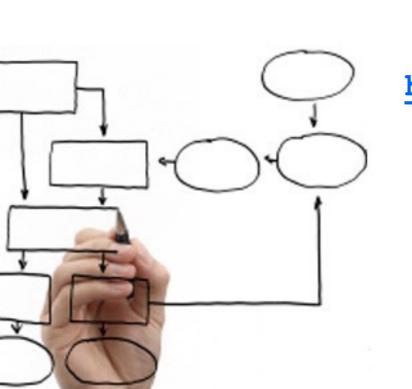
Business Processes Modelling MPB (6 cfu, 295AA)



Roberto Bruni

http://www.di.unipi.it/~bruni

* - P and NP problems

Problems and instances

A problem defines a family of related questions

For example, the factorization problem is: "given a number n, return all its prime factors"

A problem instance is one such question

An instance of the factorization problem is: "return all prime factors of 18"

Decision problem

A decision problem requires just a boolean answer

For example: "given a number n, is n prime?"

And an instance: "is 18 prime?"

Computational Complexity Theory

Computational complexity theory deals with the resources needed to solve a problem

For example, how many steps (time) or memory (space) it takes to solve a problem

P

The complexity class **P** is the set of decision problems that can be solved by a deterministic algorithm in a **P**olynomial number of steps (time) w.r.t. input size

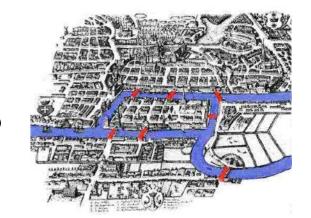
Problems in P can be (checked and) solved effectively

Eulerian circuit (P)

Given a graph G, is it possible to draw an Eulerian circuit over it? (i.e. a circuit that traverses each edge exactly once)

We have seen that it is the same problem as:

Given a graph G, is the degree of each vertex even?



The problem can be (checked and) solved effectively!

NP

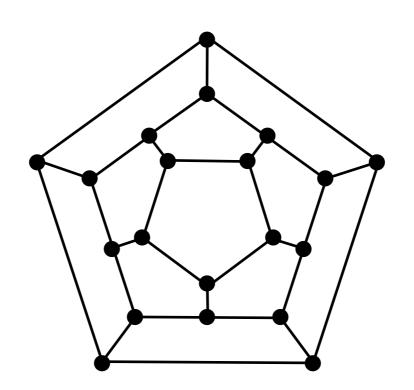
The complexity class **NP** is the set of decision problems that can be **solved** by a **N**on-deterministic algorithm in a **P**olynomial number of steps (time)

Equivalently **NP** is the set of decision problems whose solutions can be **checked** by a deterministic algorithm in a polynomial number of steps (time)

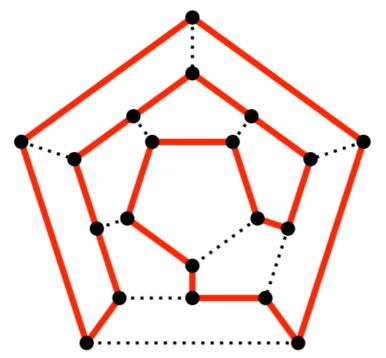
Solutions of problems in NP can be checked effectively

Hamiltonian circuit

Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)

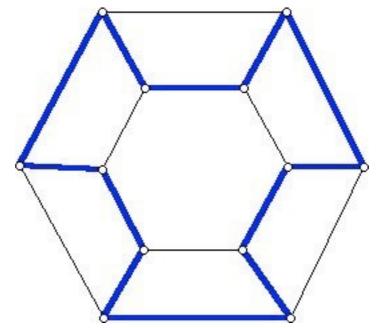


Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)



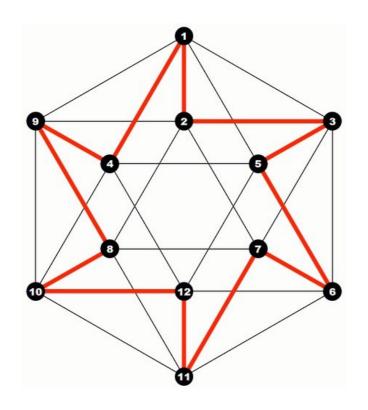
The problem can be checked effectively!

Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)



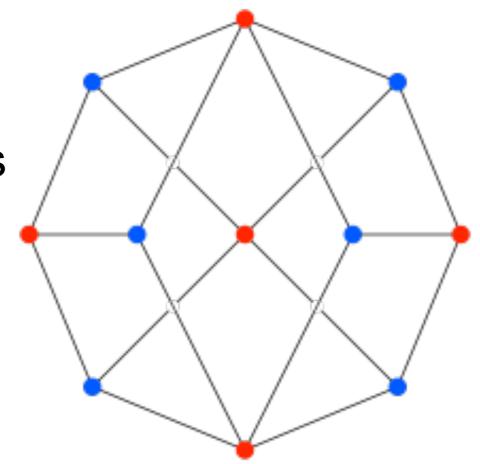
The problem can be checked effectively!

Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)

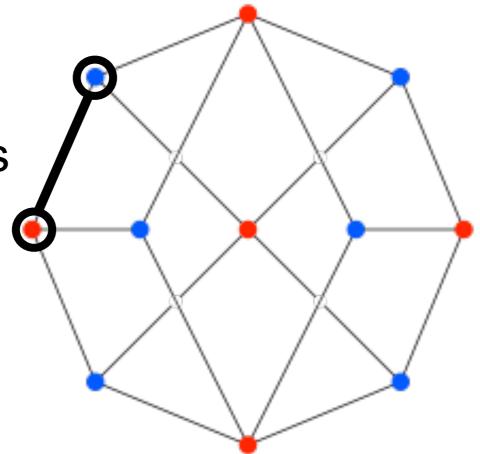


The problem can be checked effectively!

Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)

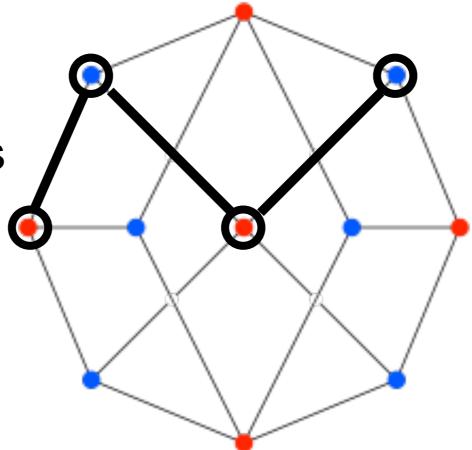


Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)



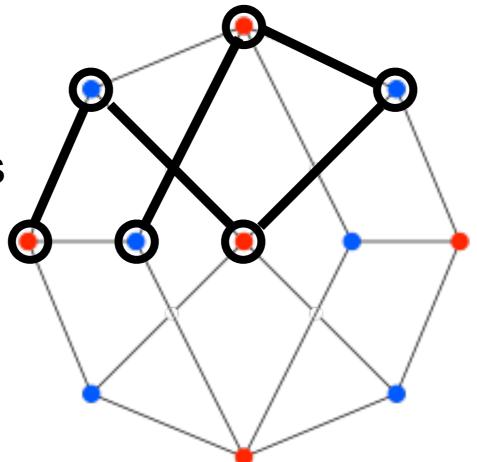
Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)

Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)

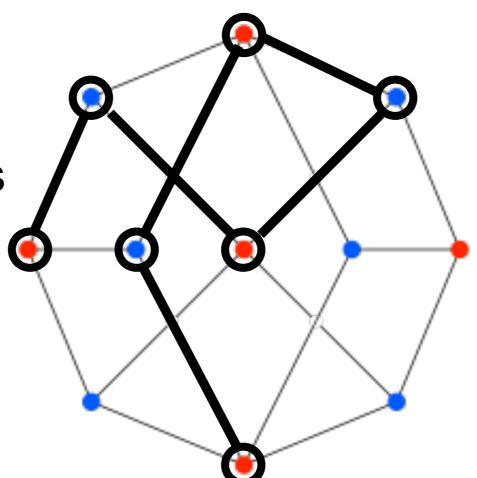


Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)

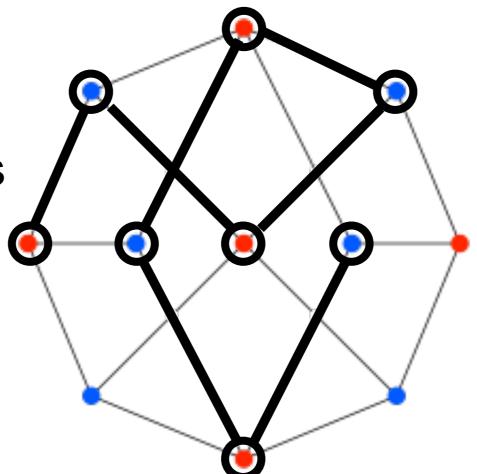
Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)



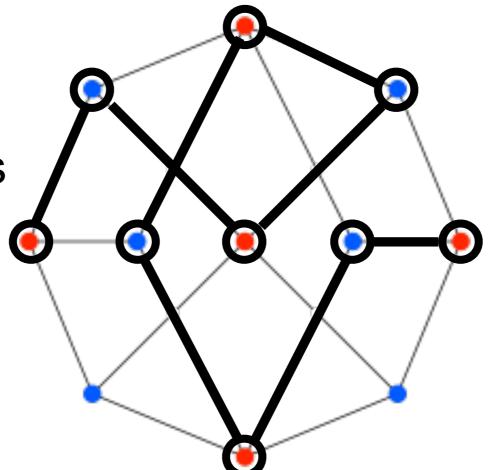
Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)



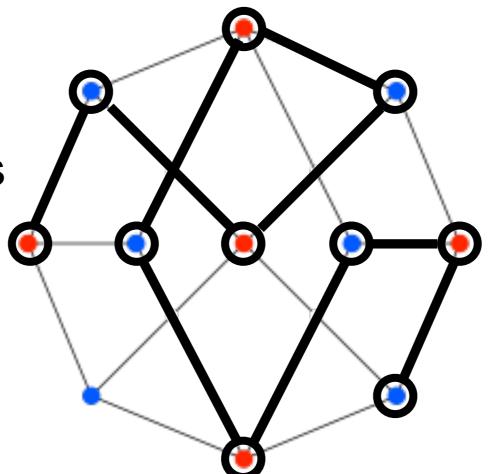
Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)



Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)



Given a graph G, is it possible to draw an Hamiltonian circuit over it? (i.e. a circuit that visits each vertex exactly once)



P vs NP

The question of whether **P** is the same set as **NP** is the most important open question in computer science

Intuitively, it is much harder to solve a problem than to check the correctness of a solution

A fact supported by our daily experience, which leads us to conjecture **P** ≠ **NP**

What if "solving" is not really harder than "checking"? what if **P** = **NP**?

NP-completeness

A problem Q in **NP** is **NP-complete** if every other problem in **NP** can be reduced to Q (in polynomial time)

(finding an effective way to solve such a problem Q would allow to solve effectively any other problem in **NP**)

