Methods for the specification and verification of business processes MPB (6 cfu, 295AA)

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14 - Sound by construction

Object

We show a technique to build sound Workflow nets

Soundness proof by construction

Idea

1. Find a suitable set of "building blocks"

they are (small) workflow nets that can be (easily) proved to be sound to be safe (1-bounded)

> 2. Define composition patterns so that by composing safe and sound WF nets we obtain safe and sound WF nets

Let N, N' be two sound and safe workflow nets



Let t be a task of N with exactly one input and one output place



Let N[N'/t] denote the net obtained by replacing the task t in N by N'



The net N[N'/t] is a sound and safe workflow net



Proof sketch

Intuitively

a sound workflow net behaves as a transition: it takes one token from its input place and it produces one token to its output place (but not atomically)

Formally

the crux of the proof is showing a bijective correspondence between markings of the composed net N[N'/t] and the pairs of markings in N and N'



implicit XOR

iteration

Some Building Blocks 2





Some Building Blocks 3



But you can define more blocks on your own















Exercise

Prove that the net below is a safe and sound workflow net



Exercise

Prove that the net below is a safe and sound workflow net (hint: "desugar" it)



Generalization

We would like to progressively refine transitions with multiple incoming and outgoing arcs



Two facts

Lemma: Let N be a sound WF net. If $(i,t) \in F$ then the pre-set of t is $\{i\}$

(otherwise t would be a dead transition)

Lemma: Let N be a sound WF net. If $(t,o) \in F$ then the post-set of t is $\{o\}$

(otherwise t would be dead or proper completion would not hold)

General replacement Let $T_{i'} = \{ u \mid \bullet u = \{i'\} \}$. Let $T_{o'} = \{ v \mid v \bullet = \{o'\} \}$.

If $(p,t) \in F_N, u \in T_{i'}$ then $(p,u) \in F_{N[N'/t]}$ If $(t,q) \in F_N, v \in T_{o'}$ then $(v,q) \in F_{N[N'/t]}$

