Business Processes Modelling
MPB (6 cfu, 295AA)
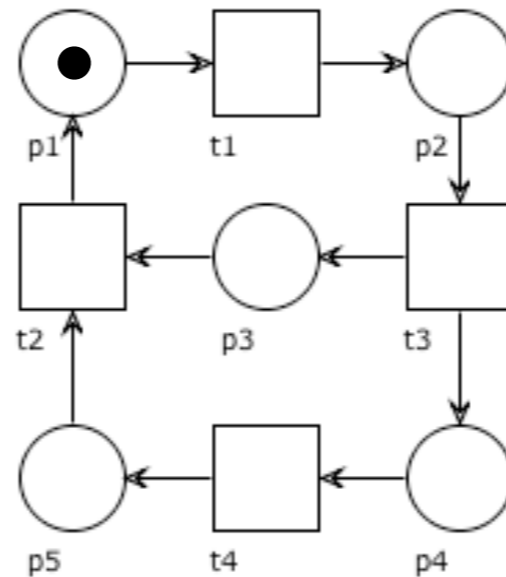
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17 - T-systems

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Object

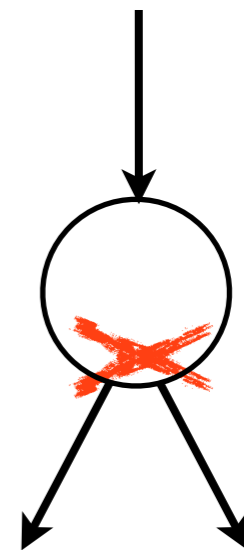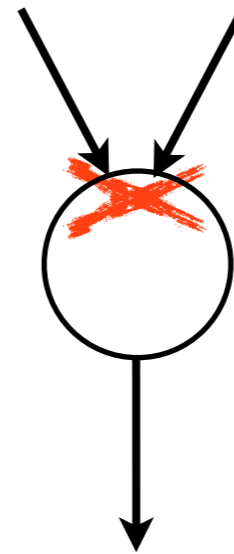We study some "good" properties of T-systems
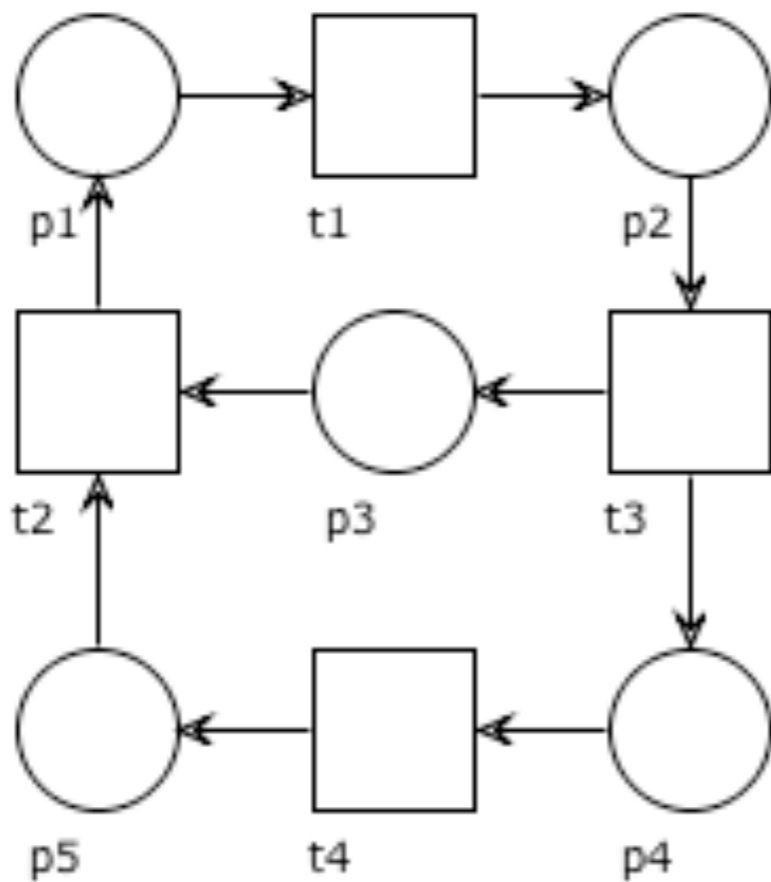
T-systems

T-system

Definition: We recall that a net N is a **T-net** if each place has exactly one input transition and exactly one output transition

$$\forall p \in P, \qquad |\bullet p| = 1 = |p \bullet|$$

A system (N,M_0) is a **T-system** if N is a T-net

T-net: example

T-systems: an observation

Notably, computation in T-systems is concurrent, but essentially deterministic:

the firing of a transition t in M cannot disable another transition t' enabled at M

T-net N*

Is the following conjecture true?
A workflow net N is a T-net
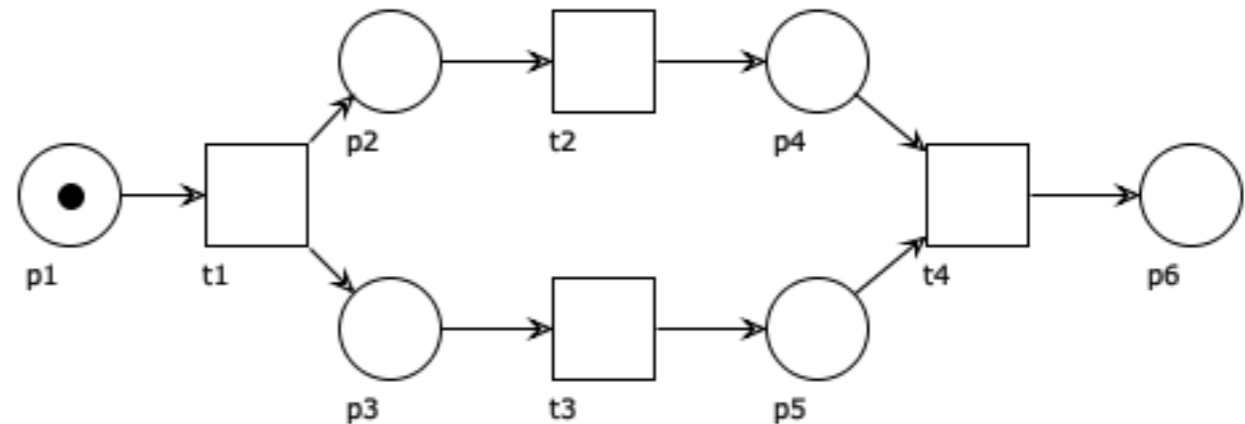iff N* is a T-net

T-net N*

Is the following conjecture true?
A workflow net N is a T-net
iff N* is a T-net

No, a workflow net cannot be a T-net because
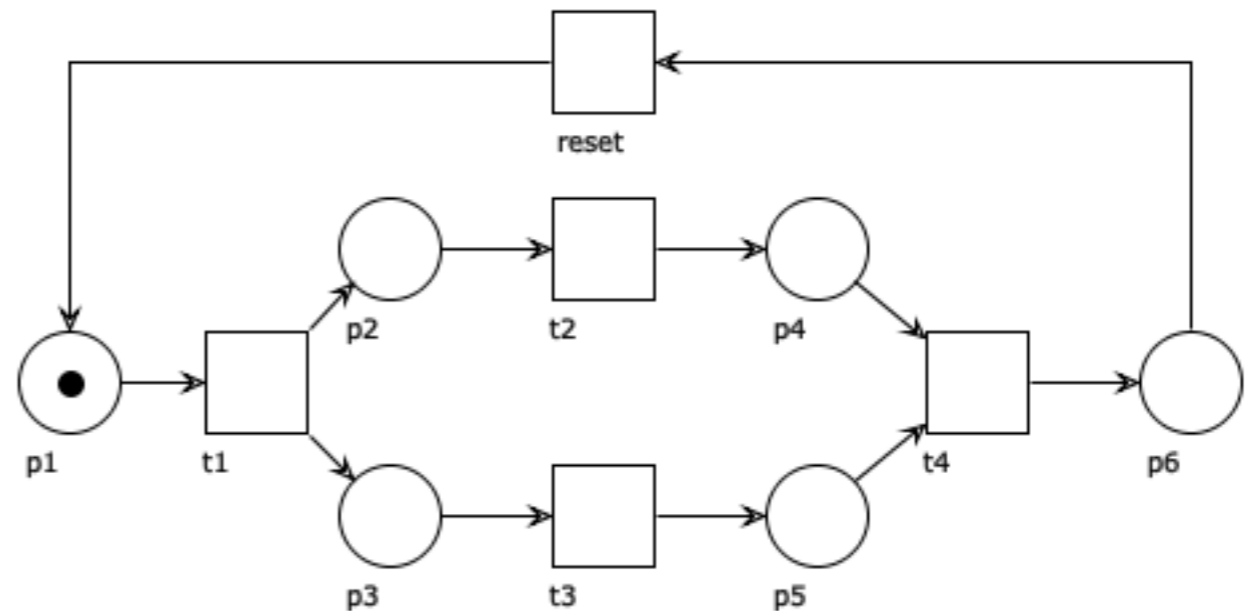the place i has no incoming arc
and the place o has no outgoing arc

(but **N* can be a T-net**)

T-net N*: example

N is a workflow net
but not a T-net

N* is a T-net
(not a workflow net)

T-systems: another observation

Determination of control:

the transitions responsible for enabling t are
one for each input place of t

Notation: token count
of a circuit

Let $\gamma = (x_1, y_1)(y_1, x_2)(x_2, y_2)...(x_n, y_n)$ be a circuit.
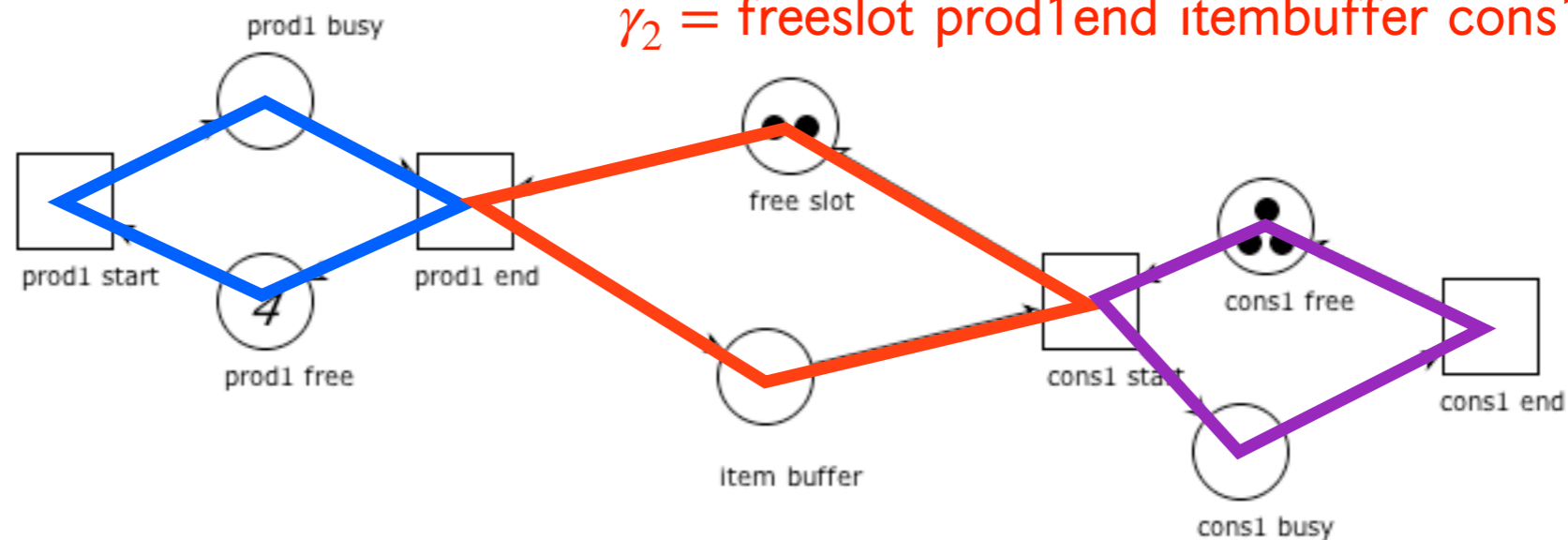
Let $P_{|\gamma} \subseteq P$ be the set of places in γ.

$$M(\gamma) = M(P_{|\gamma}) = \sum_{p \in P_{|\gamma}} M(p)$$

We say that γ is **marked at** M if $M(\gamma) > 0$

Example

γ_1 = prod1busy prod1end prod1free prod1start

γ_2 = freeslot prod1end itembuffer cons1start

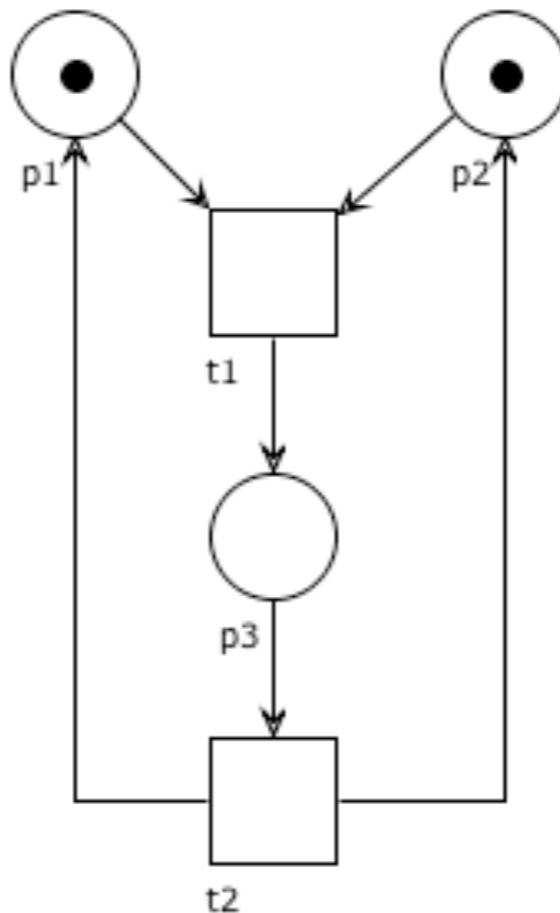γ_3 = cons1busy cons1end cons1free cons1start

$$M(\gamma_1) = 4$$

$$M(\gamma_2) = 2$$
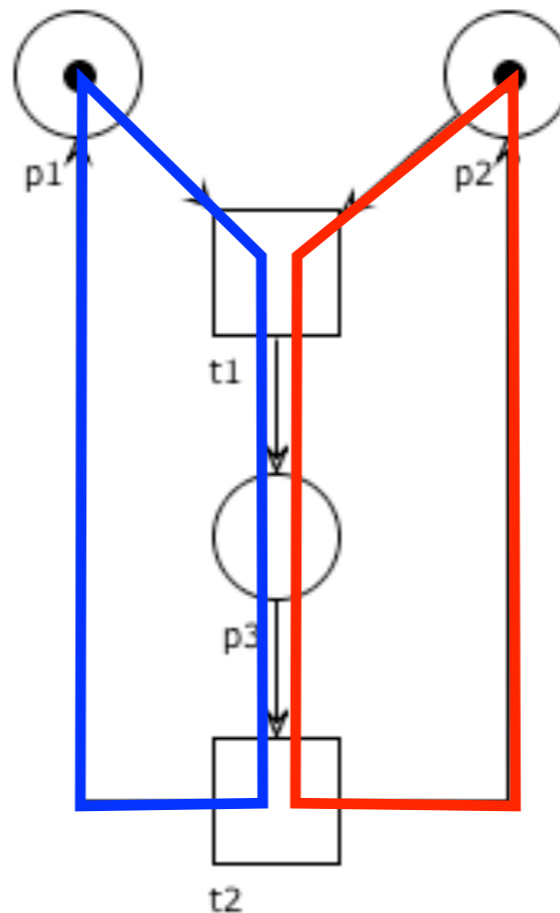
$$M(\gamma_3) = 3$$

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Question time

Trace two circuits over the T-system below

Question time

Trace two circuits over the T-system below

Fundamental property of T-systems

The token count of a circuit is invariant under any firing.

Fundamental property of T-systems

Proposition: Let γ be a circuit of a T-system (P, T, F, M_0).
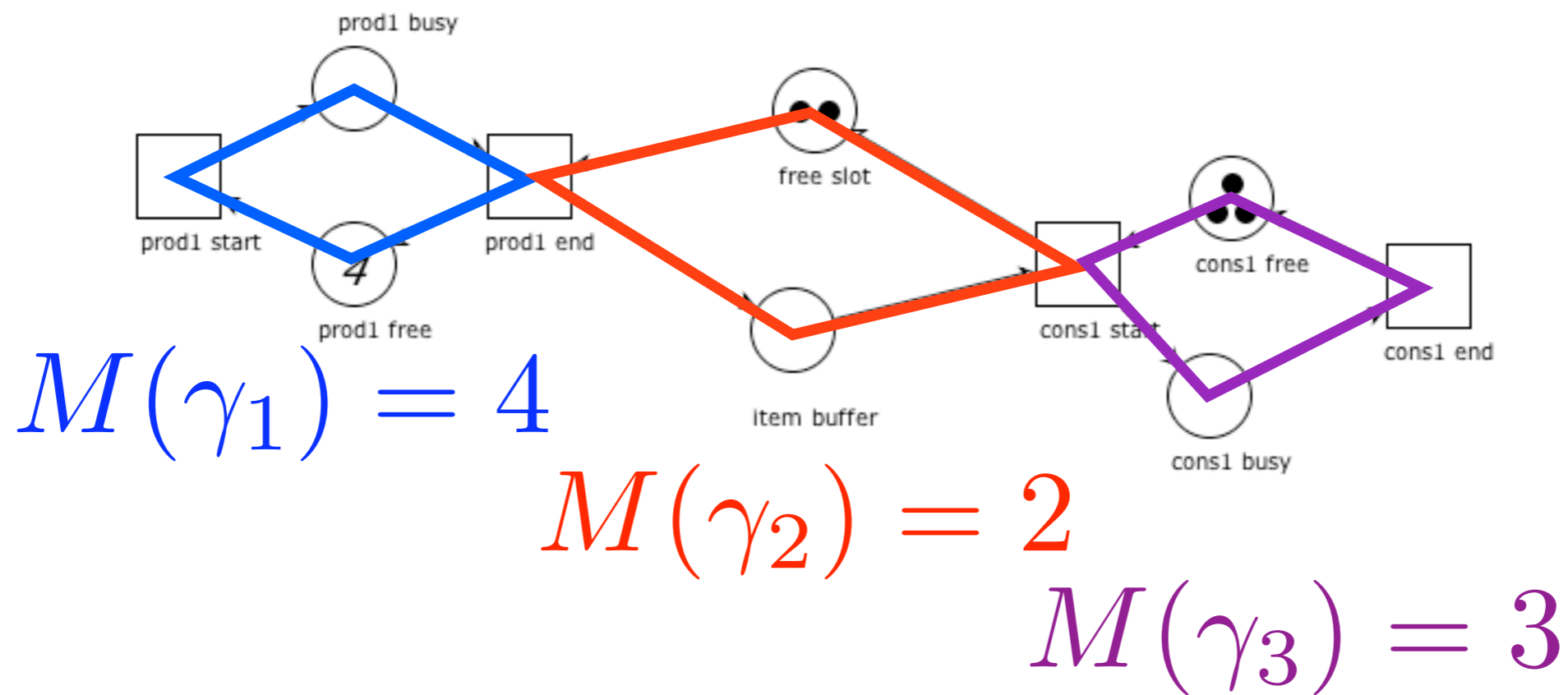If M is a reachable marking, then $M(\gamma) = M_0(\gamma)$

Take any $t \in T$: either $t \notin \gamma$ or $t \in \gamma$.

If $t \notin \gamma$, then no place in $\bullet t \cup t \bullet$ is in γ
(otherwise, by definition of T-nets, t would be in γ).
Then, an occurrence of t does not change the token count of γ.

If $t \in \gamma$, then exactly one place in $\bullet t$ and one place in $t \bullet$ are in γ.
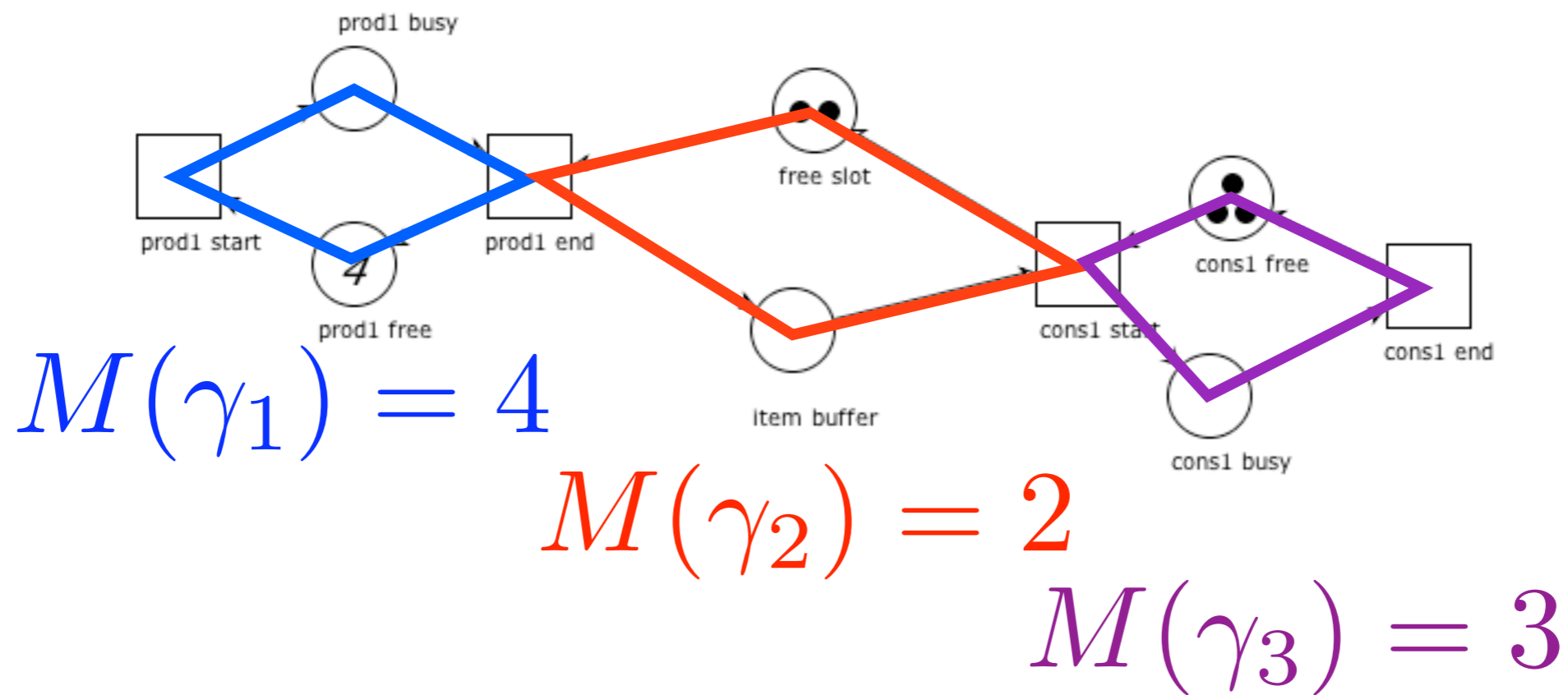Then, an occurrence of t does not change the token count of γ.

Example

$$M(\gamma_1) = 4$$

$$M(\gamma_2) = 2$$

$$M(\gamma_3) = 3$$

M₀ = [0 4 2 0 3 0]
M = [2 2 1 2 2 1]   Not reachable!

Example

$$M(\gamma_1) = 4$$

$$M(\gamma_2) = 2$$
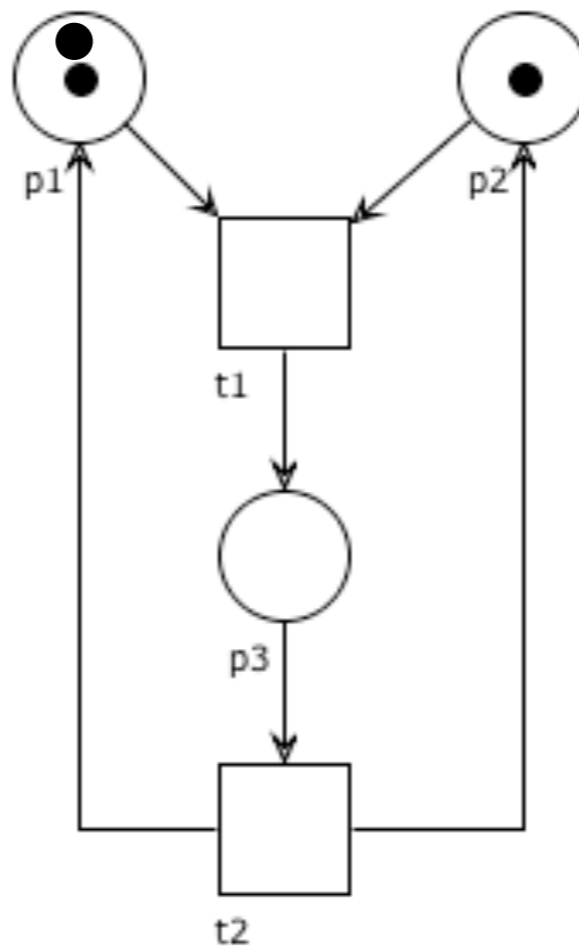
$$M(\gamma_3) = 3$$

$M_0 = [\ 0\ \ 4\ \ 2\ \ 0\ \ 3\ \ 0\ ]$

$M' = [\ 2\ \ 1\ \ 1\ \ 1\ \ 2\ \ 2\ ]$ Not reachable!

Question time

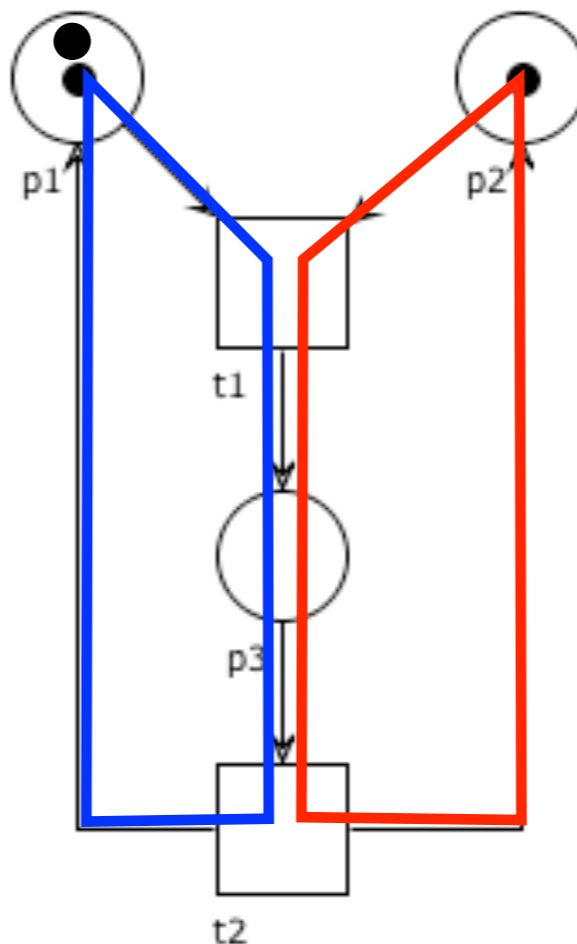Is the marking $p_1 + 2p_2$ reachable? (why?)

Question time

Is the marking $p_1 + 2p_2$ reachable? (why?)

No
the token count in
the left (blue) circuit
must remain 2
and
the token count in
the right (red) circuit
must remain 1

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T-invariants of T-nets

Proposition: Let N=(P,T,F) be a (connected) T-net. **J** is a T-invariant of N **iff J**=[k ... k] for some value k

(the proof is dual to the analogous proposition for S-invariants of S-nets)

Boundedness in strongly connected T-systems

Lemma: If a T-system (N,M$_0$) is strongly connected, then it is bounded

Let Γ be the set of the circuits of N and let $k = \max_{\gamma \in \Gamma} M_0(\gamma)$.

Since N is strongly connected, every place p belongs to some circuit γ_p.

By the fundamental property of T-systems: token count of γ_p is invariant.

Thus, for any reachable marking M, we have $M(p) \leq M(\gamma_p) = M_0(\gamma_p) \leq k$.
Hence the net is k-bounded.

Safeness in strongly connected T-systems

Corollary: If a T-system (N,M$_0$) is strongly connected and M$_0$(P)=1, then it is safe

Let Γ be the set of the circuits of N and let $k = \max_{\gamma \in \Gamma} M_0(\gamma)$=1

Since N is strongly connected, every place p belongs to some circuit γ_p.

By the fundamental property of T-systems: token count of γ_p is invariant.

Thus, for any reachable marking M, we have $M(p) \leq M(\gamma_p) = M_0(\gamma_p) \leq k$.
Hence the net is k-bounded.   =1

Safeness in strongly connected T-systems

Corollary: If a T-system (N,M_0) is strongly connected and for any circuit γ $M_0(\gamma)\leq 1$, then it is safe

Let Γ be the set of the circuits of N and let $k = \max_{\gamma \in \Gamma} M_0(\gamma)$. ≤ 1

Since N is strongly connected, every place p belongs to some circuit γ_p.
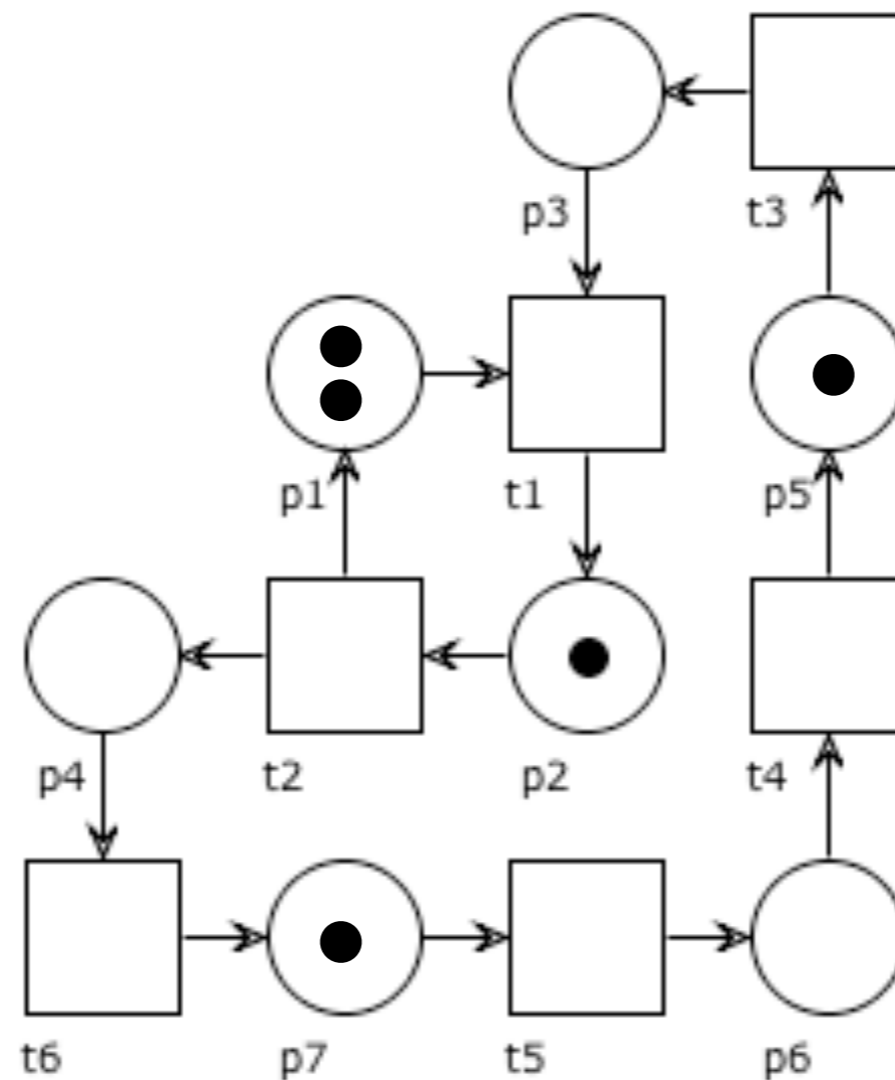
By the fundamental property of T-systems: token count of γ_p is invariant.

Thus, for any reachable marking M, we have $M(p) \leq M(\gamma_p) = M_0(\gamma_p) \leq k$. Hence the net is k-bounded. ≤ 1
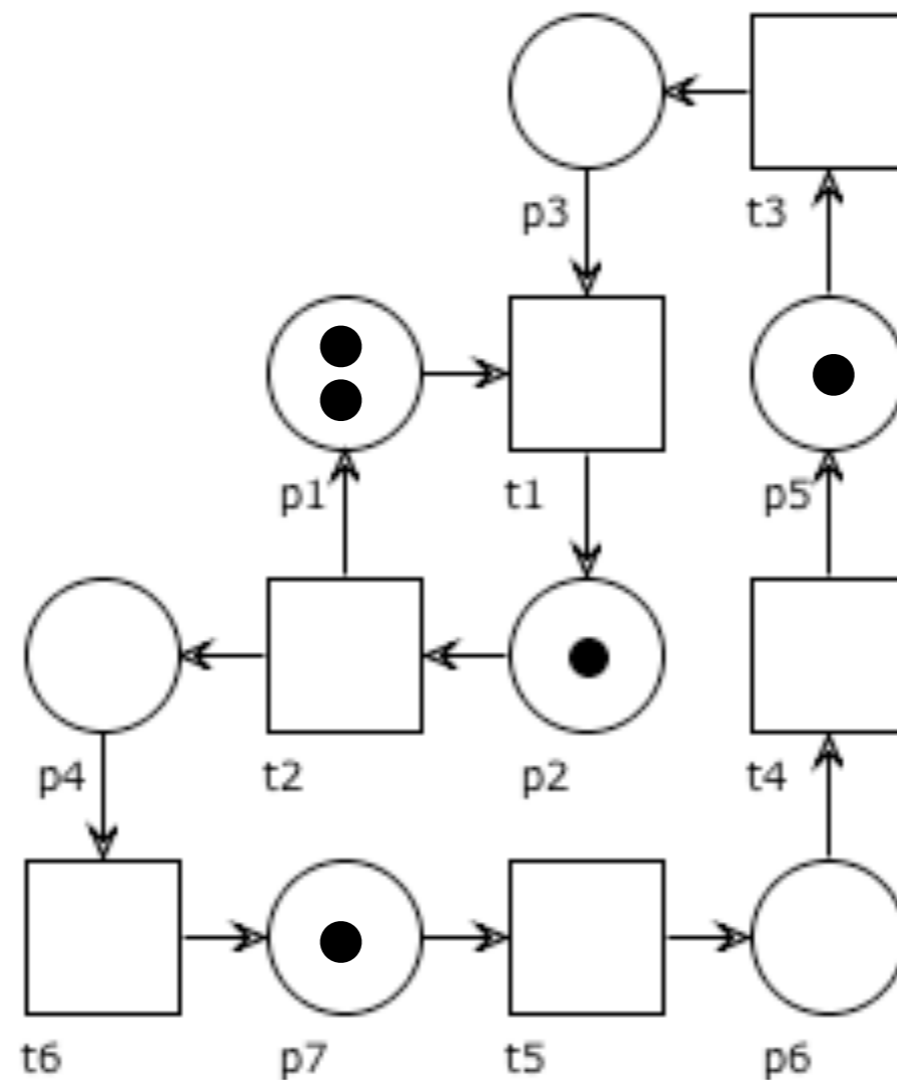
Question time

Is the T-systems below bounded? (why?)
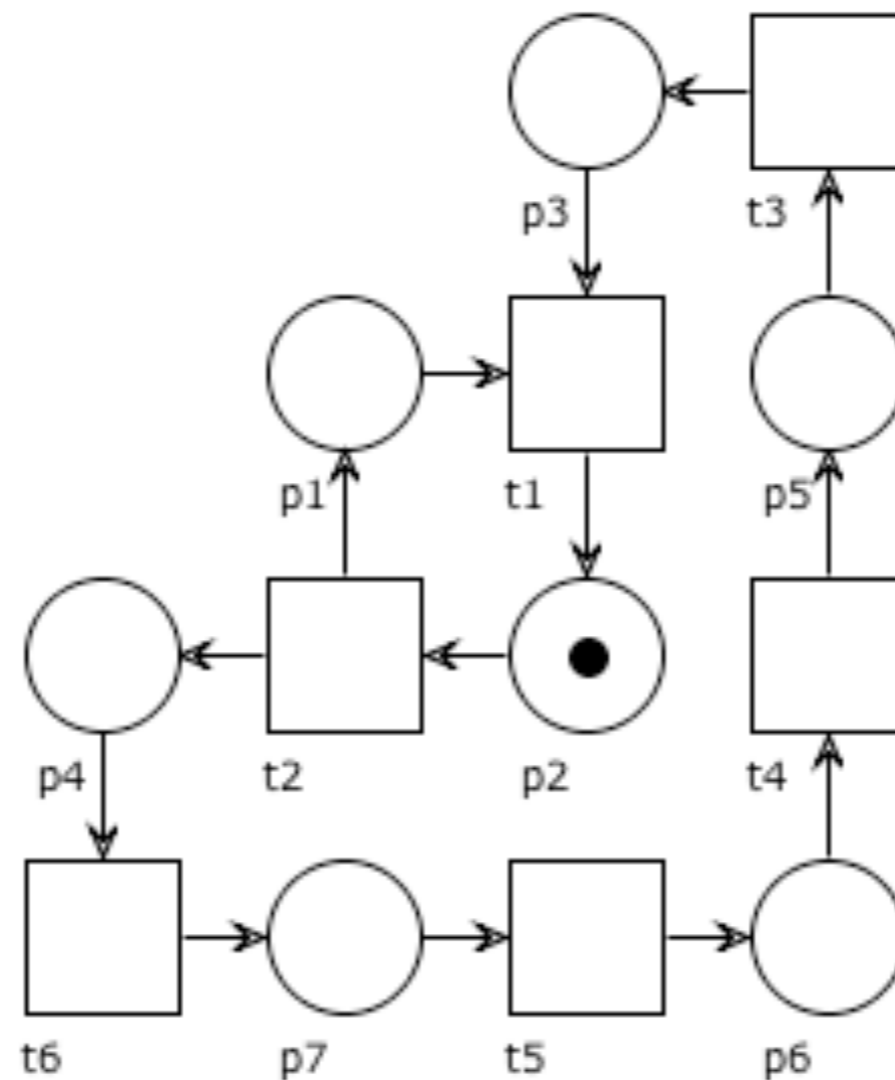
Question time

Is the T-systems below bounded? (why?)
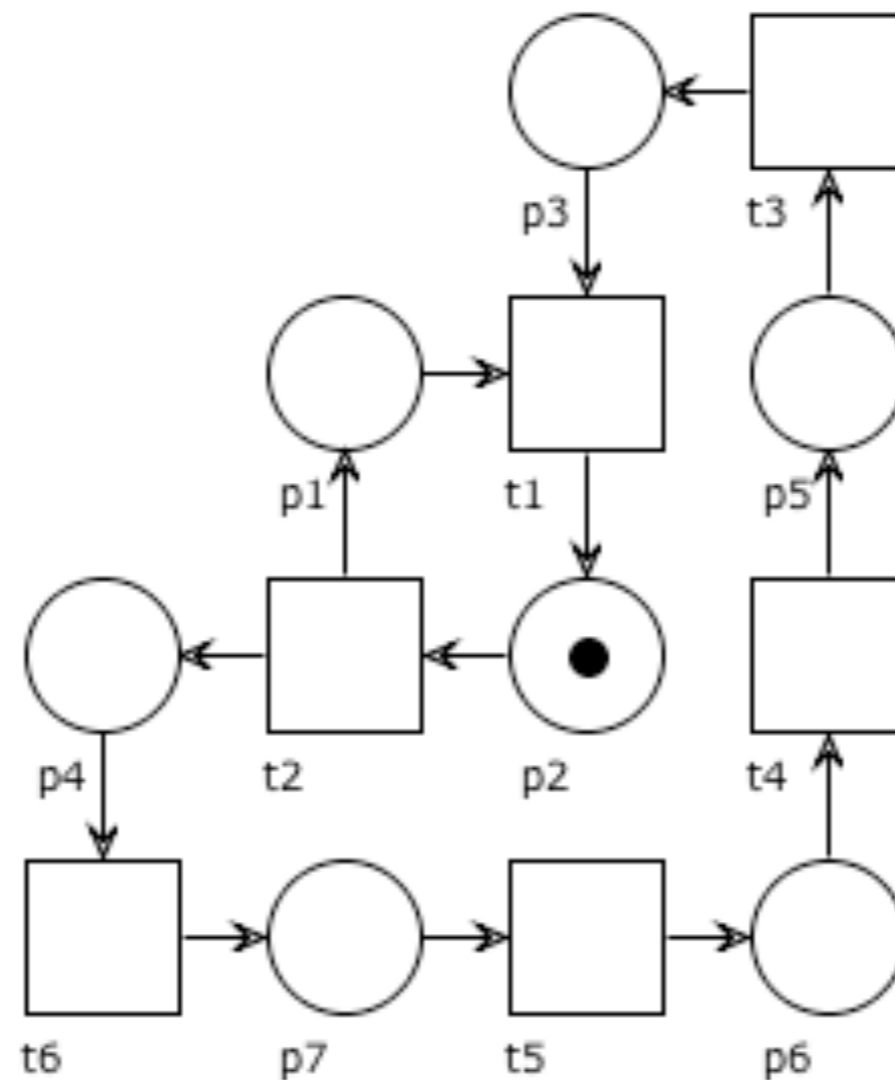

Strongly connected
=> bounded

Question time

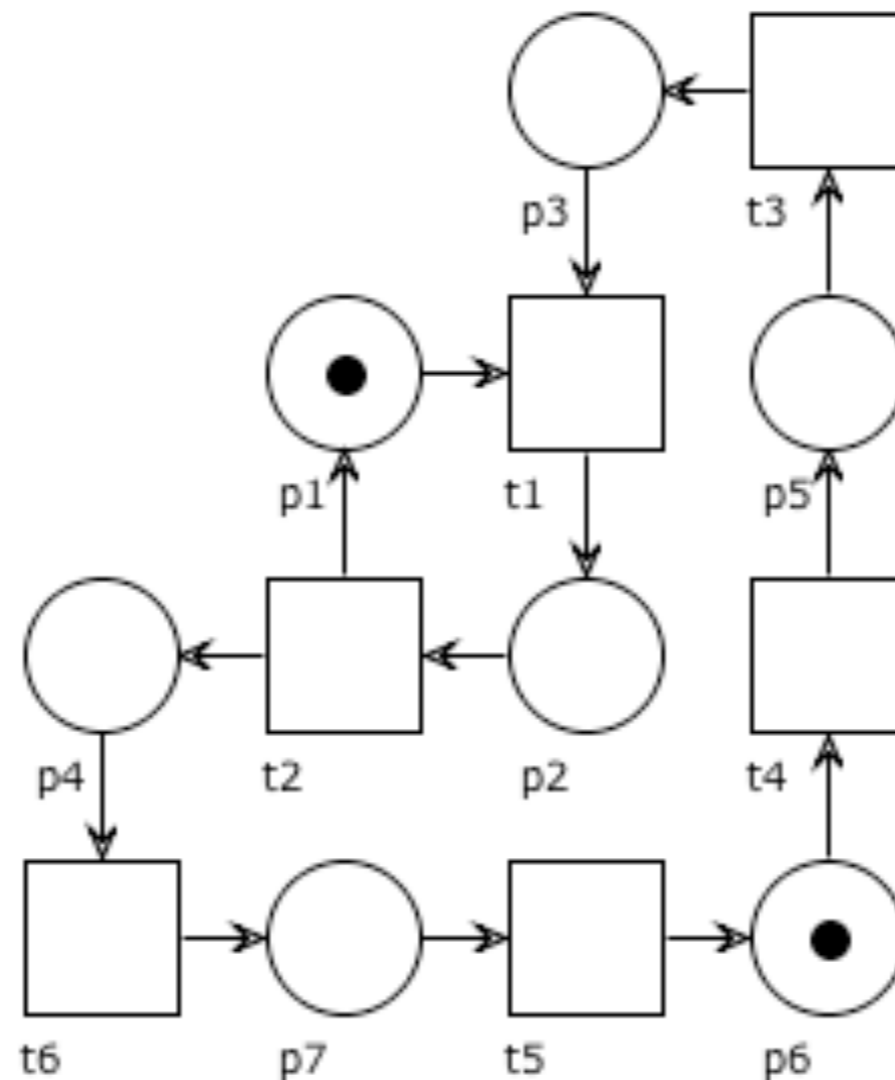Is the T-systems below safe? (why?)

Question time

Is the T-systems below safe? (why?)
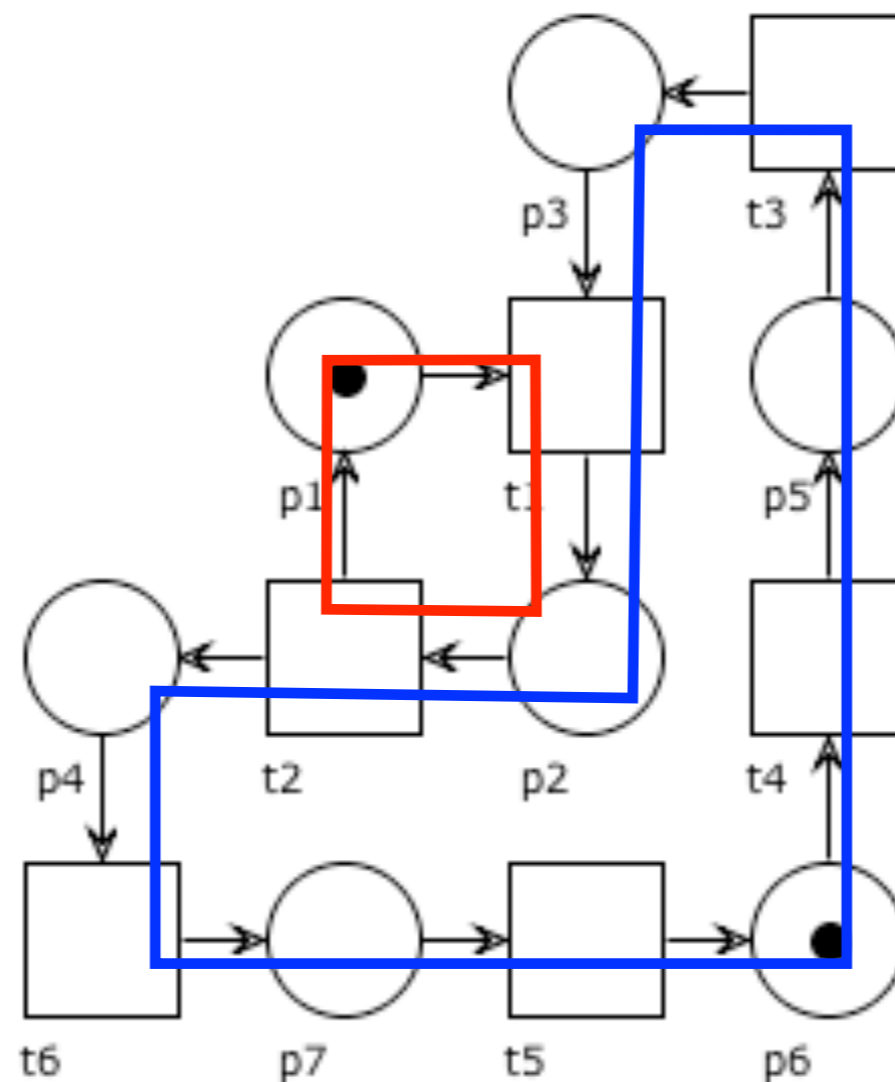

Strongly connected
+ $M_0(P)=1$
=> safe

Question time

Is the T-systems below safe? (why?)

Question time

Is the T-systems below safe? (why?)

Strongly connected
+ $\forall \gamma . M_0(\gamma) \leq 1$
=> safe

Liveness theorem for T-systems

Theorem: A T-system (N,M_0) is live
iff every circuit of N is marked at M_0

$\Rightarrow$) (quite obvious)
By contradiction, let γ be a circuit with $M_0(\gamma) = 0$.
By the fundamental property of T-systems: $\forall M \in [\, M_0 \,\rangle, \; M(\gamma) = 0$.

Take any $t \in T_{|\gamma}$ and $p \in P_{|\gamma} \cap \bullet t$.

For any $M \in [\, M_0 \,\rangle$, we have $M(p) = 0$.
Hence t is never enabled and the T-system is not live.

Liveness theorem for T-systems

Theorem: A T-system (N, M_0) is live
iff every circuit of N is marked at M_0

$\Leftarrow$) (more involved)
Take any $t \in T$ and $M \in [\, M_0 \,\rangle$.
We need to show that some marking M' reachable from M enables t.

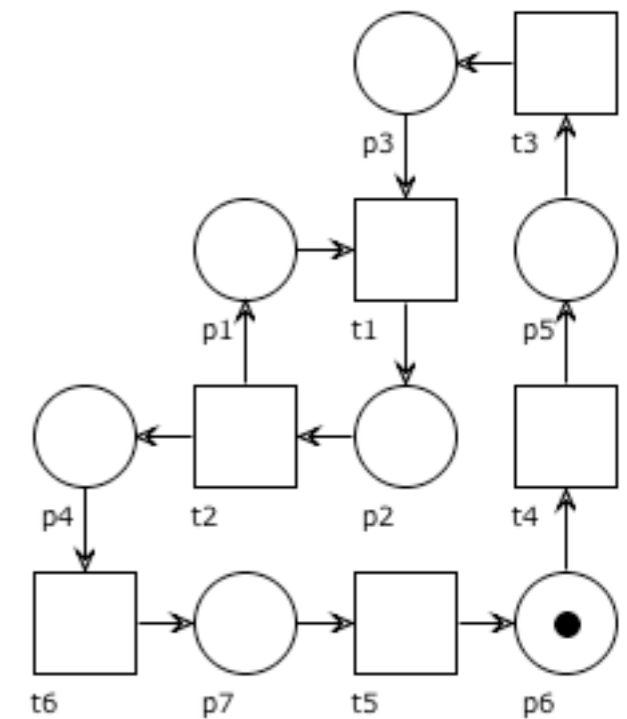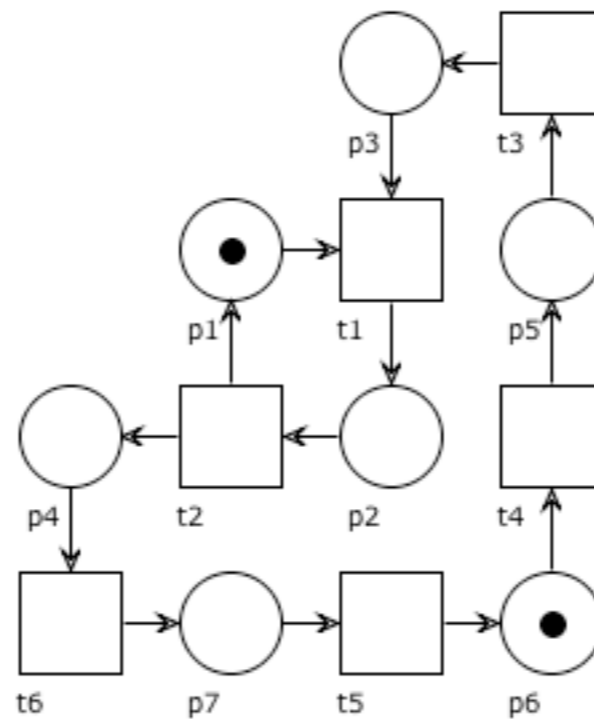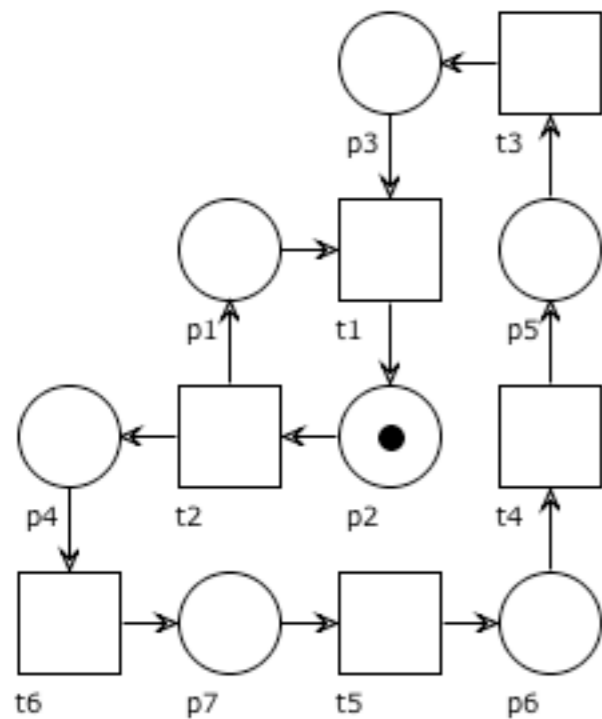The key idea is to collect the places that control the firing of t:
$p \in P_{M,t}$ if there is a path from p to t through places unmarked at M.
We then proceed by induction on the size of $P_{M,t}$.

Rest of the proof left to
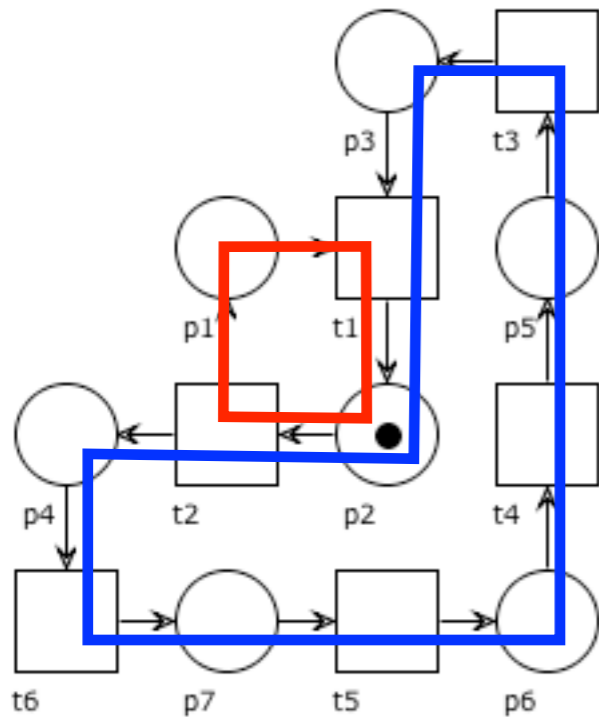your ingenuity

Question time

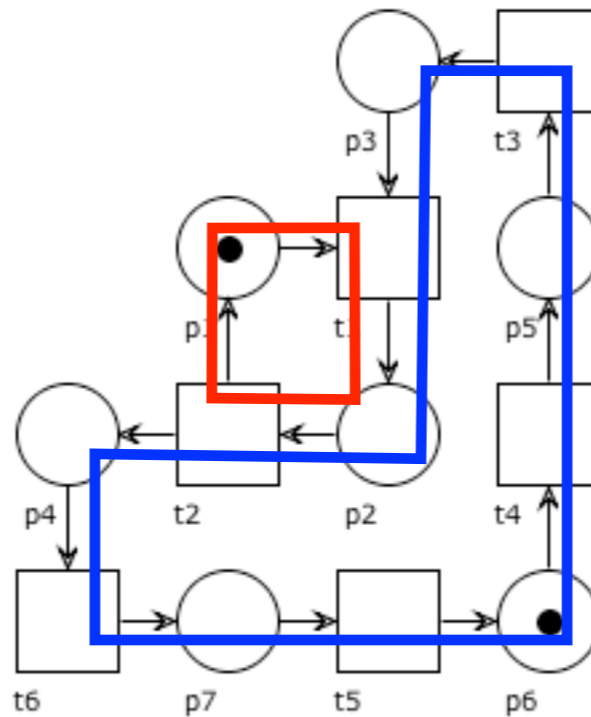Which of the T-systems below is live? (why?)

Question time
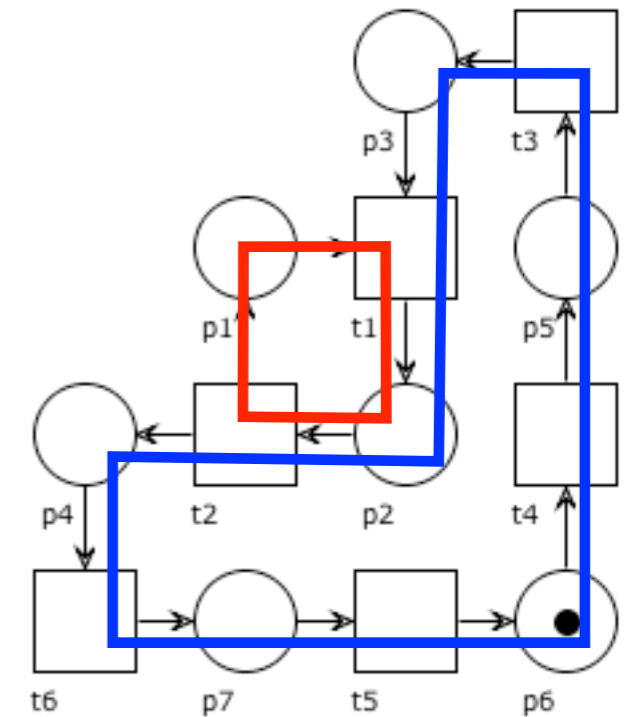
Which of the T-systems below is live? (why?)

Both circuits are marked
=> live

Both circuits are marked
=> live

red circuit is not marked
=> not live

T-systems: recap

T-system + γ circuit + M reachable   => M(γ) = M$_0$(γ)

T-system + γ circuit + M(γ)≠M$_0$(γ) => M not reachable

T-system + γ_1... γ_n circuits: $\exists$i. p$\in$ γ_i <=> p bounded

T-system: M$_0$(γ)>0 for all circuits γ <=> live

T-system: strongly connected   => bounded

T-system + live: strongly connected <=> bounded

T-system: T-invariant **J**   <=> **J** = [k k ... k]

Consequences on workflow nets

Theorem: If N is a workflow net s.t. N* is a T-system then
N is safe and sound **iff**
every circuit of N* is marked

N workflow net => N* strongly connected
N* strongly connected + T-system => N* bounded
N* strongly connected + T-system + $M_0(P)=1$ => N* safe
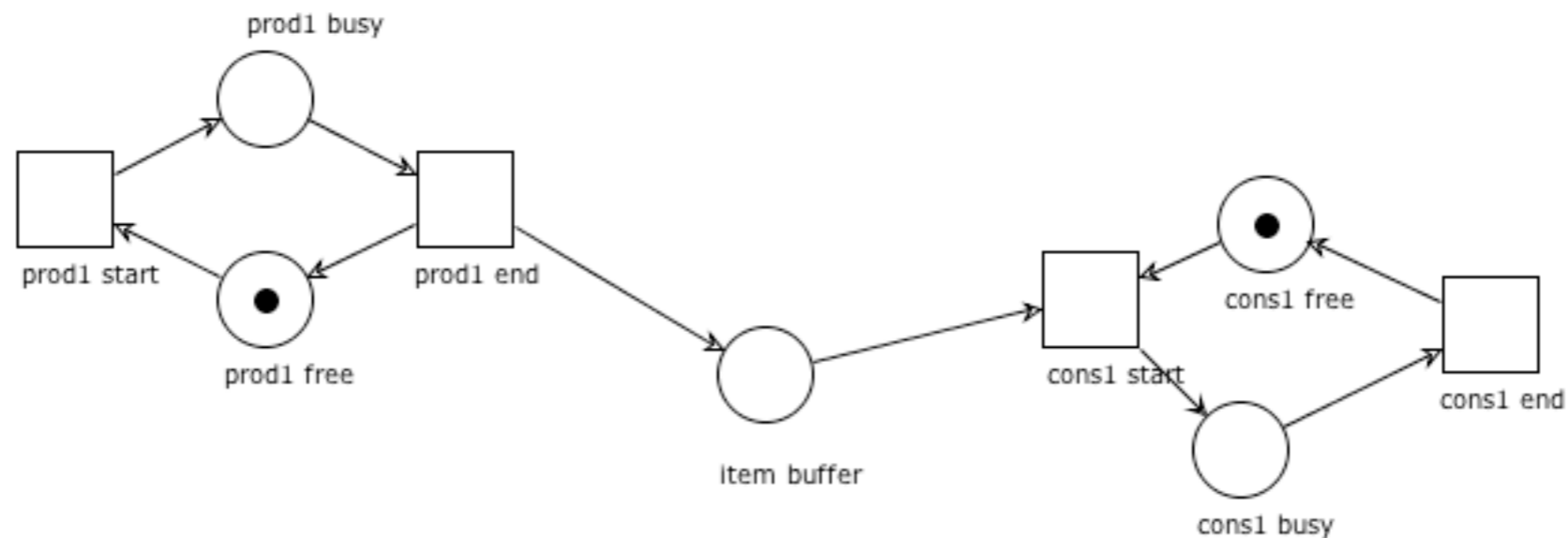
every circuits of N* is marked <=> N* live

Exercises

Which are the circuits of the T-system below?
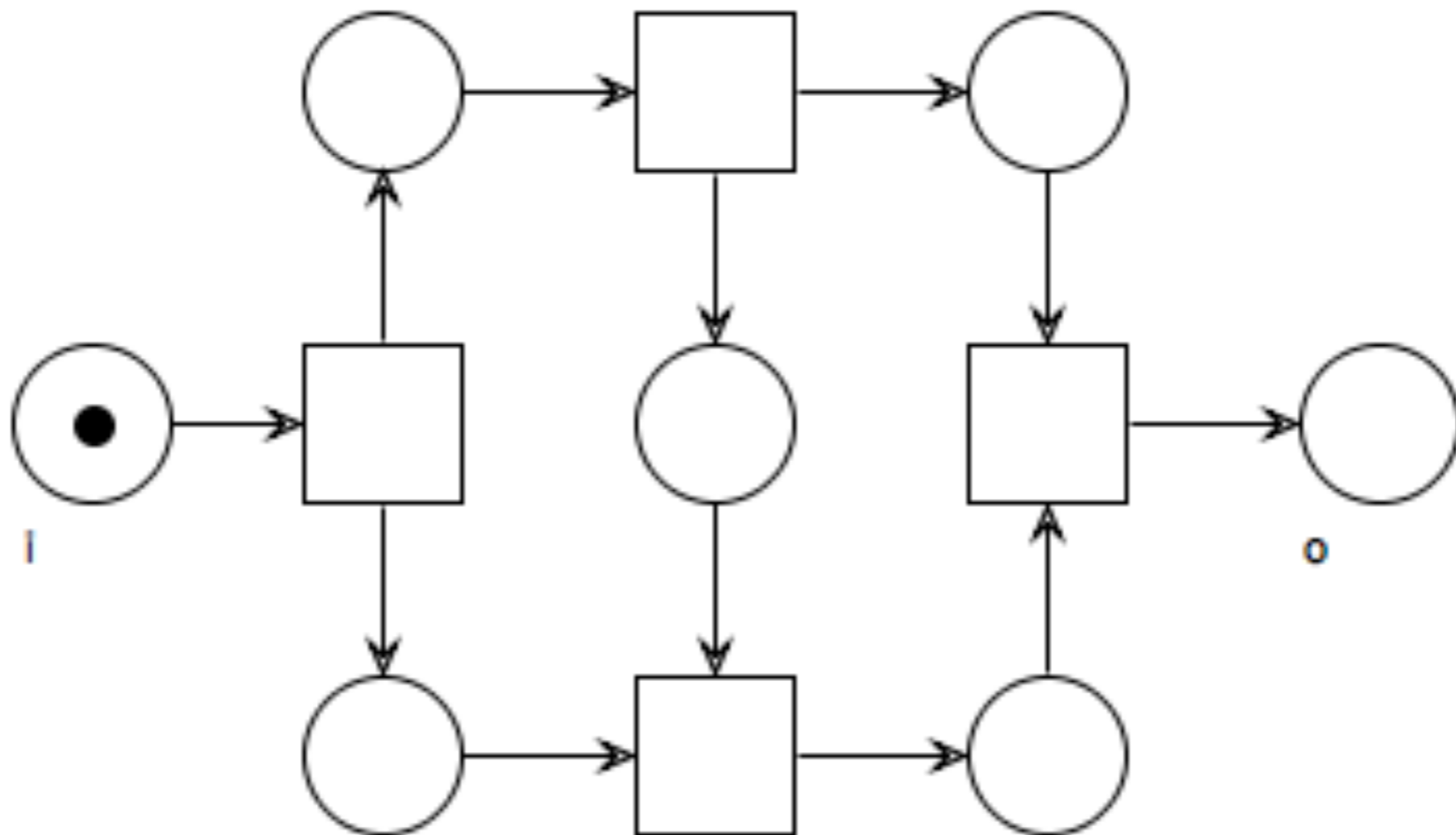Is the T-system below live? (why?)
Which places are bounded? (why?)
Assign a bound to each bounded place.

Exercise

Is the workflow net N below a T-net?
Is it sound?

Exercise

Is the workflow net below a T-net?
Is it sound?