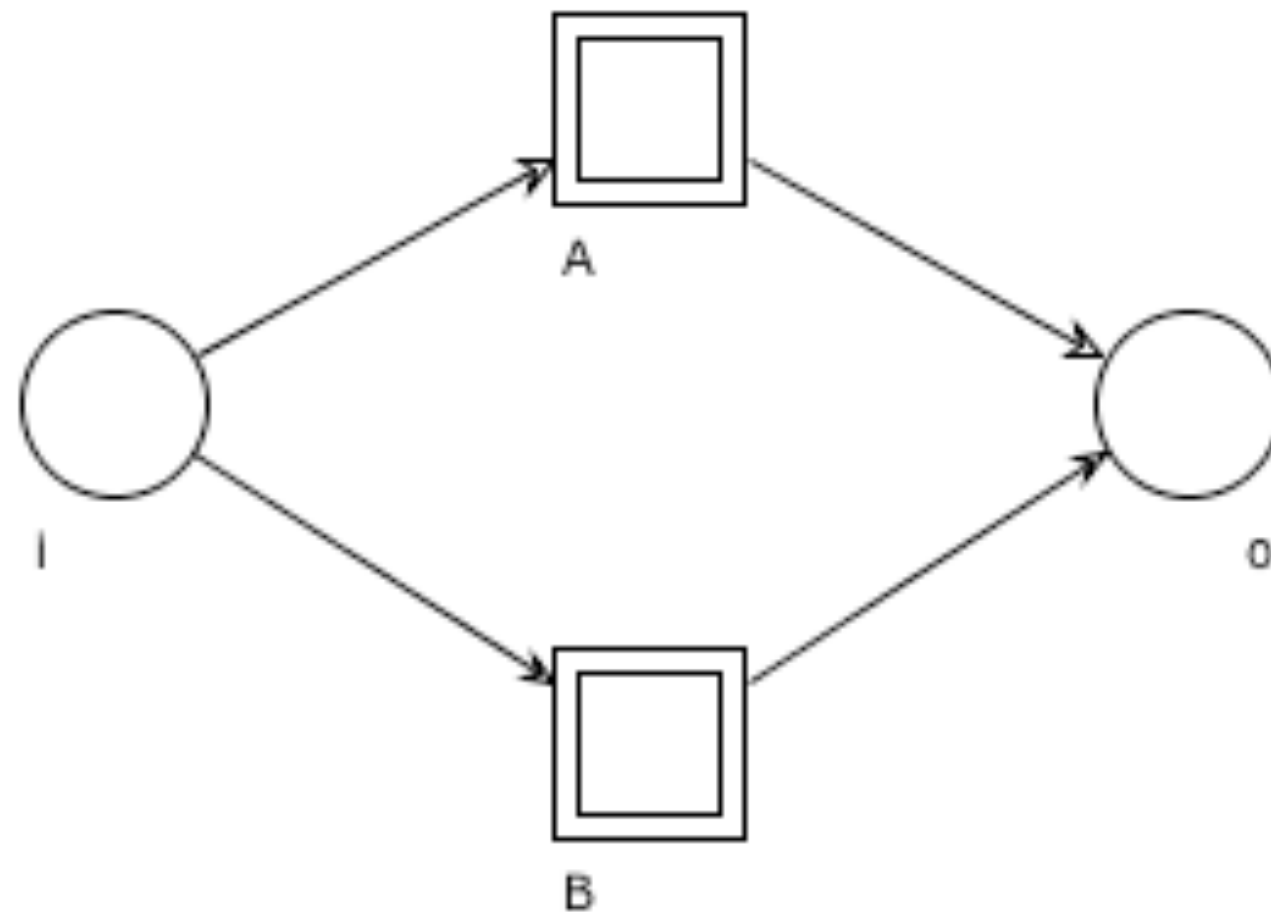
Decorated version for business process stakeholders





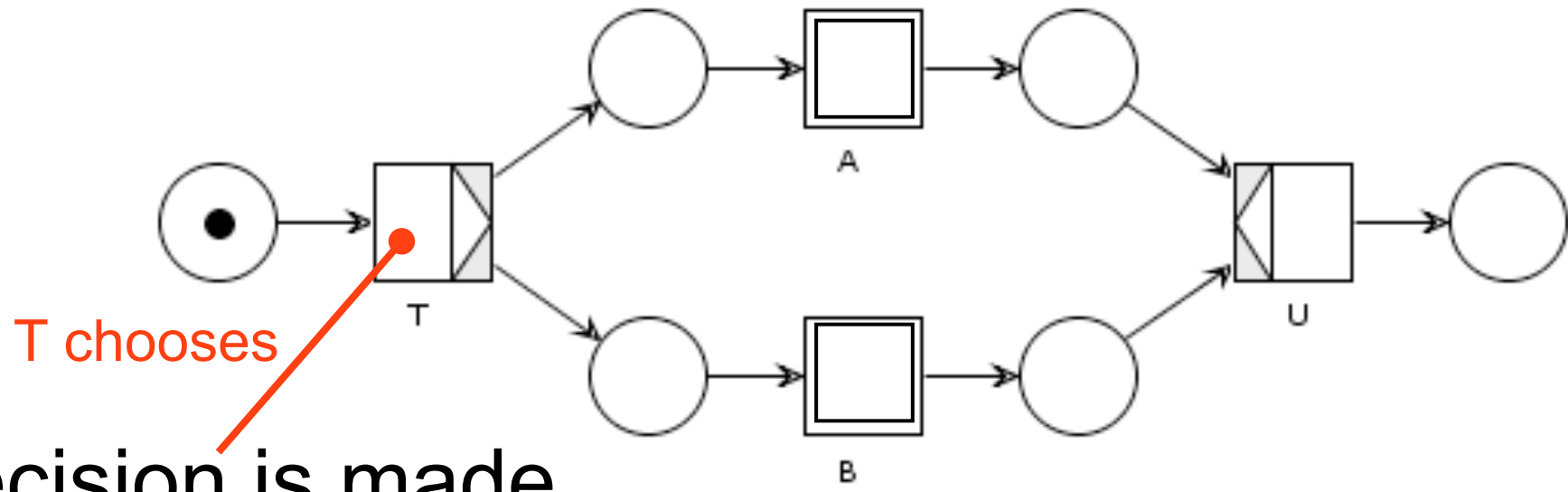
# Deferred choice

Either A or B is executed (choice is **implicit**)

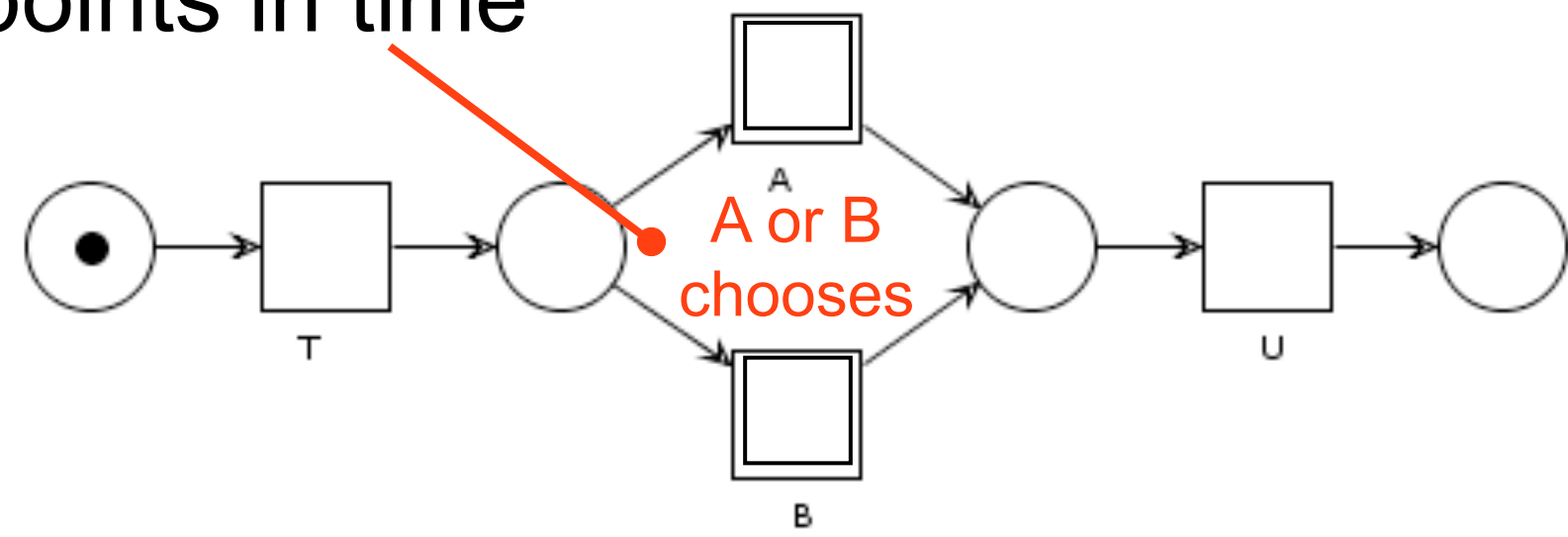


# Remember

Explicit choice  $\neq$  Implicit choice

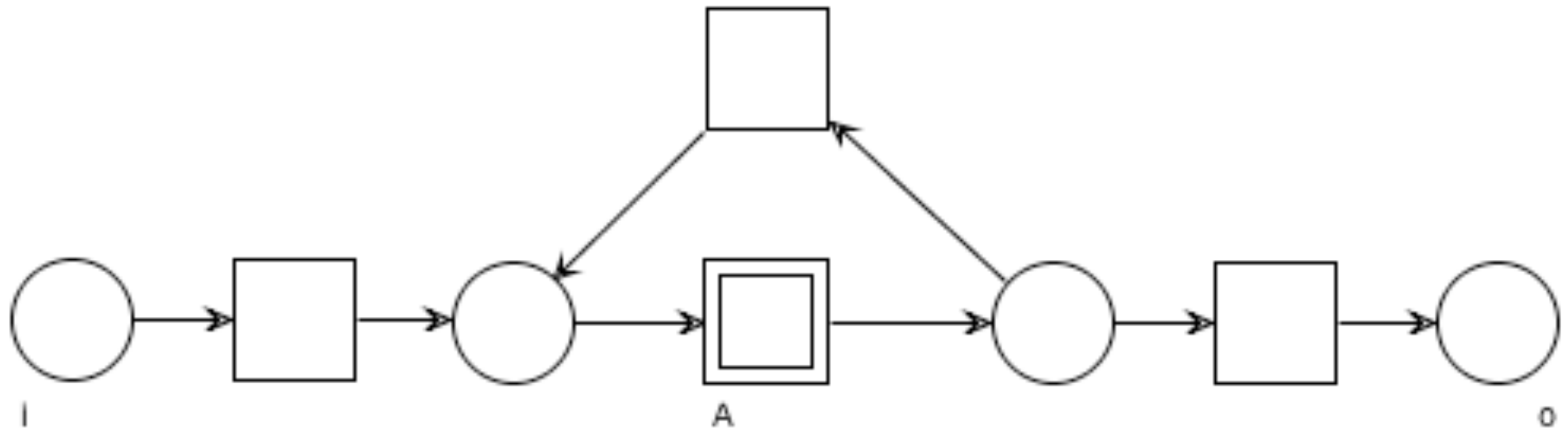


The decision is made at different points in time



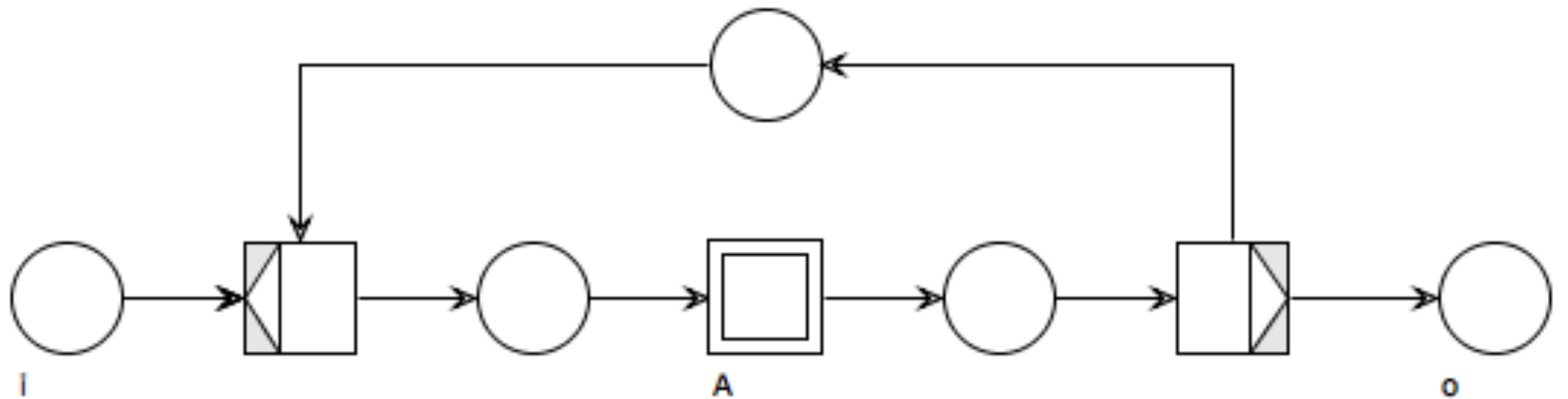
# Iteration (one or more times)

A is executed 1 or more times



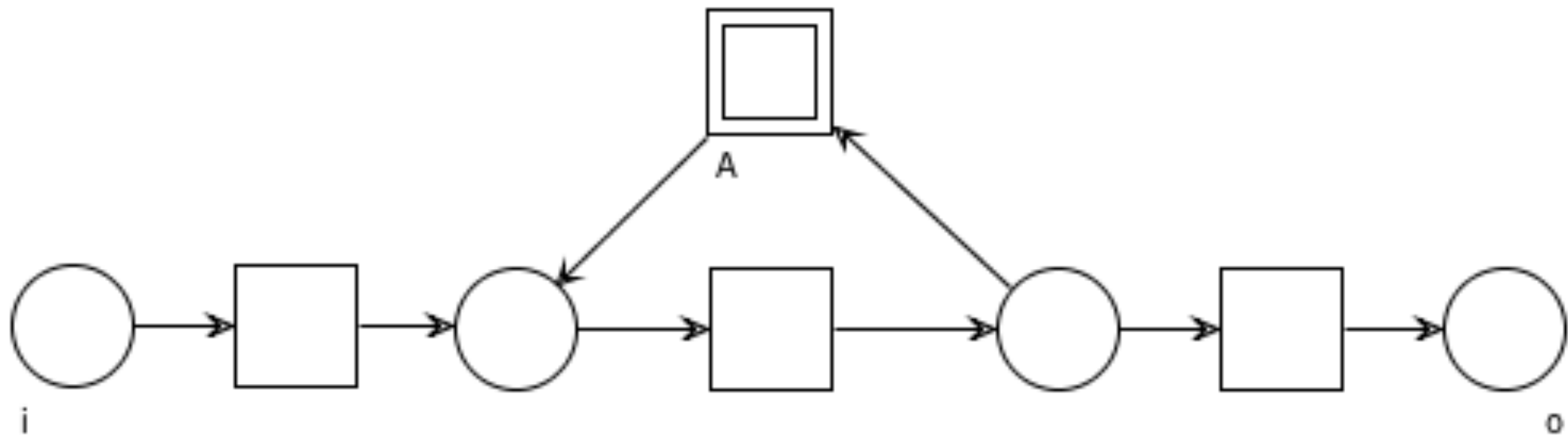
# One-or-more iteration ("sugared" version)

Decorated version for business process stakeholders



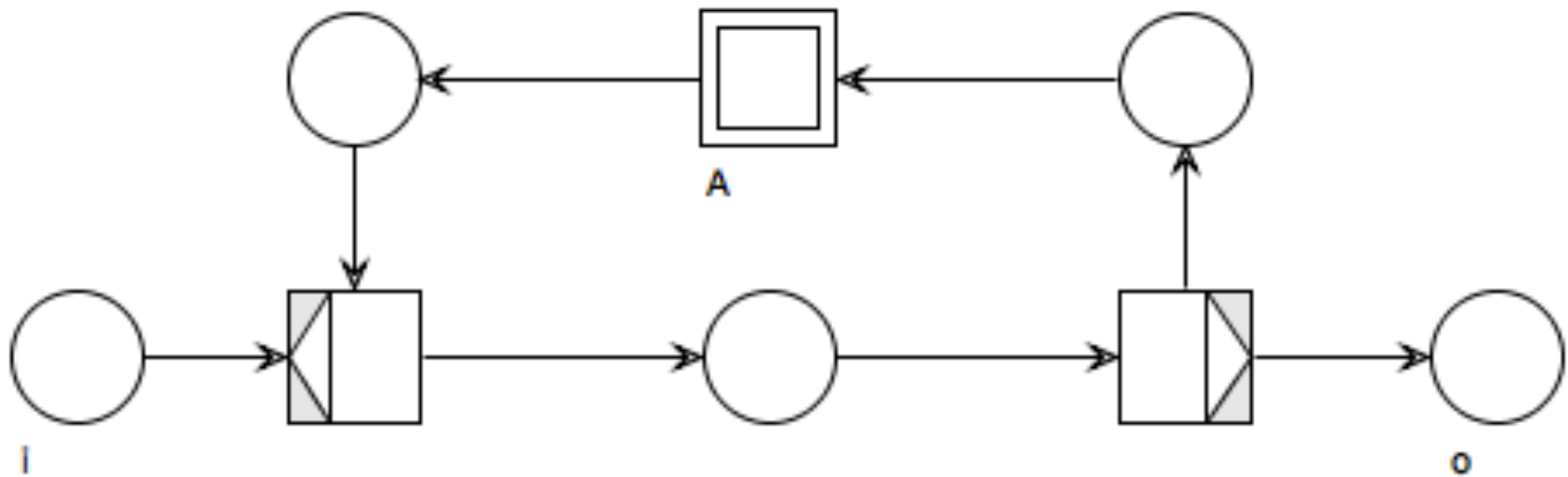
# Iteration (zero or more times)

A is executed 0 or more times



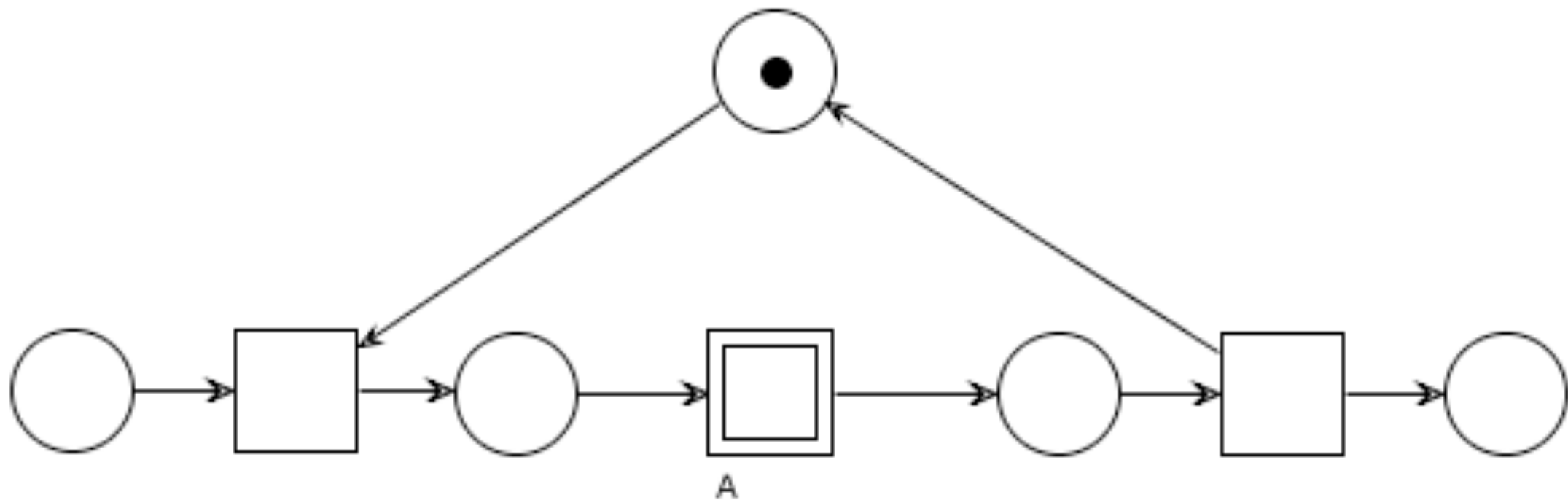
# Zero-or-more iteration ("sugared" version)

Decorated version for business process stakeholders



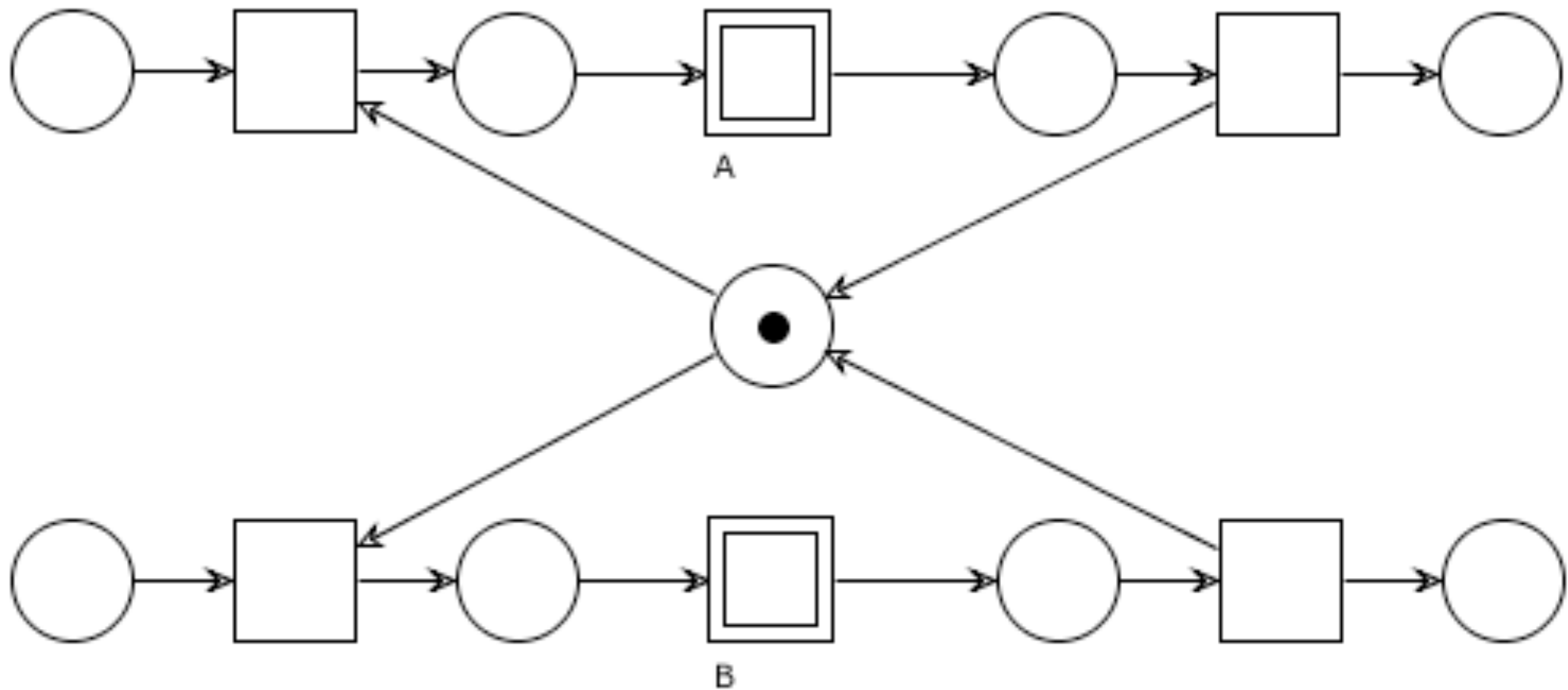
# One serve per time

Multiple activations are handled one by one



# Mutual exclusion

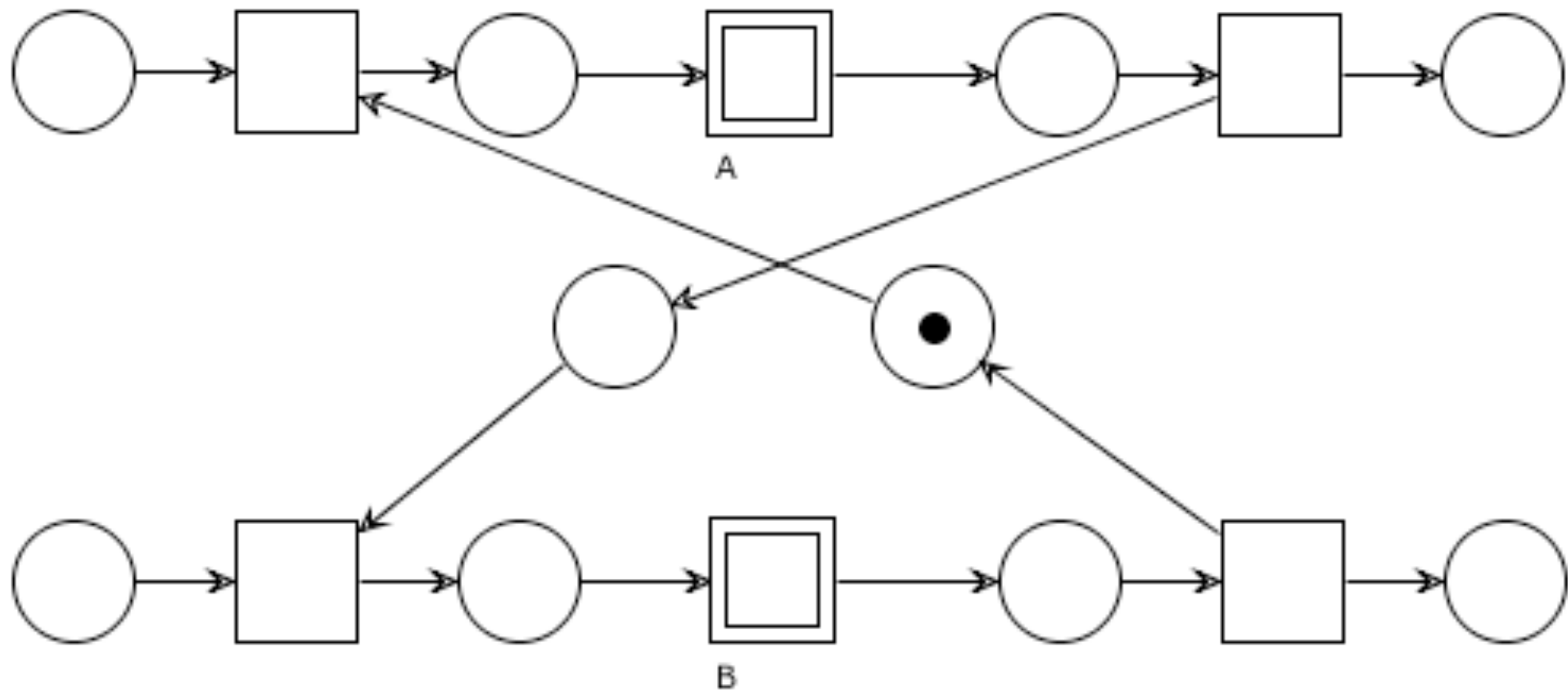
A and B cannot execute concurrently





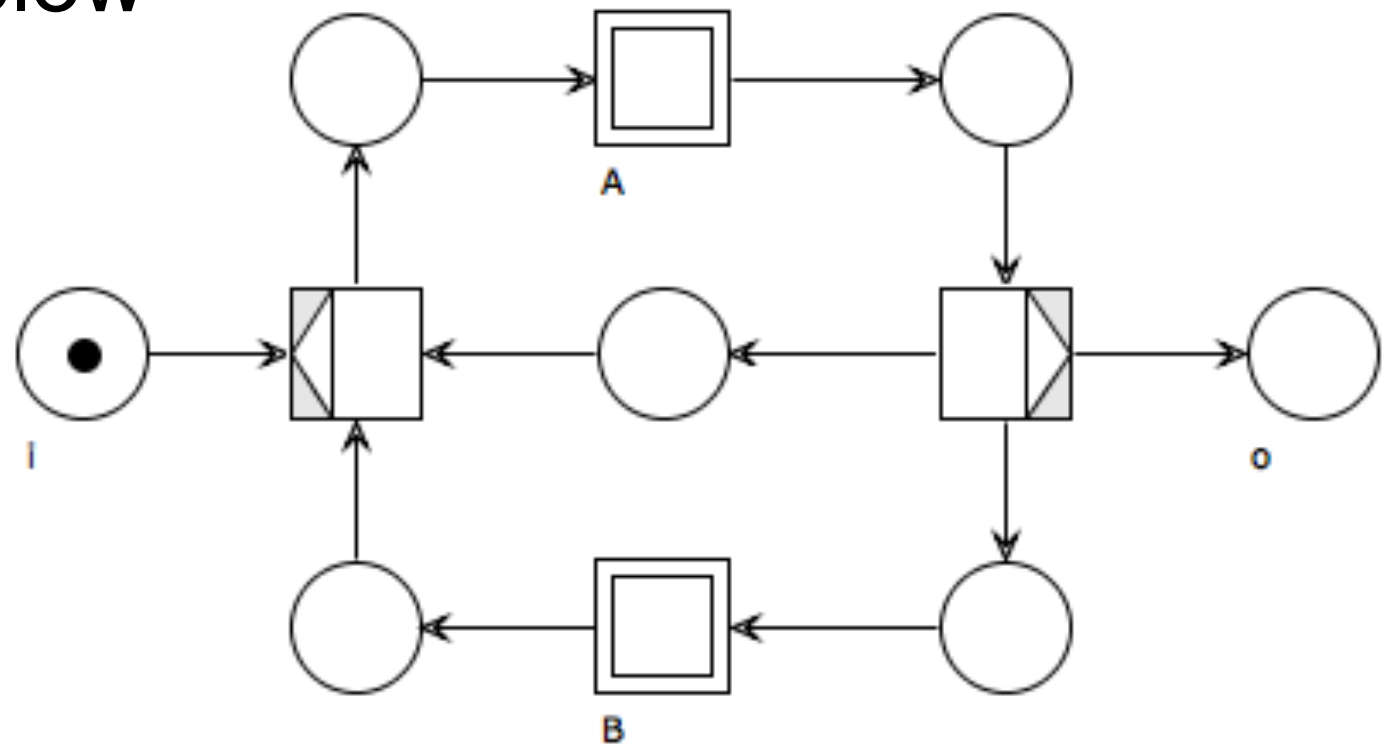
# Alternation

A and B execute one time each (A first)



# Question time

Consider the workflow net below



How many times can A be executed?

How many times can B be executed?

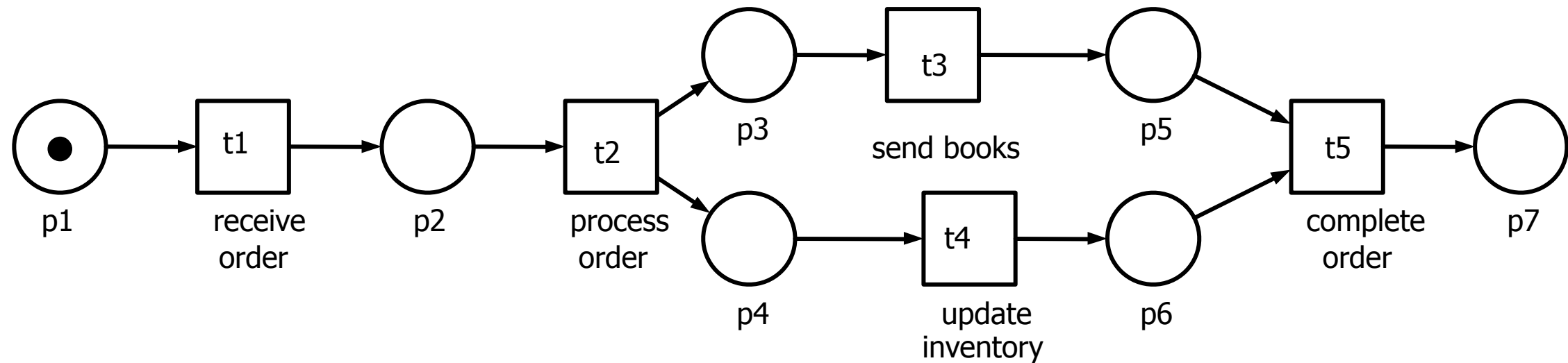
Can a firing sequence contain two As in a row?

Can a firing sequence contain two Bs in a row?

Can a firing sequence contain more Bs than As?

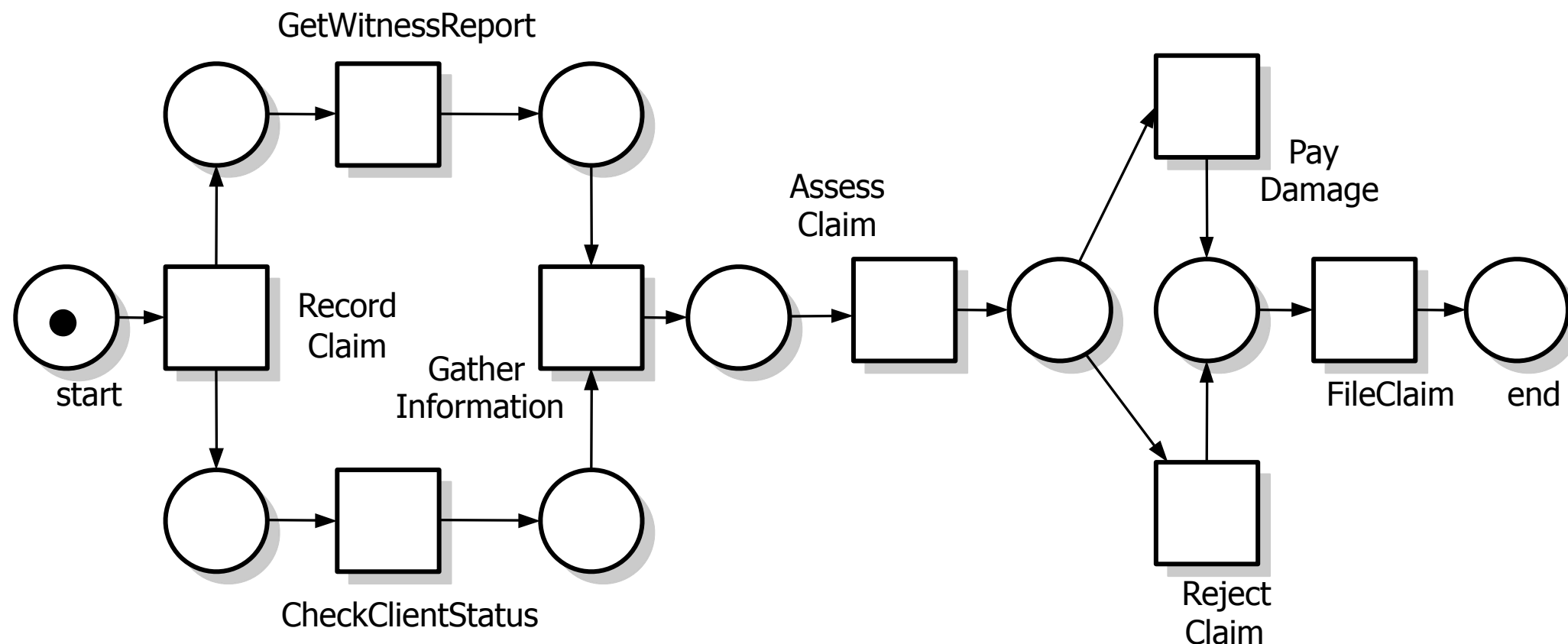
# Exercises

- Which "patterns" can be found in the workflow net below?
- Draw the corresponding Reachability Graph
- What is its language?



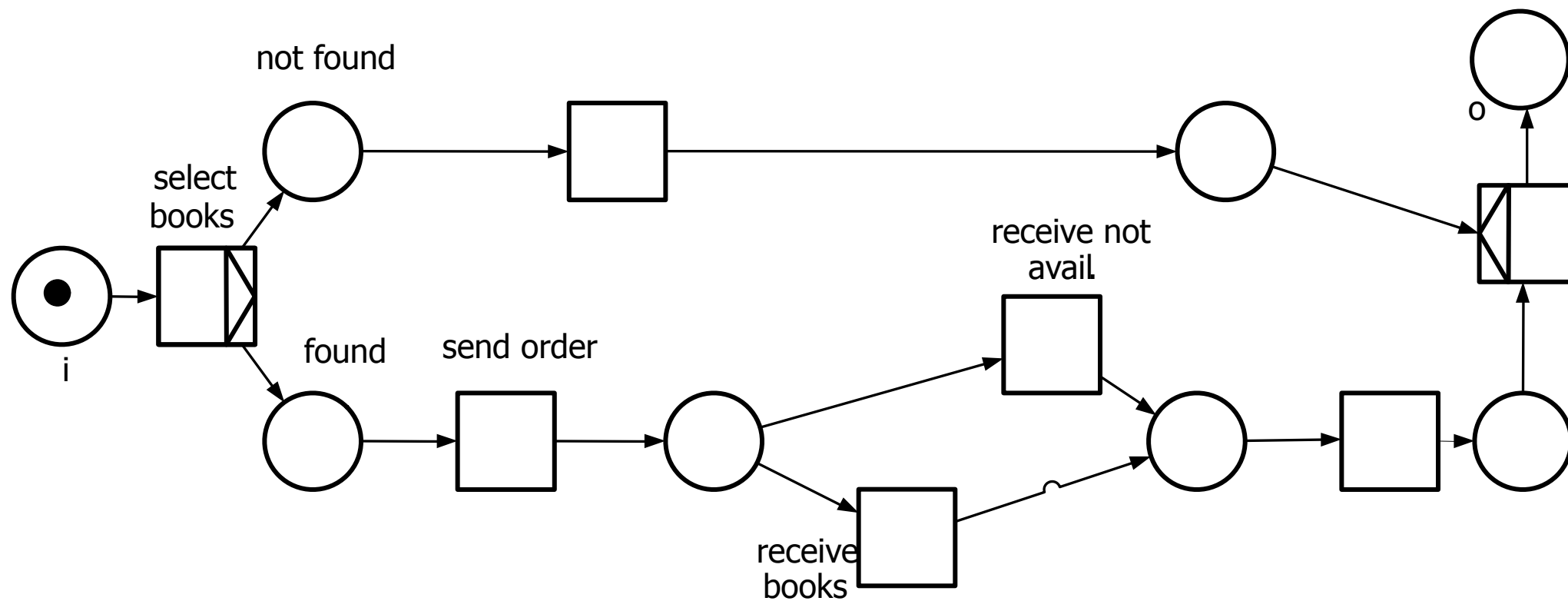
# Exercises

- Which "patterns" can be found in the workflow net below?
- "Sugarize" the net (where it makes sense)
- Name all places and draw the Reachability Graph
- What is its language?



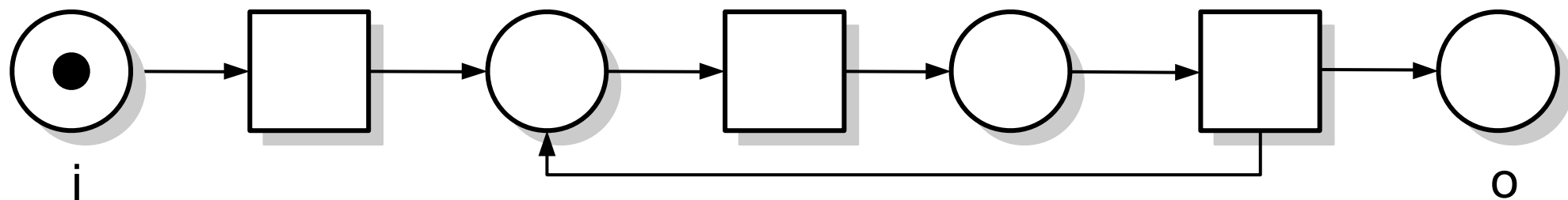
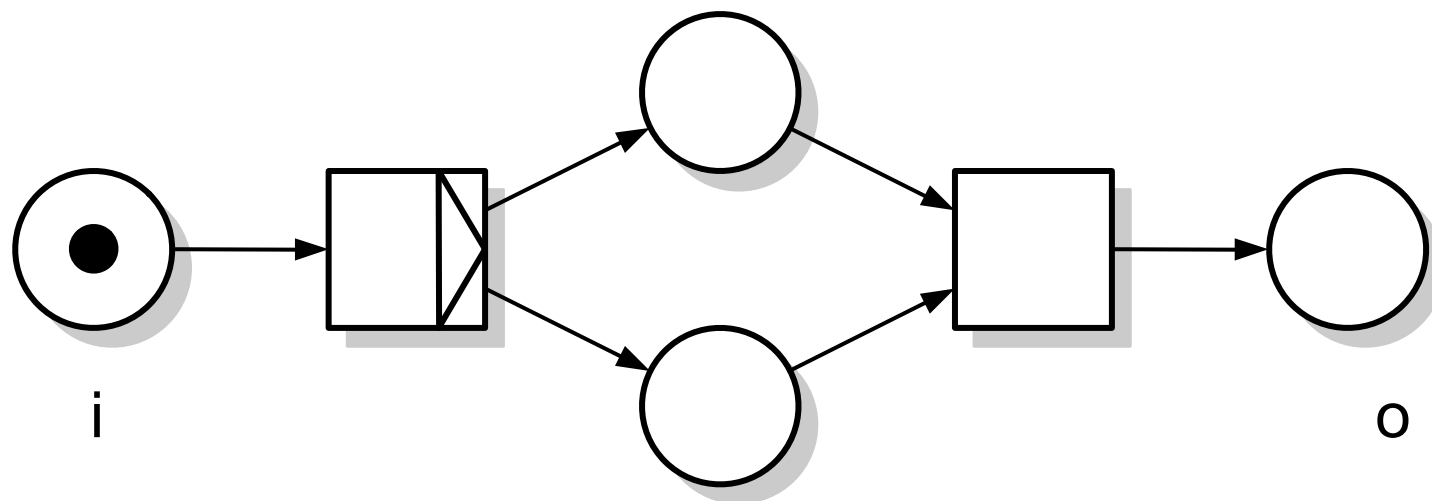
# Exercises

- "Desugarize" the workflow net below
- Name all nodes and draw the Reachability Graph
- What is its language?



# Exercises

- "Desugarize" the workflow nets below
- Name all nodes and draw the Reachability Graphs
- What are their languages?



# Triggers

Execution constraints can depend on the environment in which processes are enacted.

In workflow nets, transitions can be decorated with the information on who (or what) is responsible for the "firing" of that task.

Such annotations are called **triggers**

# Triggers

Triggers can be:

a human interaction

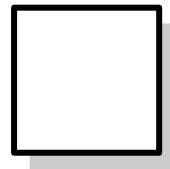
the receipt of a message

the expiration of a time-out

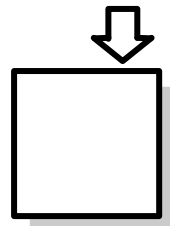
Transitions with no trigger can fire automatically



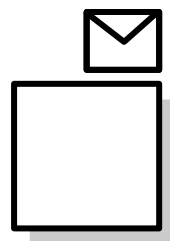
# Symbols for triggers



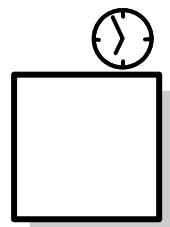
Automatic Trigger: Task enacted automatically



User Trigger: A human user takes initiative and starts activity

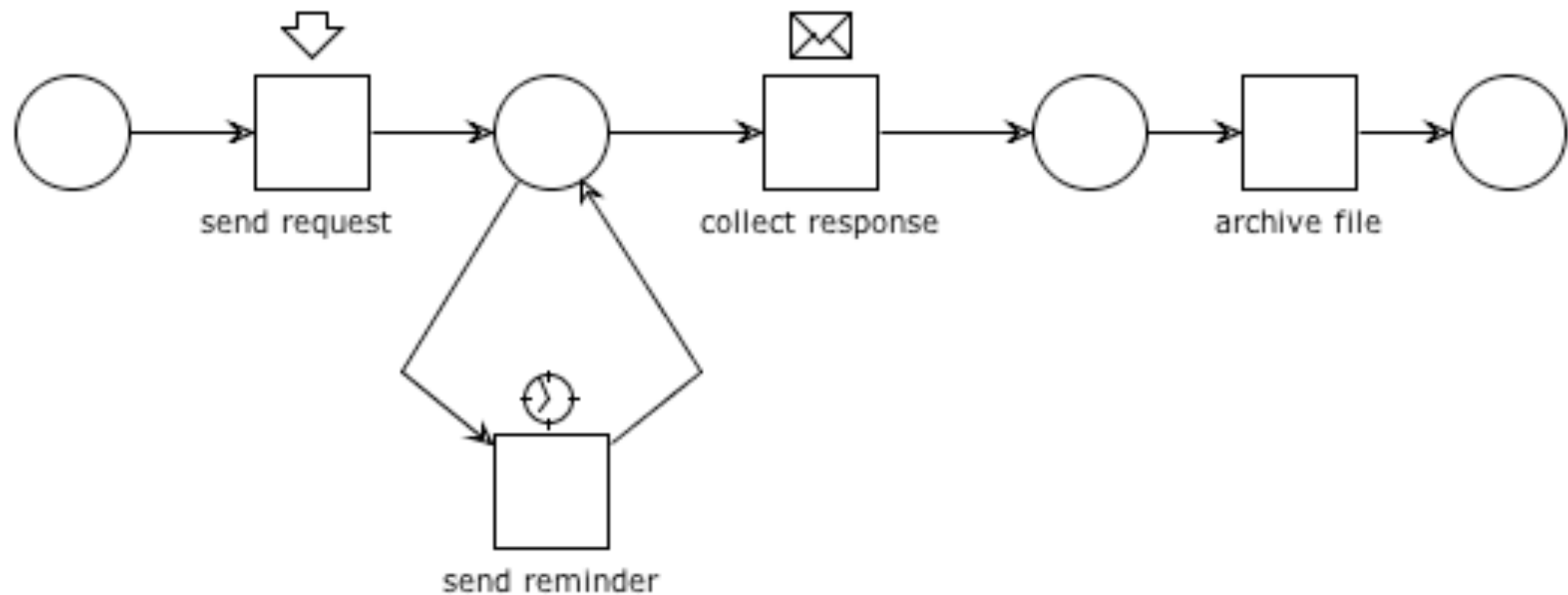


External Trigger: External event required to start activity

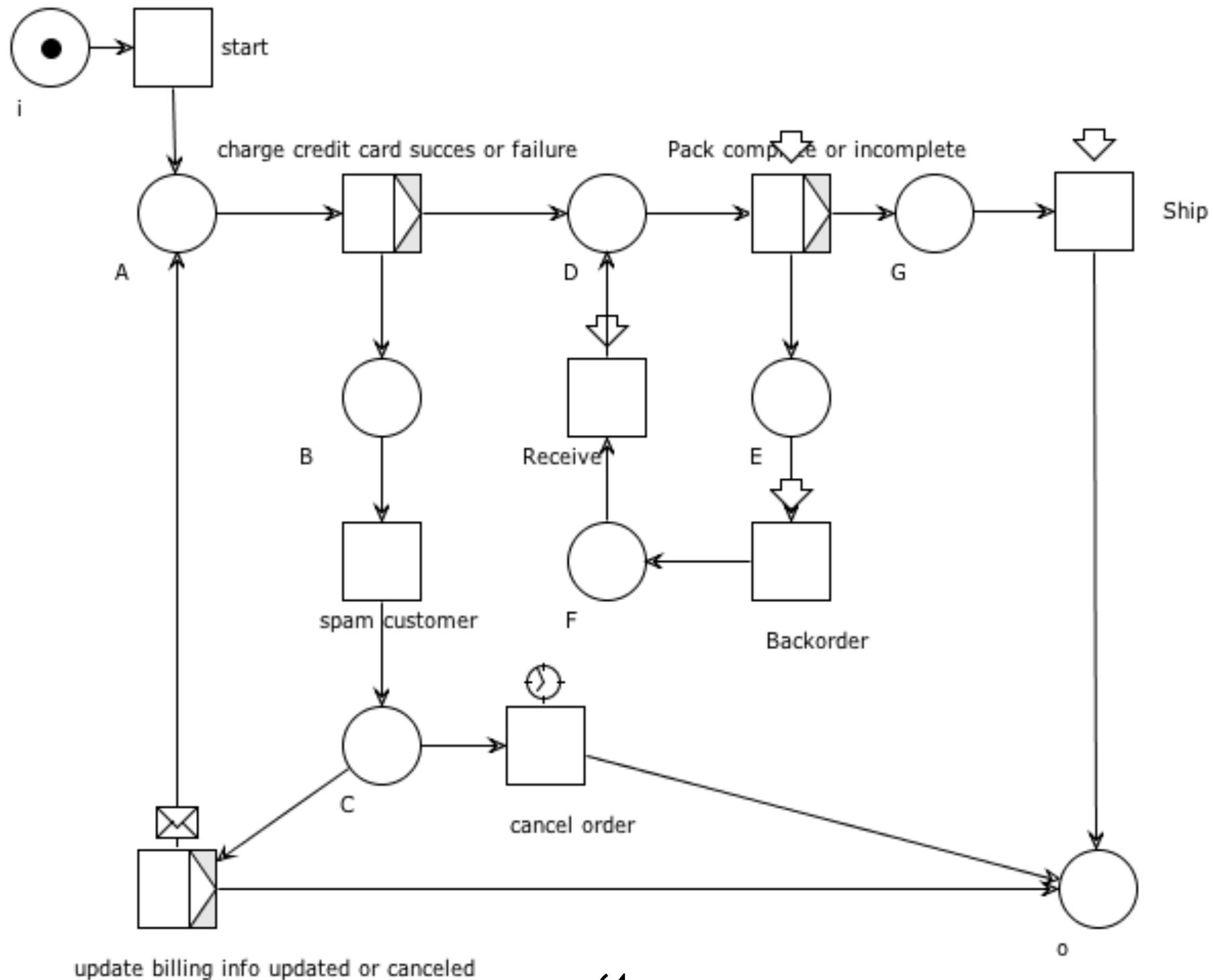


Time Trigger: Activity started when timer elapses

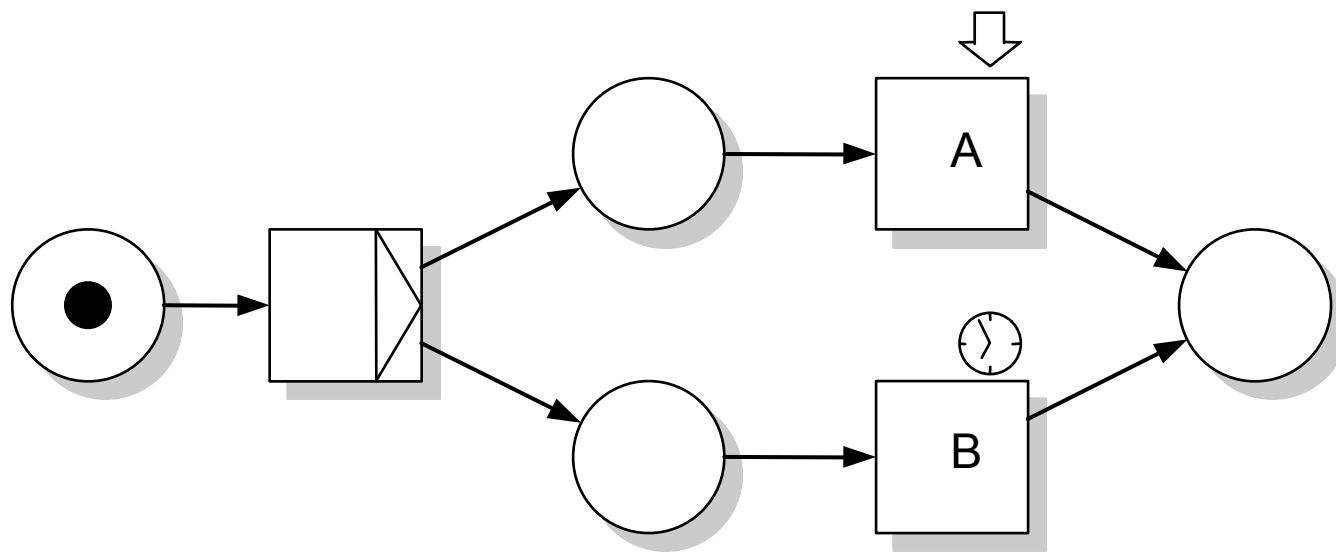
# Triggers: example



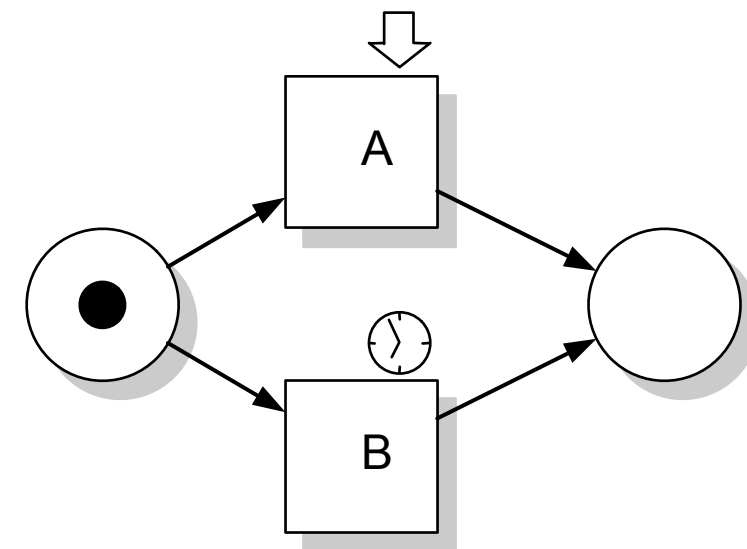
# Triggers: example



# Explicit vs Implicit XOR-split



(a) *Explicit xor split* does not enable A and B concurrently



(b) *Implicit xor split* enables A and B concurrently

# Motivation for the analysis

$L(N)$  shows the correct ways to run the process  
if it is empty there is clearly some problem

Are we guaranteed that nothing can go wrong?  
Are we guaranteed that once a case is started  
it will reach an end?

BPs are large, with increasing complexity  
flawed situations are frequent

