

# Methods for the specification and verification of business processes

MPB (6 cfu, 295AA)

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24 - Quantitative Analysis



# Object



We overview some techniques for the quantitative analysis of business processes

Ch.7 of Fundamental of Business Process Management. M. Dumas et al.  
(inspired by slides available at <https://courses.cs.ut.ee/2014/bpm/>)

# Performance Analysis

## **Validation**

is concerned with the relation between the model and reality

## **Verification**

is typically used to answer **qualitative** questions

Is there a deadlock possible? It is possible to successfully handle a specific case? Will all cases terminate eventually?

It is possible to execute two tasks in any order?

## **Performance analysis**

is typically used to answer **quantitative** questions

How many cases can be handled in one hour? What is the average flow time? How many extra resources are required?

How many cases are handled within 2 days?

# Performance dimensions

Any company would like to make its processes

faster,

cheaper,

and better.

# Performance dimensions

Any company would like to make its processes

faster, (time)

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Any company would like to make its processes

faster, (time)

cheaper, (finance)

and better.

# Performance dimensions

Any company would like to make its processes

faster, (time)

cheaper, (finance)

and better. (quality)

# KPI

To estimate the performance along any dimension we need to fix something that can be measured

**A process performance measure is**  
a quantity that can be unambiguously determined  
for a given business process

Time, finance, quality  
can be refined to a number of  
**Key Performance Indicators (KPI)**



# 1. Time

## **Cycle time:**

the time needed to handle one case from start to end.

## **Processing time (also service time):**

the time that resources spend on actually handling the case

## **Waiting time:**

the time that a case spends in *idle* mode  
(e.g., it includes **queueing time** due to unavailability of  
resources to handle the case)

# Time objectives

One can aim to reduce the *average* cycle time

or to reduce the *maximal* cycle time

or to *meet a cycle time* negotiated with the customer

## 2. Finance

Cost, turnover, yield or revenue are all concerned with finance-related performance dimensions.

A yield increase may have the same effect as a cost decrease w.r.t. the organization profit.

Business process redesign is typically concerned with cost.

# Aggregation functions

There are several types of cost:  
cost of production, of delivery, of human resources, ...

Each type can be refined into performance measures by  
selecting an **aggregation function** such as:  
count, average, variance, minimum, maximum, ...

Example: average delivery cost per item

# Cost types

## **Fixed cost:**

overhead costs not affected by the intensity of processing  
(e.g., use of infrastructure, maintenance costs).

## **Variable cost:**

positively correlated with some variable quantity  
(e.g. the level of sales, the number of purchased goods,  
the number of new hires)

## **Operational cost:**

closer to productivity,  
often directly related to the output of a business process  
(e.g. labor cost in producing a good or delivering a service)

# Operational cost

Process redesign is often aimed to **reduce operational cost**, particularly labor cost

Although task automation may reduce labor cost, it may cause incidental cost involved with developing the respective application and fixed maintenance cost for it

# 3. Quality

**External quality:** from the viewpoint of the client  
(e.g. client satisfaction with the delivered product or with the way the process has been executed)

Important factors: amount, relevance, quality and timeliness of the information a client receives as process progresses

**Internal quality:** from the viewpoint of process participants  
Important factors: the level of control of the work performed, of variation experienced, of challenges faced

# Quality vs time

External process quality is often measured in terms of time (e.g. the average cycle time or the percentage of cases where deadlines are missed)

In the following we assume that any performance measure where time is involved is classified under the time dimension, even if it is related to quality



# Deriving performance measures

One possible method for deriving performance measures for a given process is the following:

1. Formulate performance objectives of the process at the high level, in the form of a desirable state that the process should ideally reach
2. For each performance objective, identify the relevant performance dimensions and aggregation functions and derive one or more KPI for the objective
3. Define a target objective for each KPI

# Deriving performance measures: example

A restaurant has recently lost many customers due to **poor customer service**.

The management team has decided to address this issue first of all by **focussing on the delivery of meals**.

The team gathered data by asking customers about how quickly they liked to receive their meals and what they considered as an **acceptable wait**.

The data suggested that half of the customers would prefer their meals to be served in **15 minutes or less**.

All customers agreed that a waiting time of **30 minutes or more** is unacceptable.

# Deriving performance measures: example

1. Formulate performance objectives of the process at the high level, in the form of a desirable state that the process should ideally reach (e.g., customers should be served in less than 30 minutes)
2. For each performance objective, identify the relevant performance dimensions and aggregation functions and derive one or more KPI for the objective (e.g., time dimension,  $ST_{30}$  be the percentage of customers served in less than 30 minutes)
3. Define a target objective for each KPI (e.g.,  $ST_{30} \geq 97\%$ )

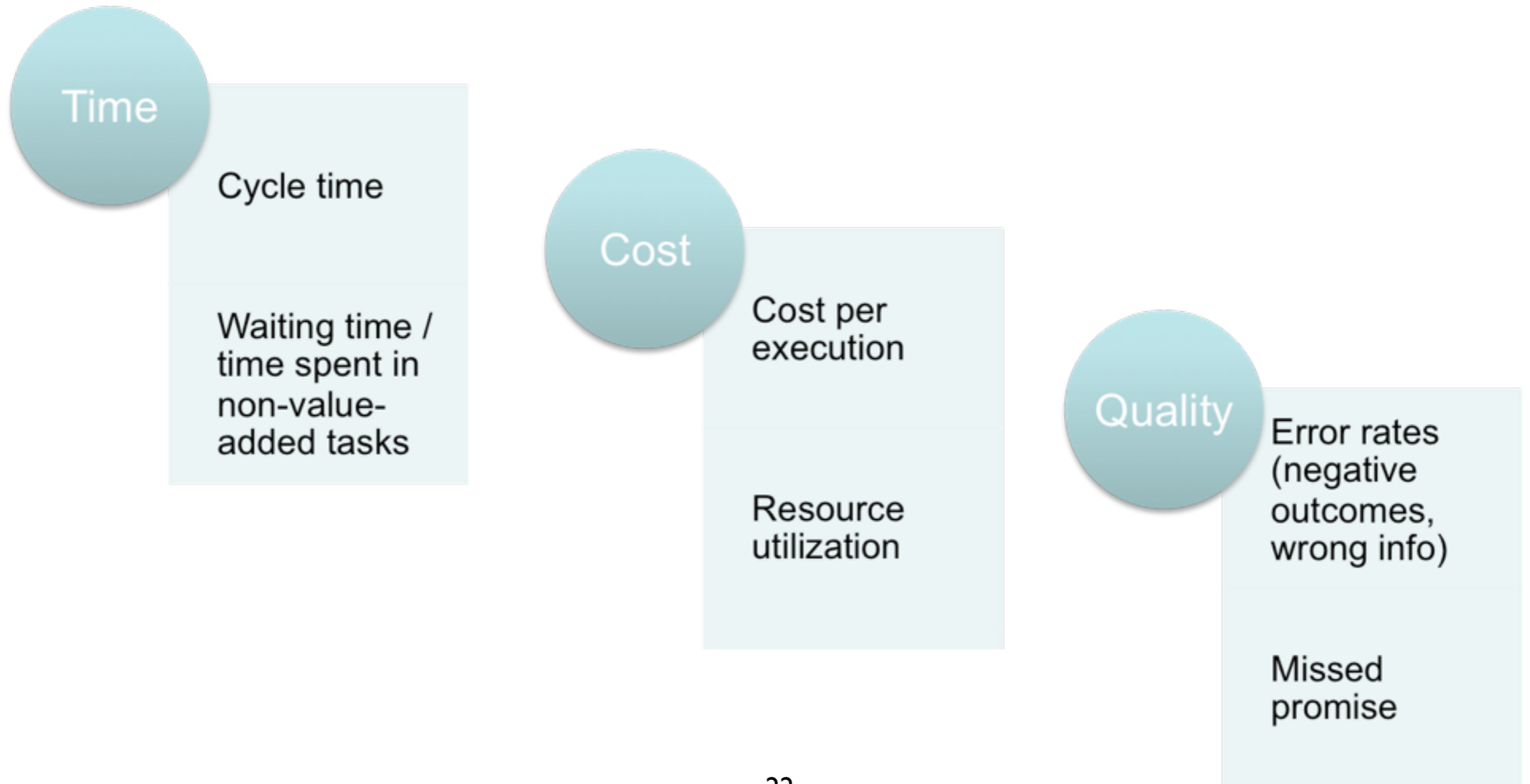
# Deriving performance measures: example

1. Formulate performance objectives of the process at the high level, in the form of a desirable state that the process should ideally reach (e.g., customers should be served in about 15 minutes)
2. For each performance objective, identify the relevant performance dimensions and aggregation functions and derive one or more KPI for the objective (e.g., time dimension,  $ST_{15}$  be the percentage of customers served in less than 15 minutes)
3. Define a target objective for each KPI (e.g.,  $ST_{15} \geq 85\%$ )

# Deriving performance measures: example

1. Formulate performance objectives of the process at the high level, in the form of a desirable state that the process should ideally reach (e.g., customers should be served in about 15 minutes)
2. For each performance objective, identify the relevant performance dimensions and aggregation functions and derive one or more KPI for the objective (e.g., time dimension, AMDT be the average meal delivery time)
3. Define a target objective for each KPI (e.g.,  $AMDT \leq 15'$ )

# Typical process performance measures



# Flow analysis

Flow analysis is a family of techniques to estimate the overall performance of a process given some knowledge about the performance of its activities

Examples:

we calculate the min/max/average **cycle time** of a process given the min/max/average cycle time of each activity

we compute the average **cost** of a process knowing the cost-per-execution of each activity.

we calculate the **error rate** of a process given the error rate of each activity.

# Cycle time analysis



# Cycle time analysis

**Cycle time** = difference between the start time (ready to be executed) and the end time (completion) of a case

**Cycle time analysis** = the task of calculating the average cycle time of an entire process or some process fragment

**Assumption:** average activity times are available for all the activities involved in the process

**Activity time** = waiting time + processing time

# Flow patterns

The simplest case is that of a single activity, but then we can take into account different structure patterns that frequently occur:

paths composed in sequence

alternative paths (XOR split and join)

parallel paths (AND split and join)

rework (1-or-more cycles, 0-or-more cycles)

# Notation

We denote the **average cycle time** by **CT**  
and call it simply *cycle time*

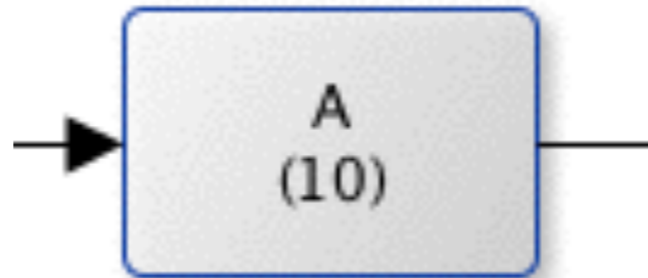
When several (sub)processes  $P_1, P_2, \dots, P_n$  are involved  
we refer to their cycle times by  $CT_1, CT_2, \dots, CT_n$

Similarly, if activities  $A, B, \dots$  are involved  
we refer to their cycle times by  $CT_A, CT_B, \dots$

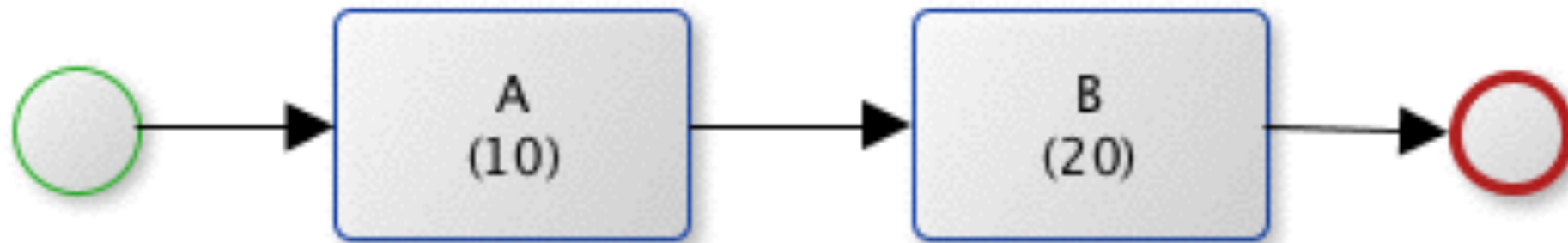
# Notation

In diagrams, we will often write activity cycle time within parentheses

$$CT_A = 10 \text{ units of time}$$

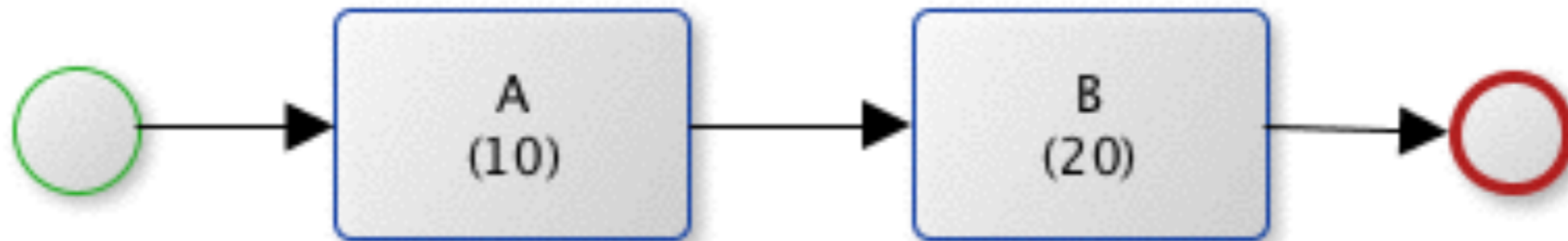


# Sequence



CT = ?

# Sequence



$$CT = CT_A + CT_B = 10 + 20 = 30$$

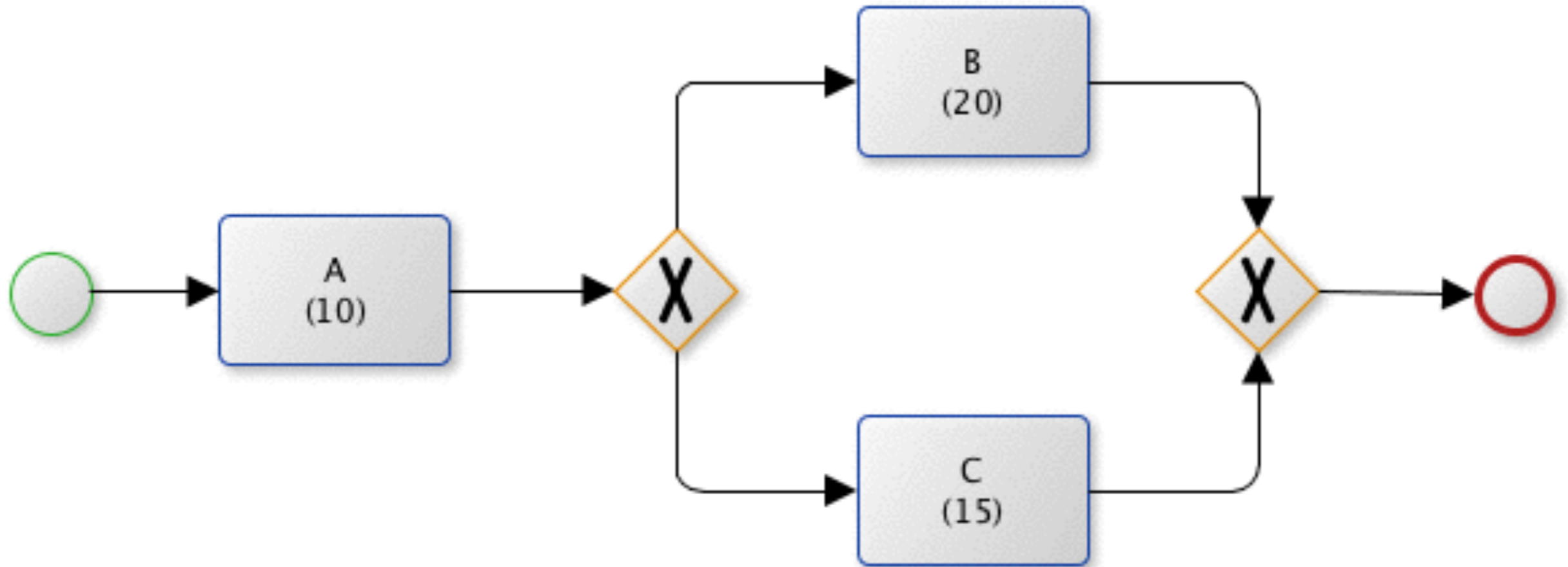
# Sequence

The cycle time of a purely sequential fragment of a process is the sum of the cycle times of the activities in the fragment



$$CT = \sum_{i=1}^n CT_i$$

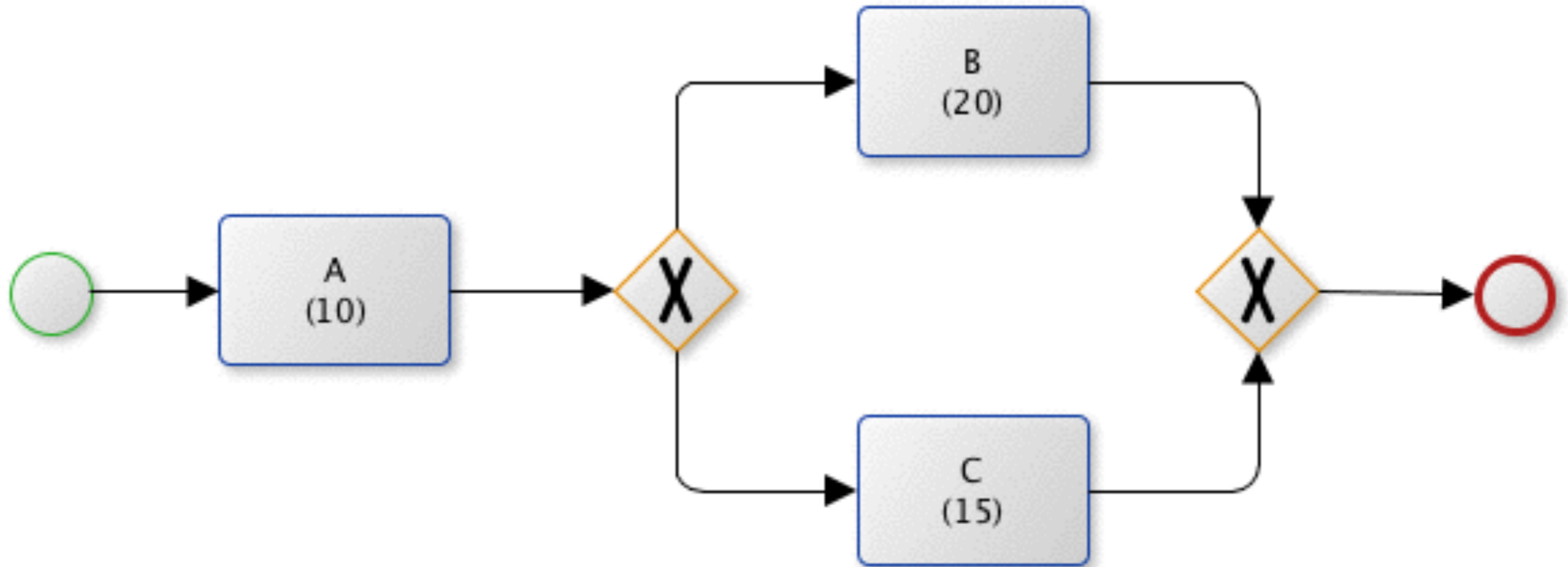
# Alternative paths



CT = ?



# Alternative paths

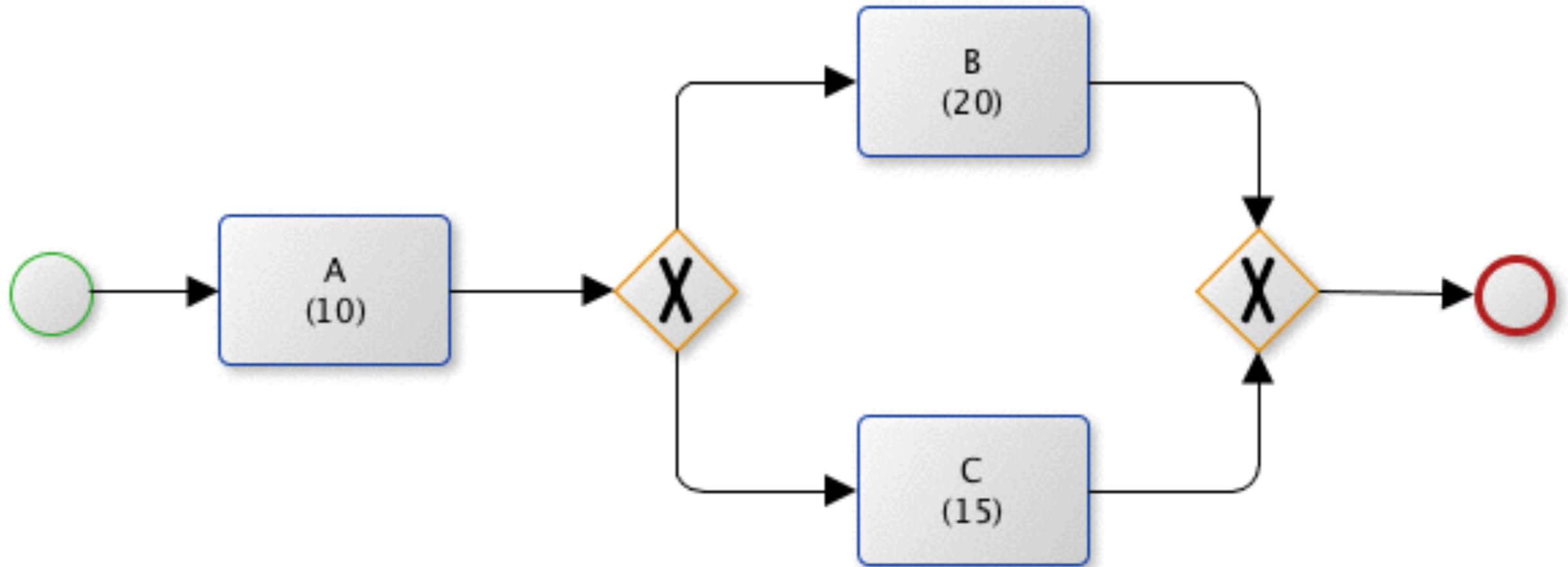


in some cases  $CT = CT_A + CT_B = 10 + 20 = 30$

in other cases  $CT = CT_A + CT_C = 10 + 15 = 25$

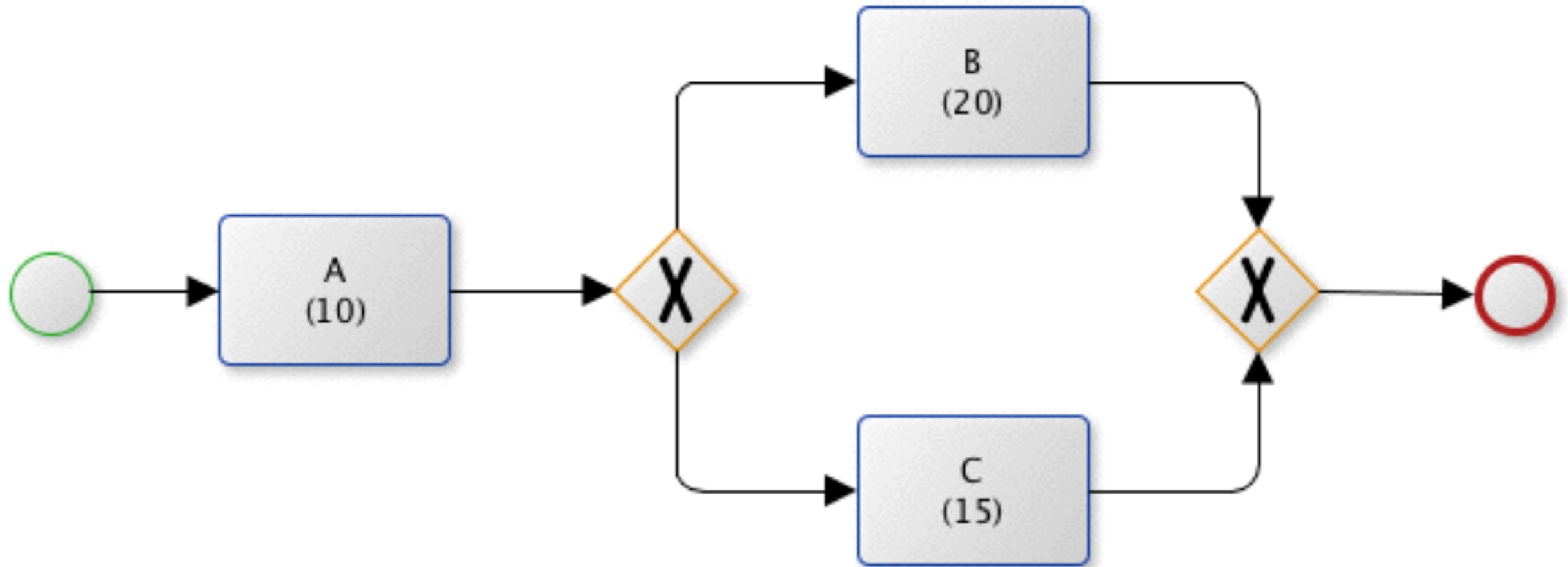
whether the average is closer to 25 or to 30  
depends on **how frequently** each branch is taken

# Alternative paths



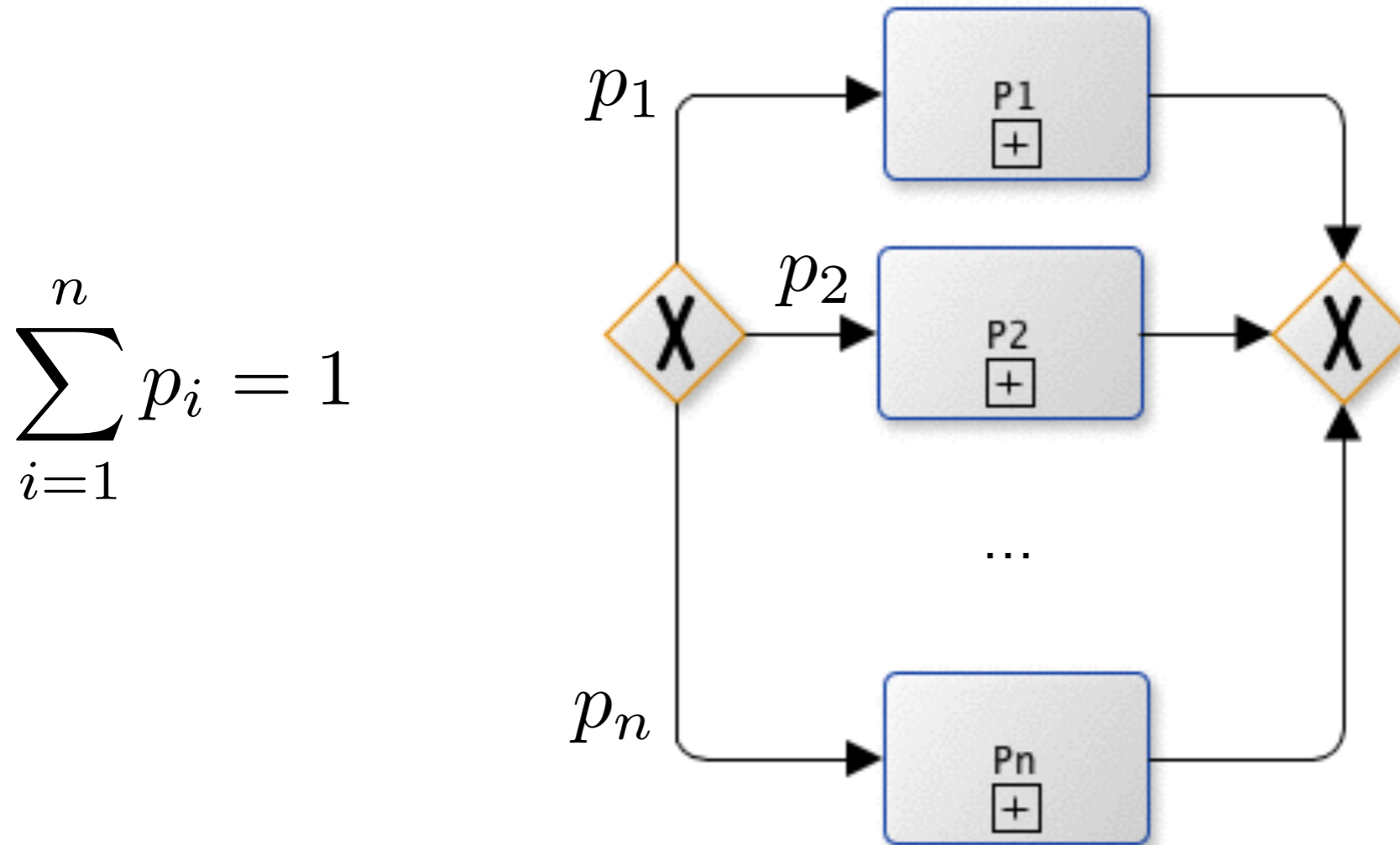
if B and C are taken in 50% of the cases each, then the average sits in the middle between 25 and 30

# Alternative paths



if B is taken in 90% of the cases, then  
the average is closer to 30 than 25

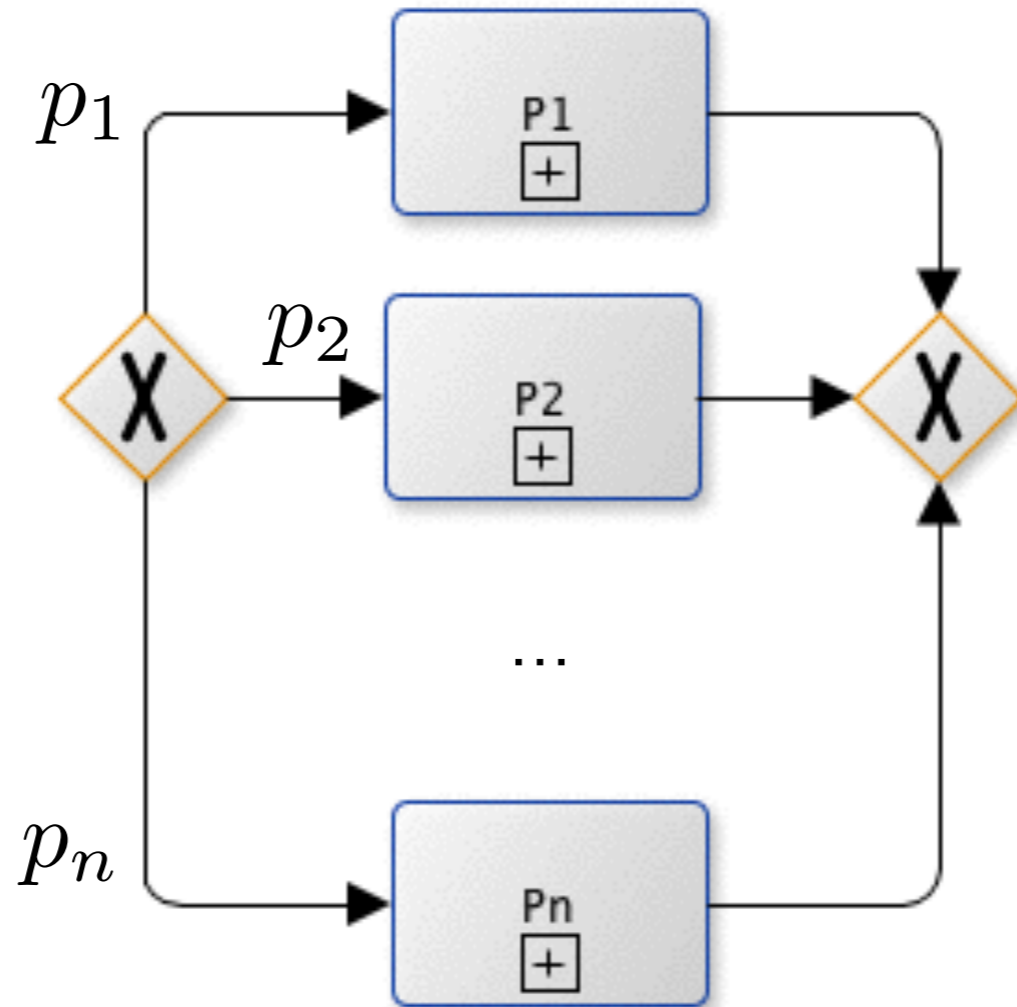
# Branching probability



**Branching probability  $p_i$ :**  
is the frequency with  
which a given branch of a decision gateway is taken

# Alternative paths

$$\sum_{i=1}^n p_i = 1$$

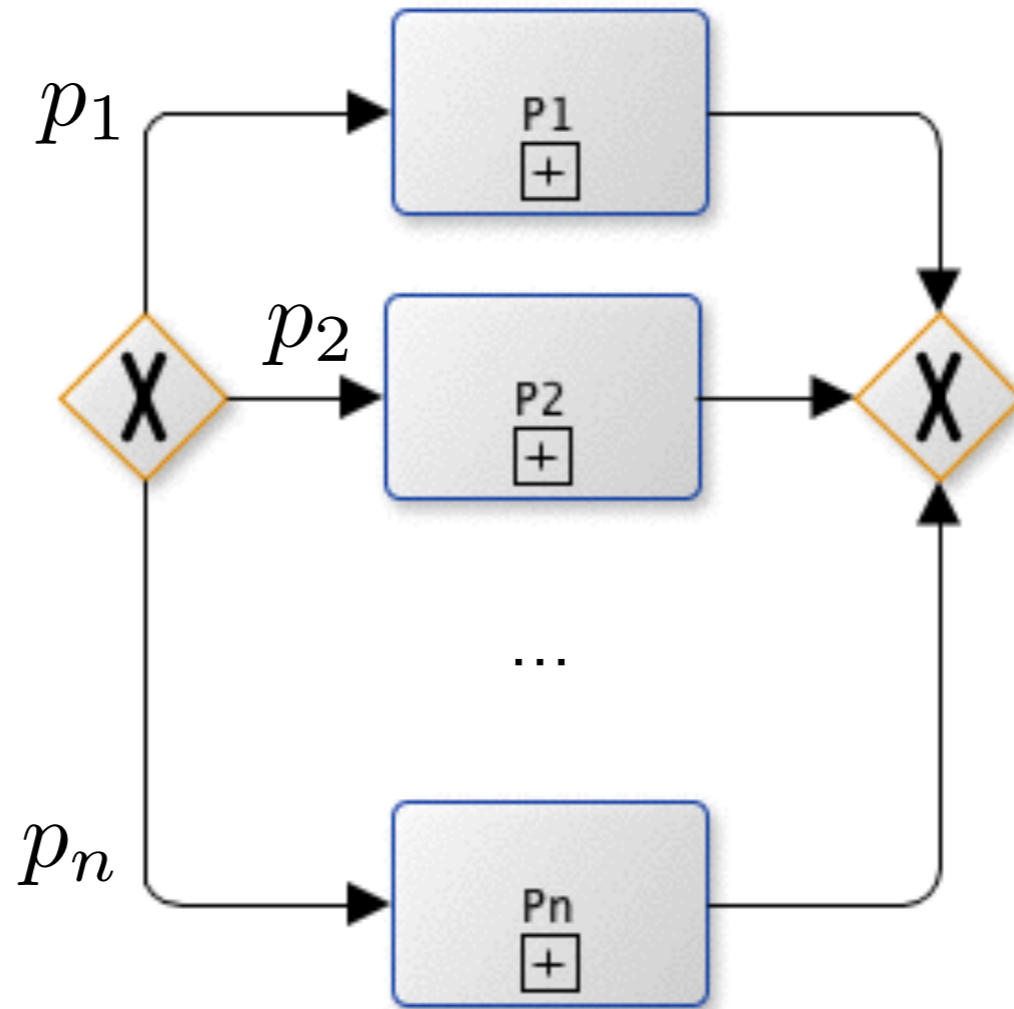


The fragment between the XOR split and join is called **XOR-block**

The cycle time of a XOR-block fragment is the **weighted average** of the cycle times of the branches

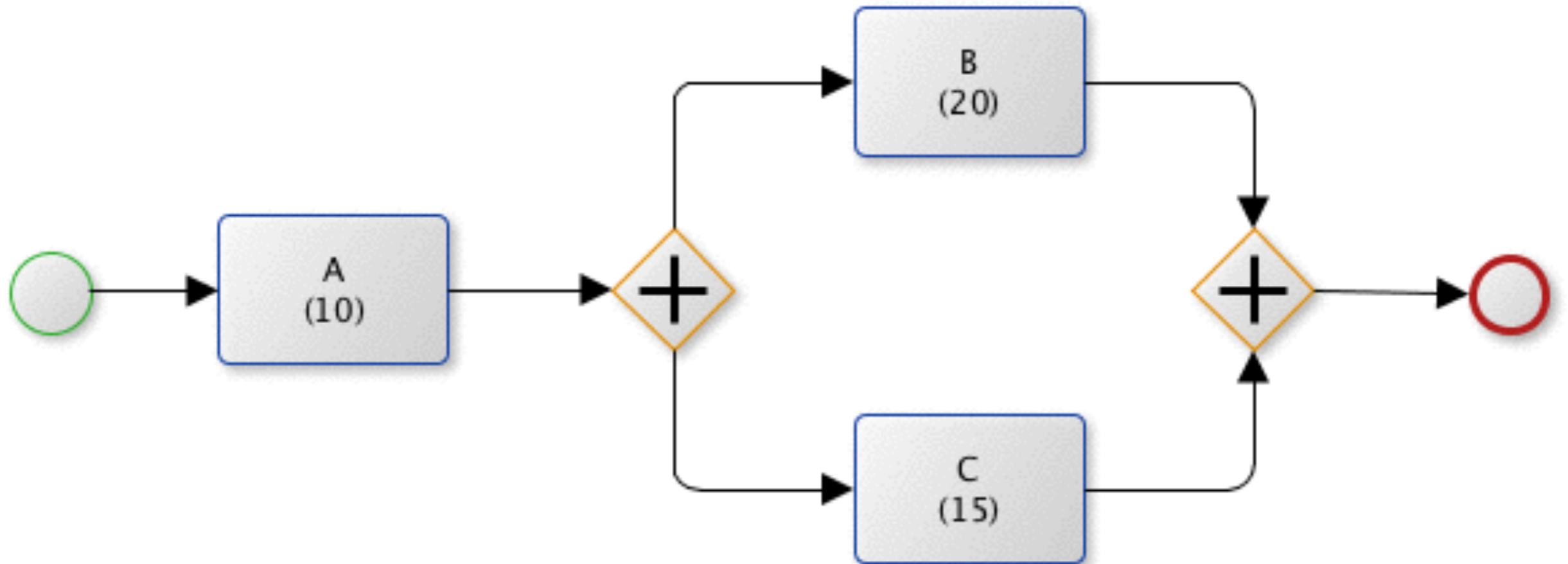
# Alternative paths

$$\sum_{i=1}^n p_i = 1$$



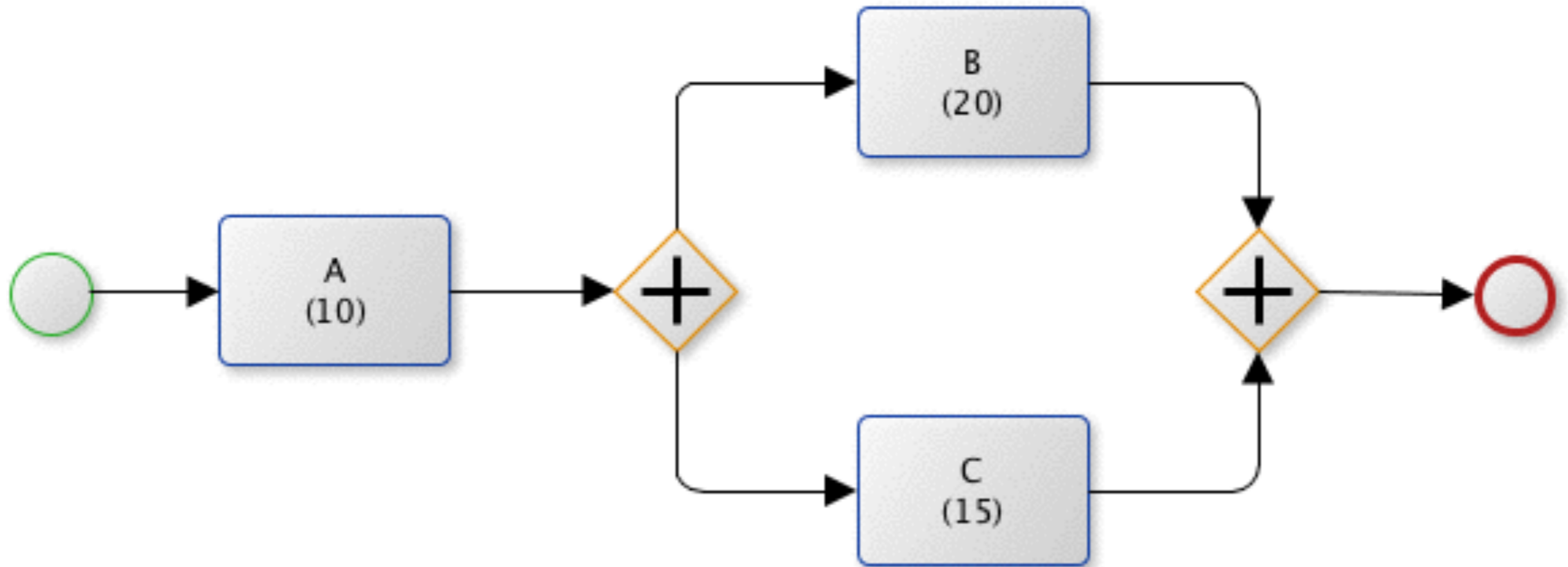
$$CT = \sum_{i=1}^n p_i \cdot CT_i$$

# Parallel paths



$$CT = CT_A + CT_B + CT_C = 10 + 20 + 15 = 45 ?$$

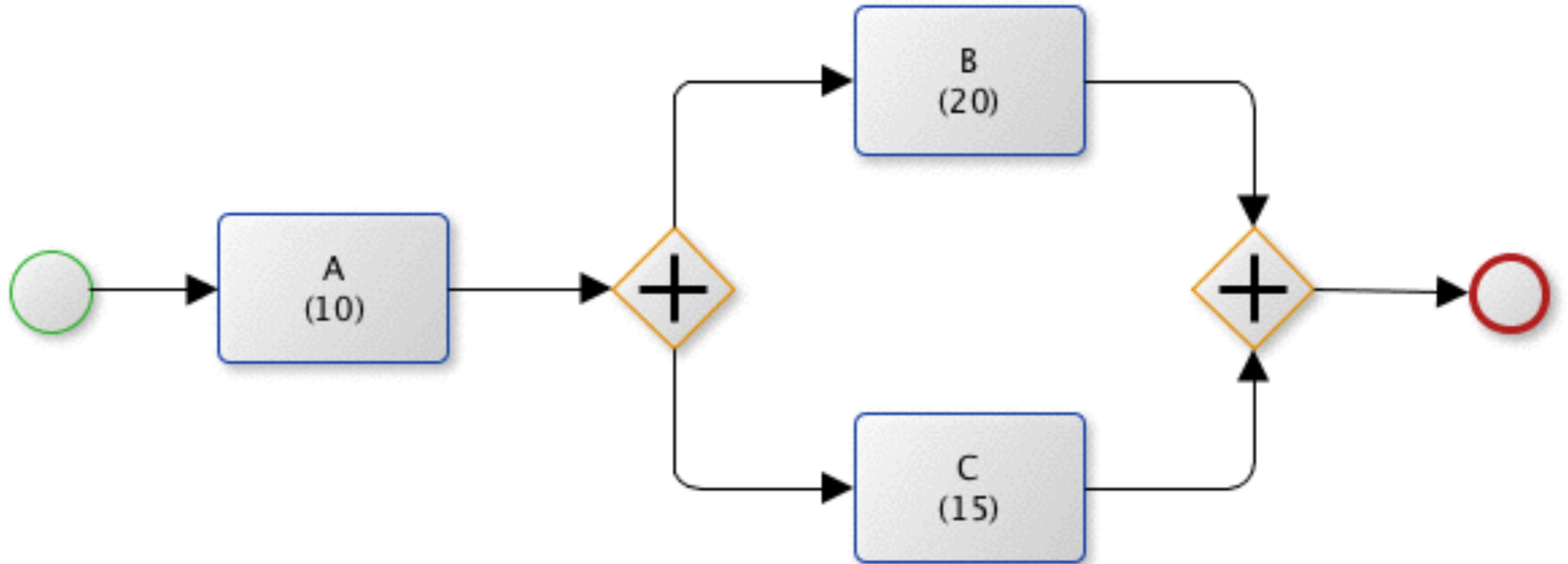
# Parallel paths



$CT = CT_A + CT_B + CT_C = 10 + 20 + 15 = 45 ?$   
(but while B is executed, also C is executed,  
and B takes longer than C)

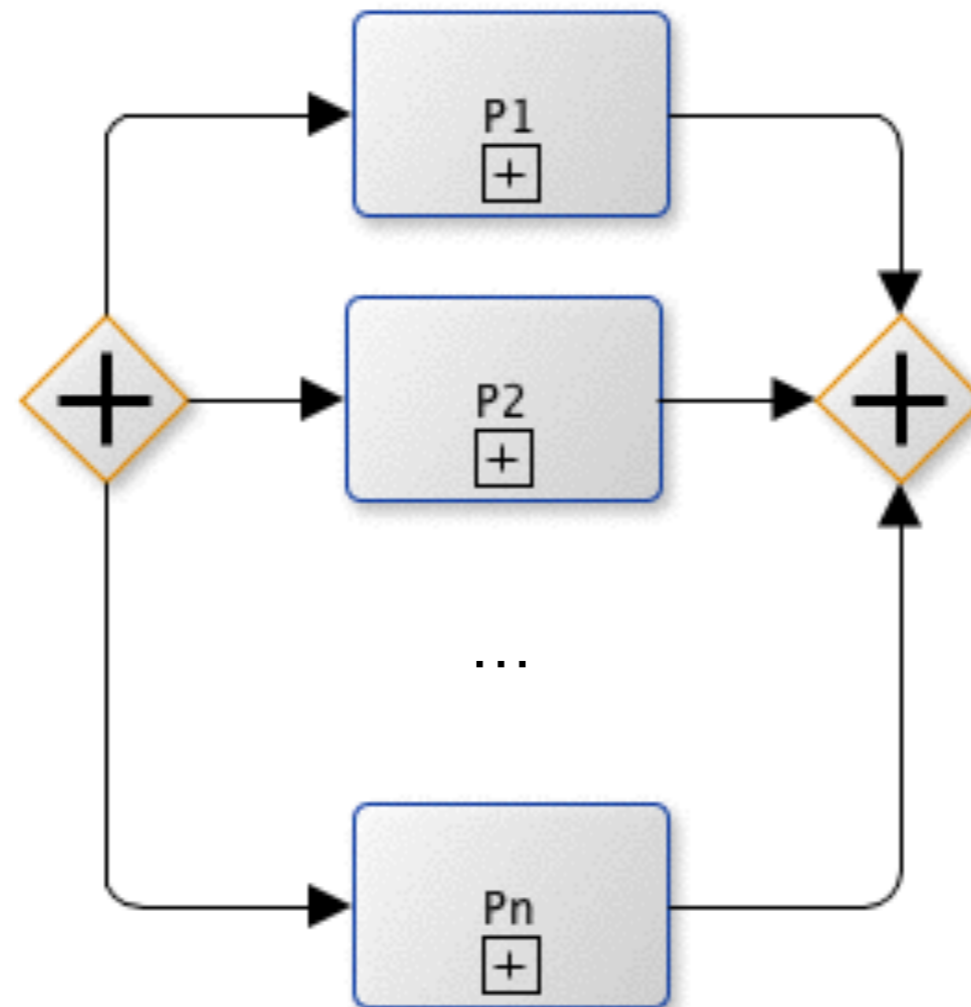


# Parallel paths



$$CT = 10 + \max \{20, 15\} = 10 + 20 = 30$$

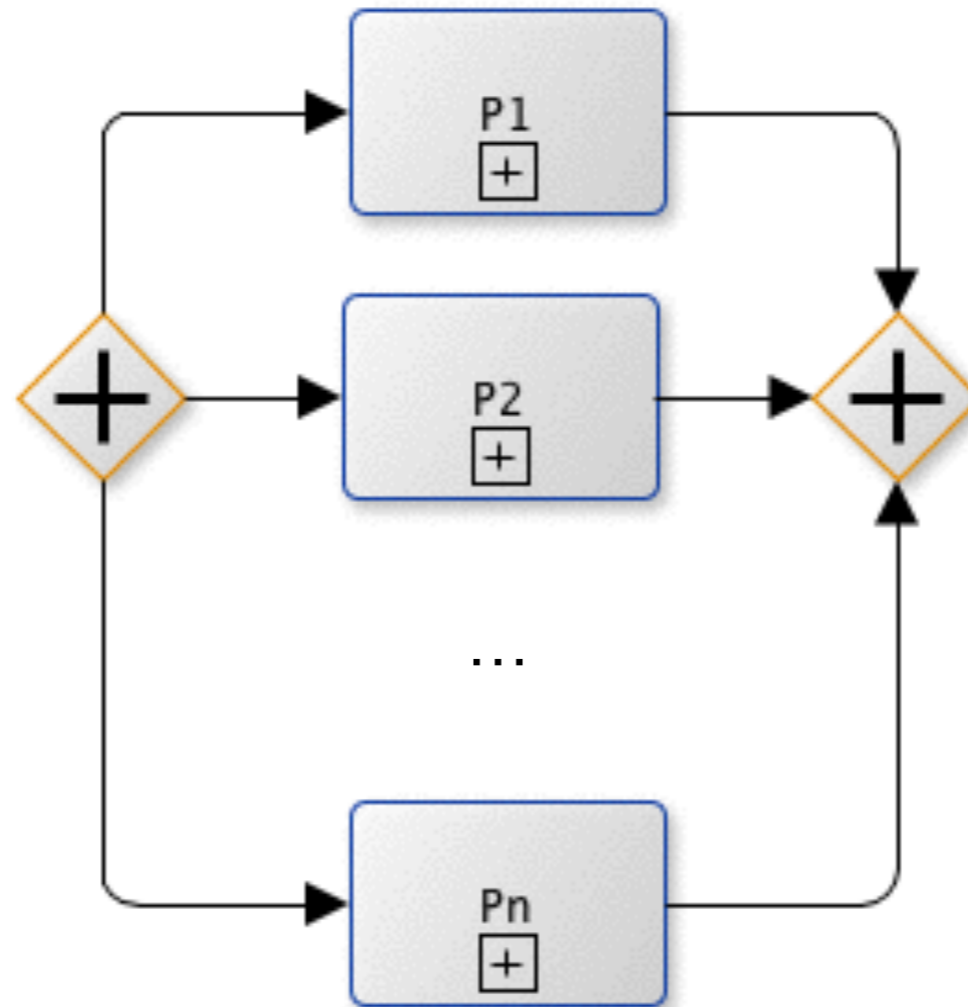
# Parallel paths



The fragment between the AND split and join is called **AND-block**

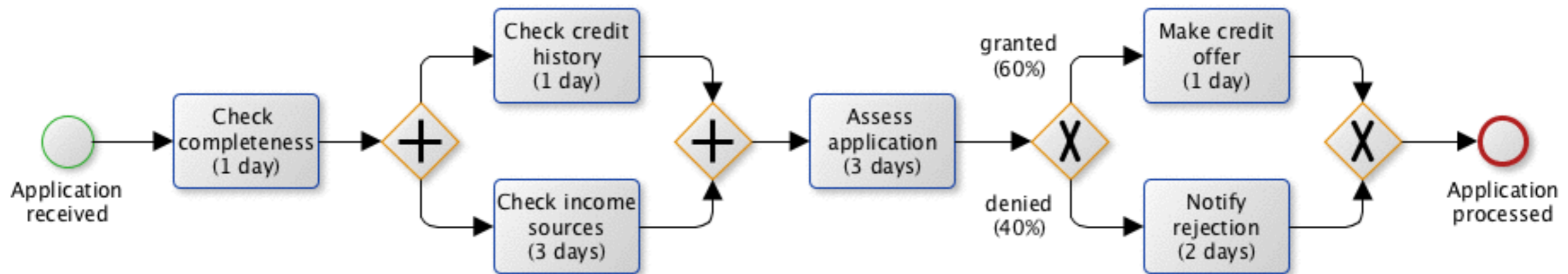
The cycle time of an AND-block fragment is the cycle time of the **slowest** branch

# Parallel paths



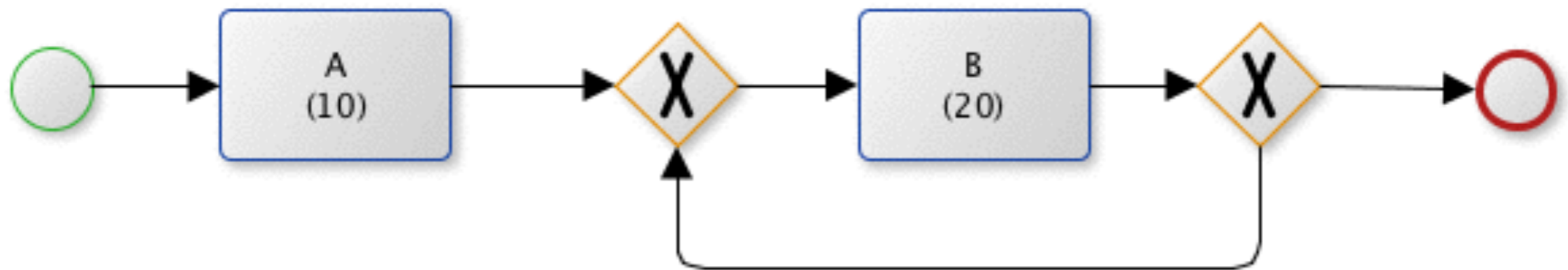
$$CT = \max_i \{CT_i\} = \max\{CT_1, CT_2, \dots, CT_n\}$$

# Question time



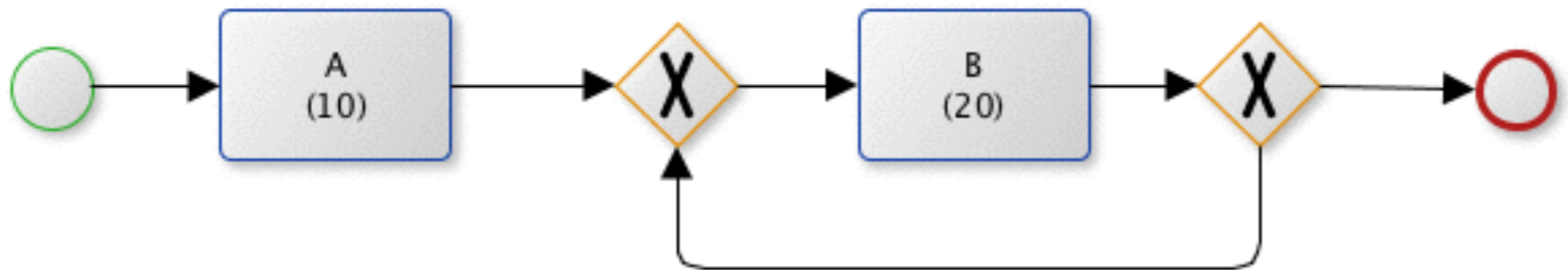
CT = ?

# Rework loop (1 or more times)



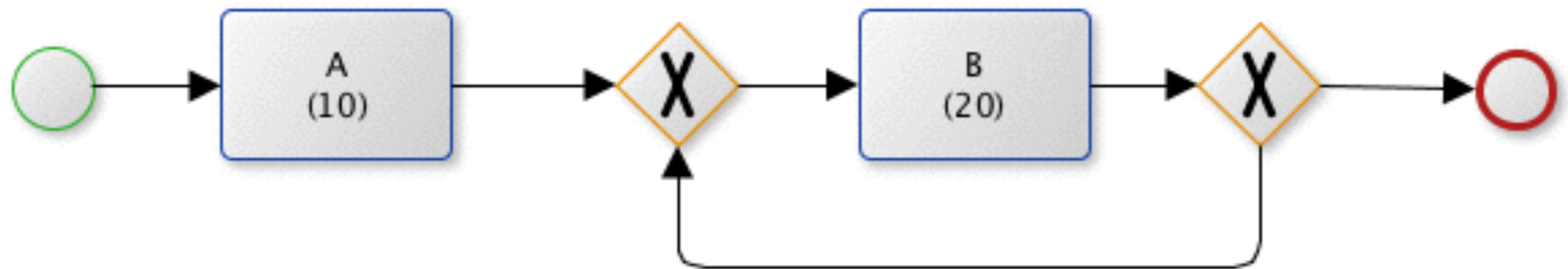
CT = ?

# Rework loop (1 or more times)



$$CT = 10 + 20 + 20 + 20 + \dots ?$$

# Rework loop (1 or more times)



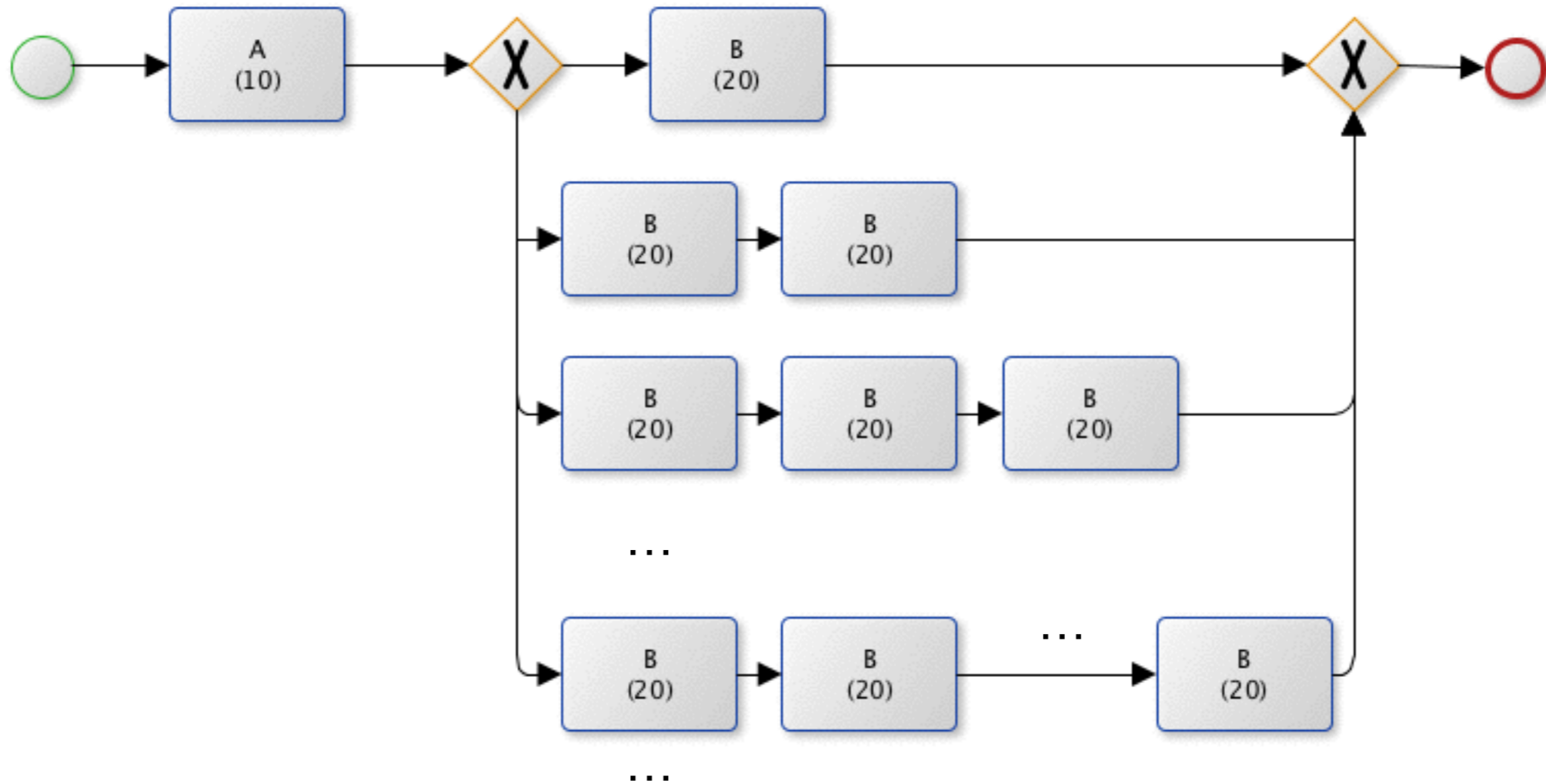
For sure we can say that B will be executed once.

$$(CT \geq 10 + 20 = 30)$$

Then, depending on a choice, B can be executed twice.

Then, a third time, and so on ...

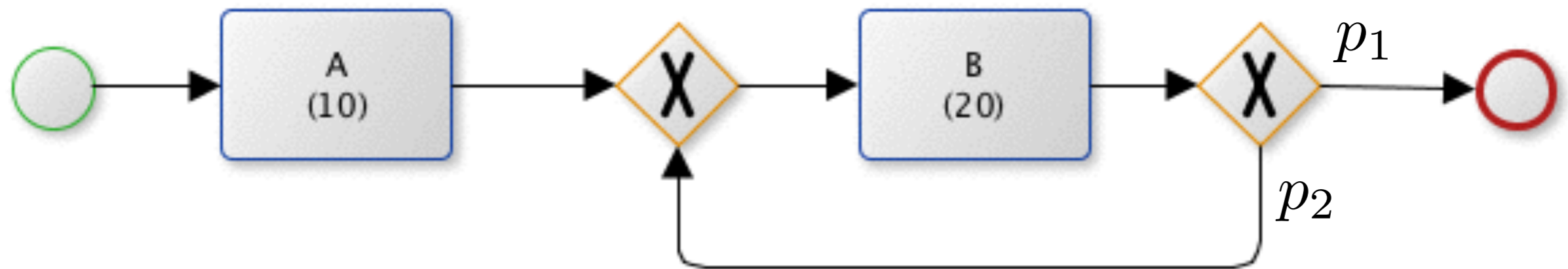
# Rework loop



For sure we can say that B will be executed once.  
Then, depending on a choice, B can be executed twice.  
Then, a third time, and so on ...



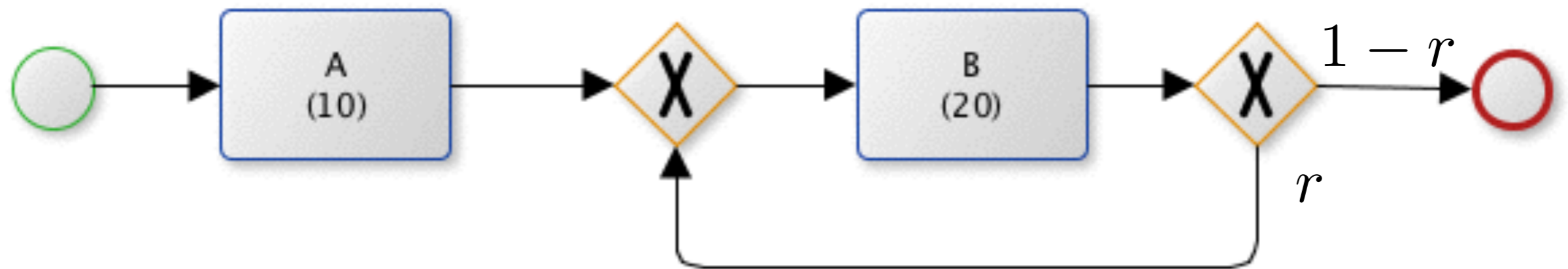
# Branching probability, again...



$$p_1 + p_2 = 1$$

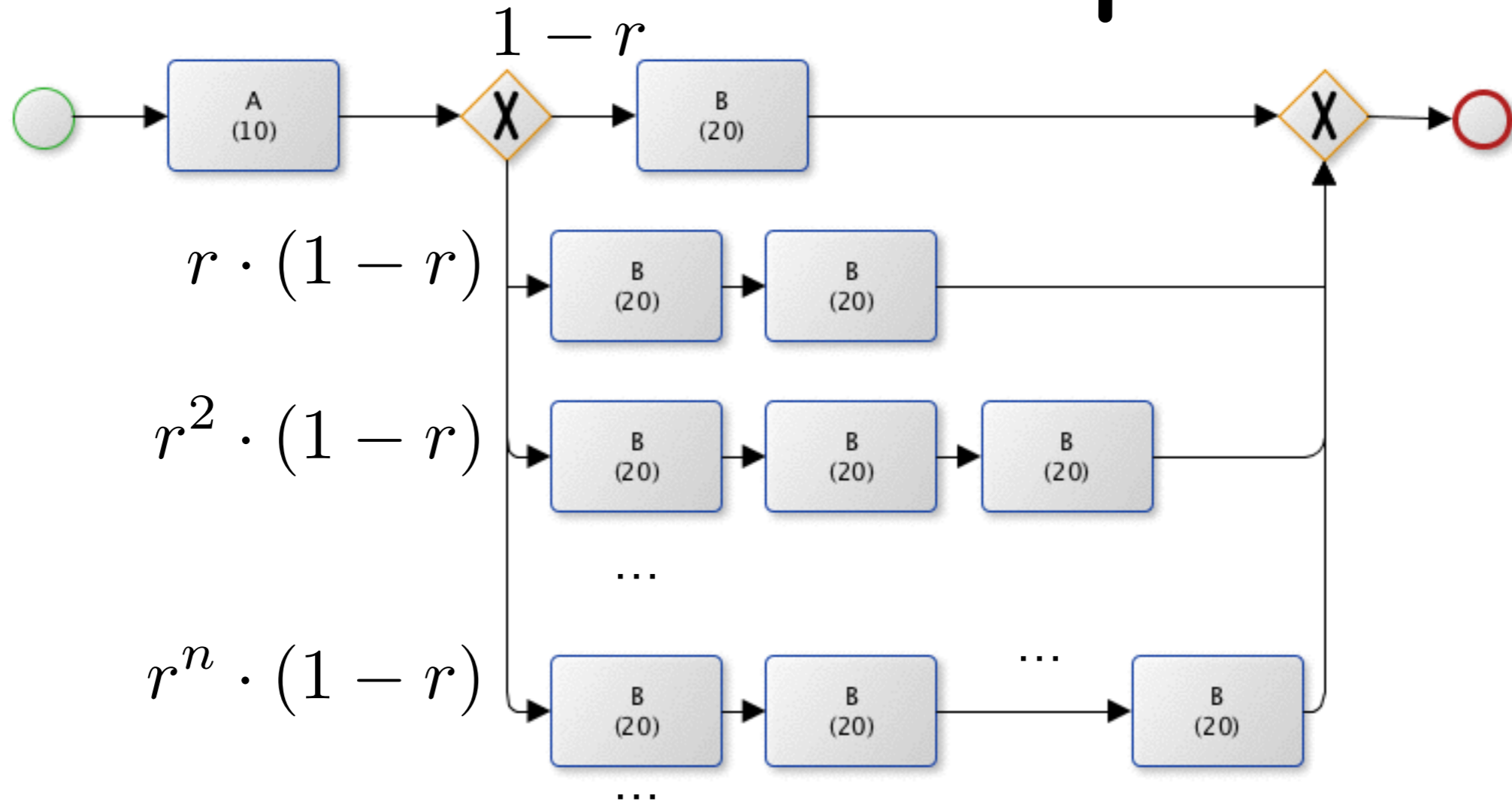
Branching probability  $p_i$ : is the frequency with which a given branch of a decision gateway is taken

# Rework probability



**Rework probability  $r$ :**  
is the frequency with which the task is reworked

# Rework loop

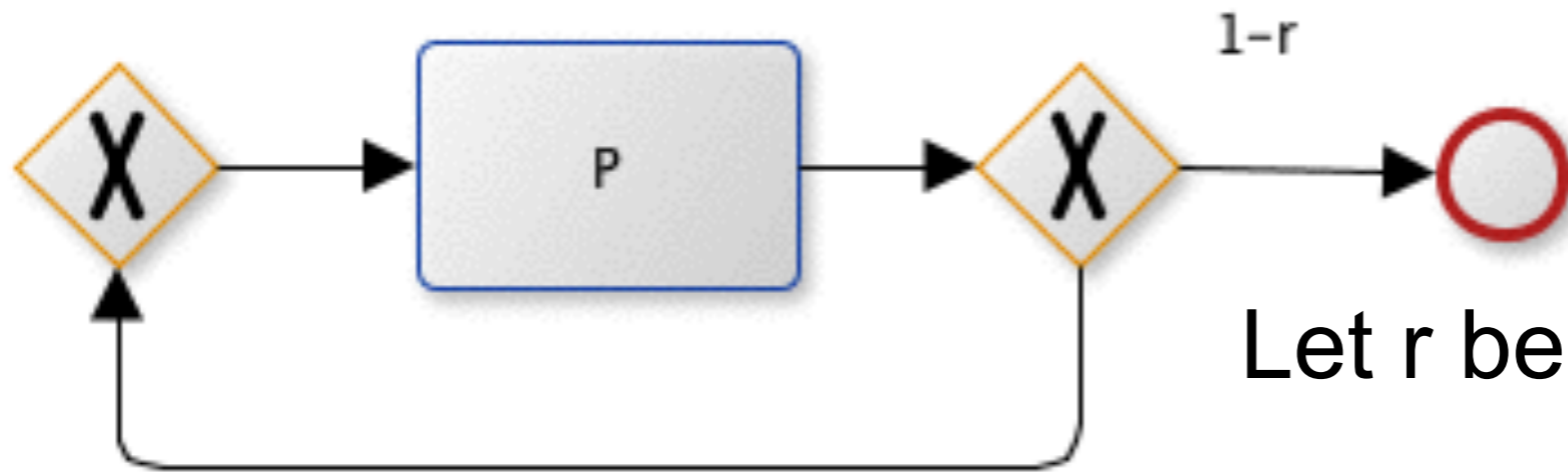


For sure we can say that B will be executed once.  
Then, depending on a choice, B can be executed twice.

Then, a third time, and so on ...

**but always with less and less probability**

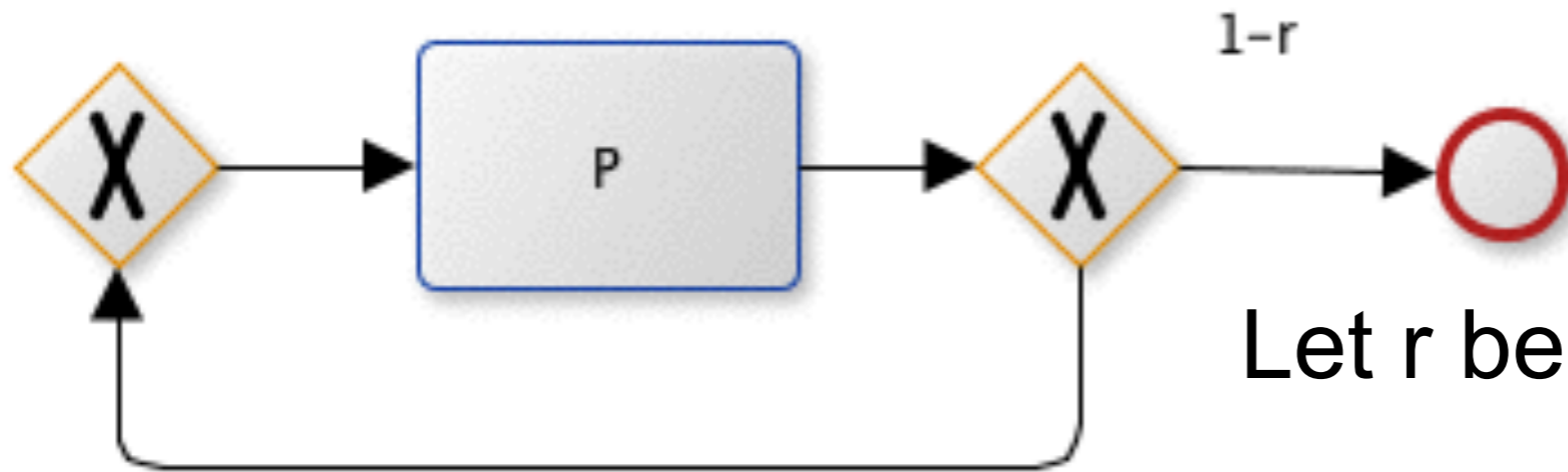
# Rework loop



Let  $r$  be the *rework probability*,  
then:

$$\begin{aligned} CT &= 1 \cdot CT_P \cdot r^0 \cdot (1 - r) \\ &+ 2 \cdot CT_P \cdot r^1 \cdot (1 - r) \\ &+ \dots \\ &+ n \cdot CT_P \cdot r^{n-1} \cdot (1 - r) \\ &+ \dots \\ &= \sum_{i=1}^{\infty} i \cdot CT_P \cdot r^{i-1} \cdot (1 - r) \end{aligned}$$

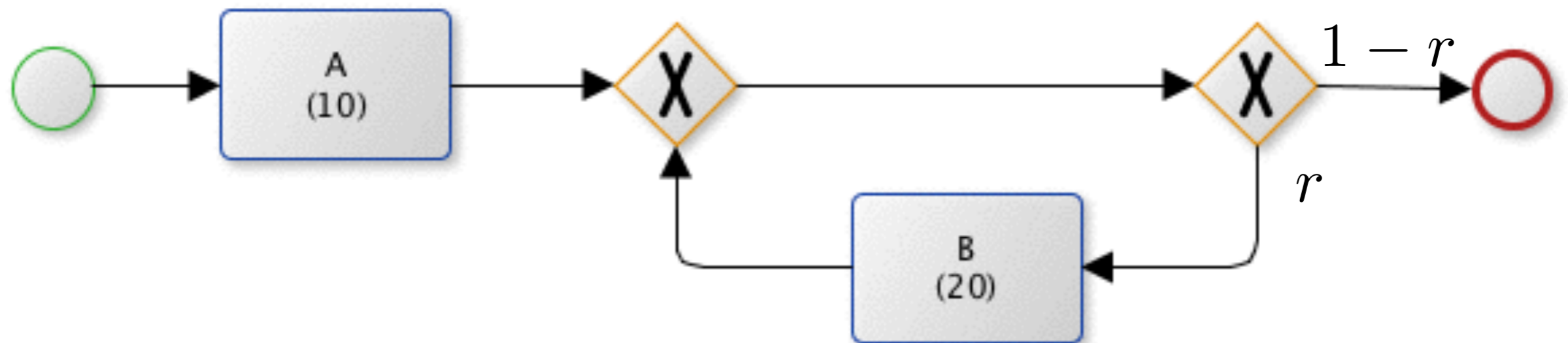
# Rework loop



Let  $r$  be the *rework probability*, then:

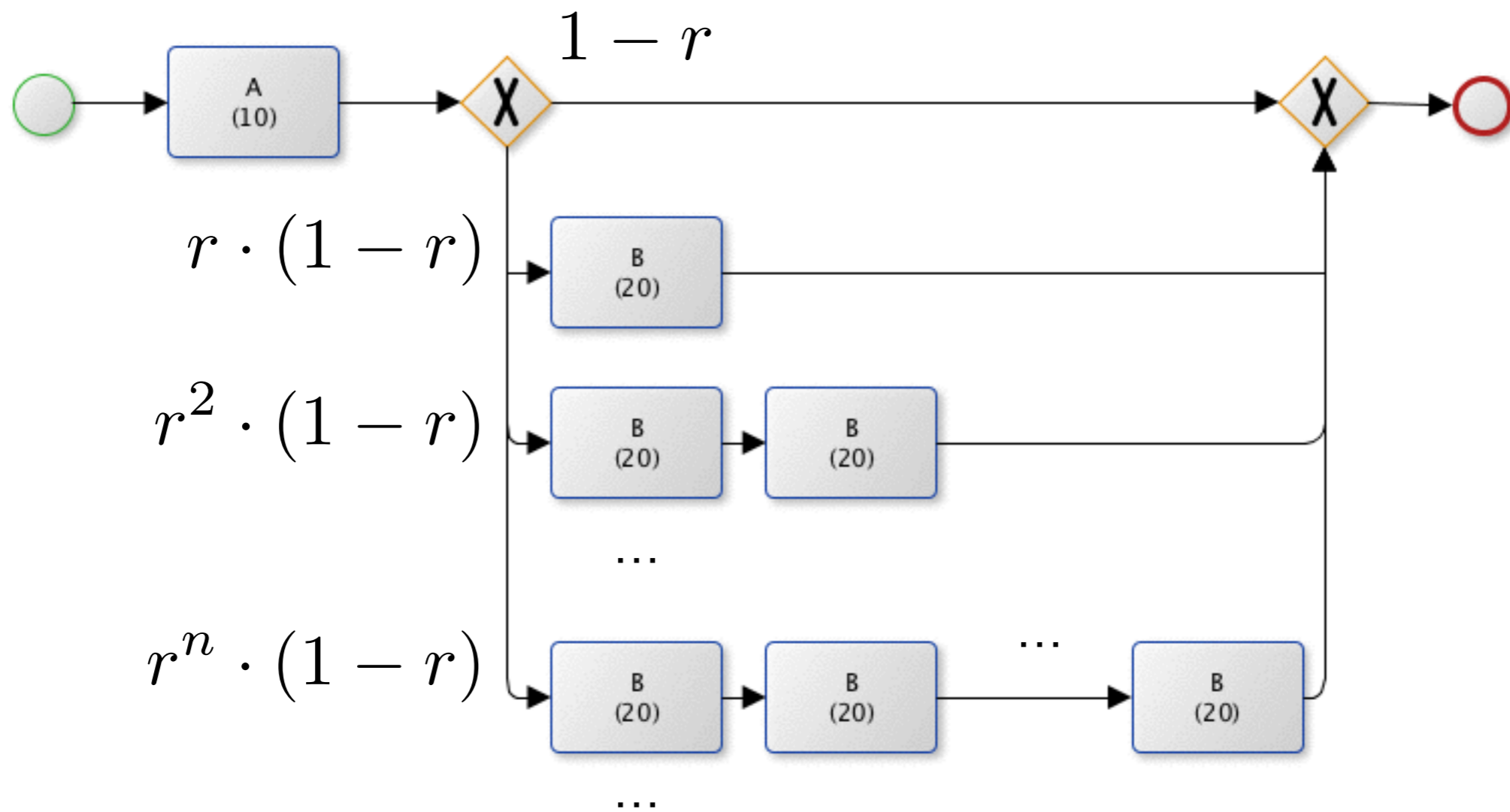
$$\begin{aligned} CT &= \sum_{i=1}^{\infty} i \cdot CT_P \cdot r^{i-1} \cdot (1-r) \\ &= CT_P \cdot (1-r) \cdot \sum_{i=1}^{\infty} i \cdot r^{i-1} \\ &= CT_P \cdot (1-r) \cdot \frac{1}{(1-r)^2} \\ &= \frac{CT_P}{1-r} \end{aligned}$$

# Rework loop (0 or more times)

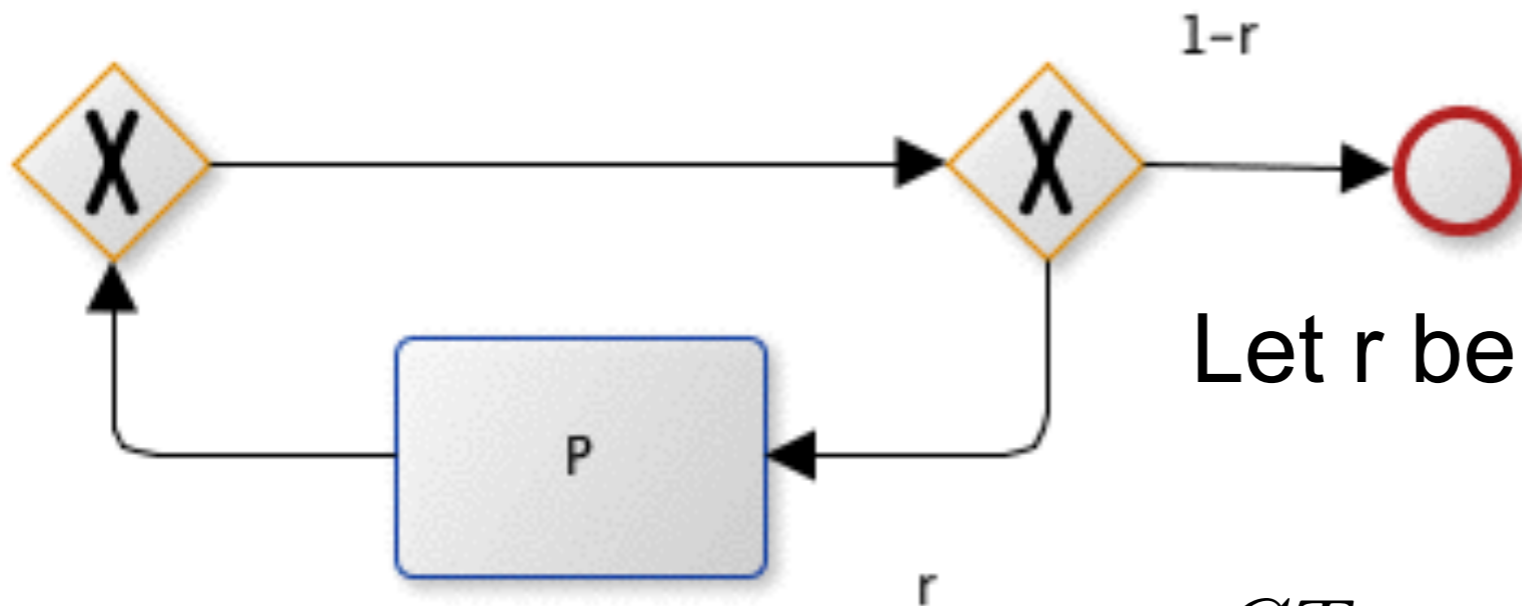


CT = ?

# Rework loop



# Rework loop

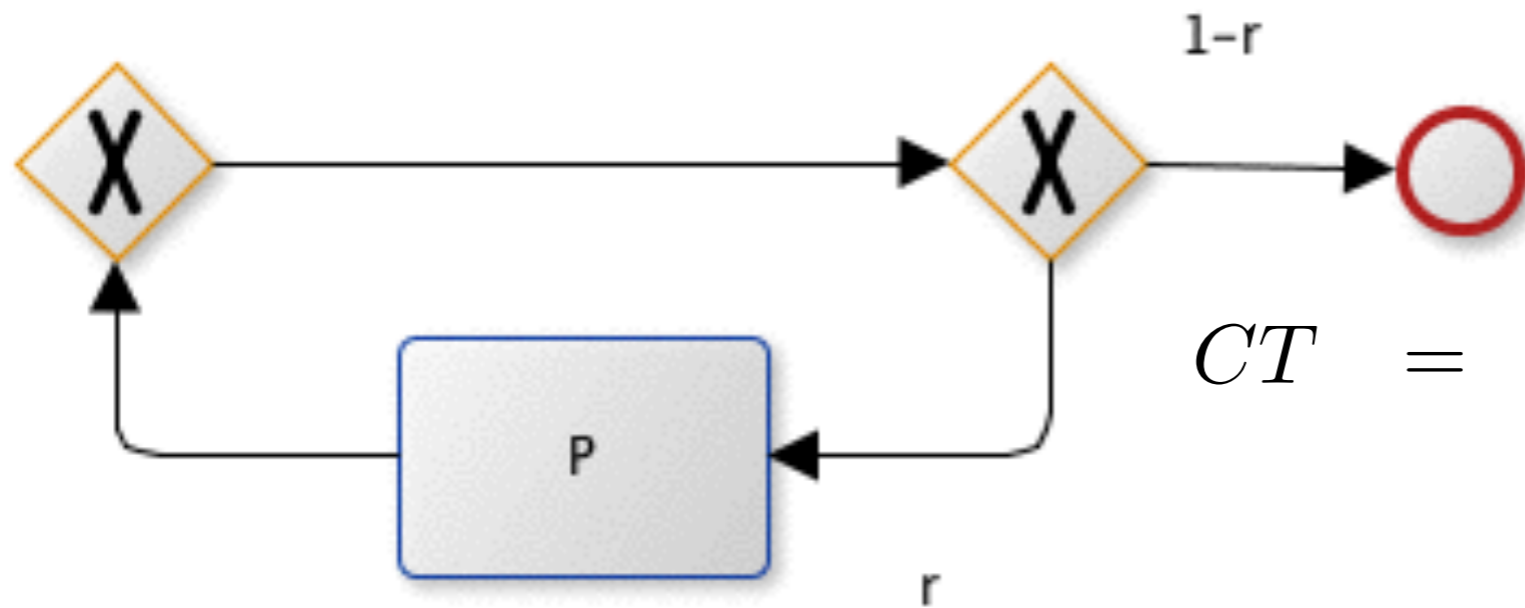


Let  $r$  be the *rework probability*, then:

$$\begin{aligned} CT &= 0 \cdot CT_P \cdot r^0 \cdot (1 - r) \\ &+ 1 \cdot CT_P \cdot r^1 \cdot (1 - r) \\ &+ \dots \\ &+ n \cdot CT_P \cdot r^{n-1} \cdot (1 - r) \\ &+ \dots \\ &= \sum_{i=0}^{\infty} i \cdot CT_P \cdot r^i \cdot (1 - r) \end{aligned}$$

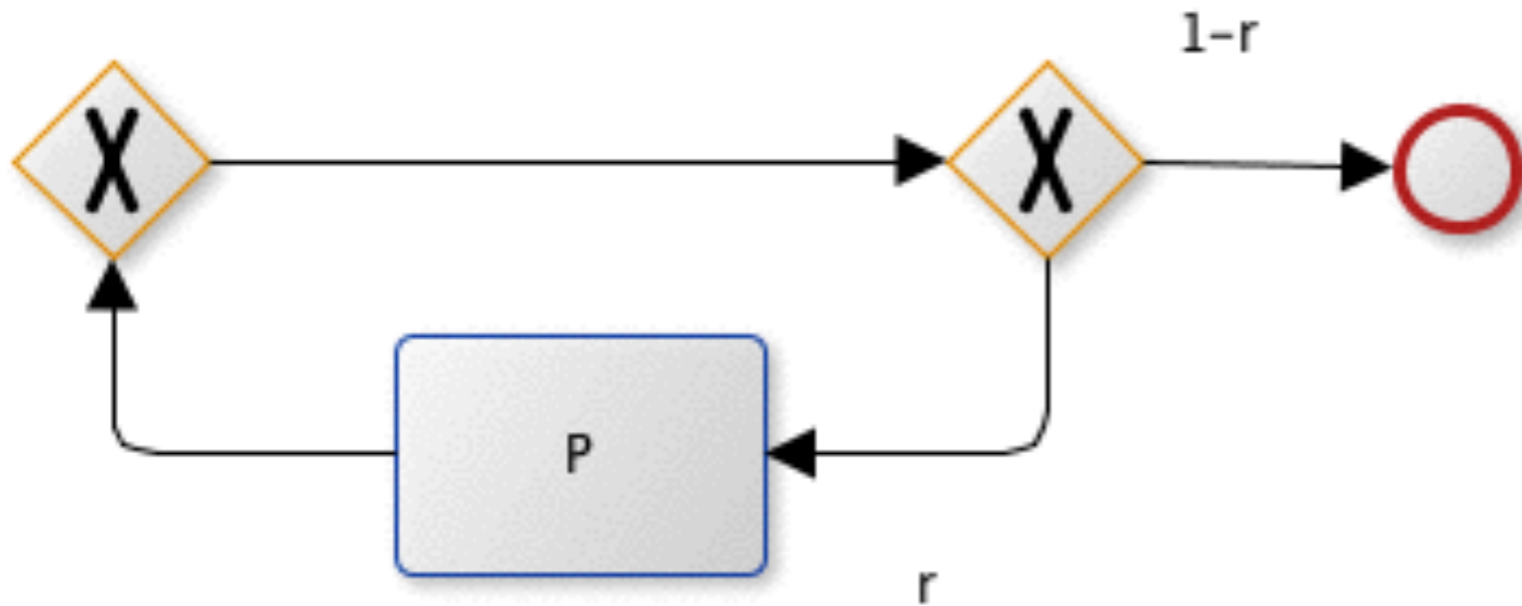


# Rework loop



$$\begin{aligned}
 CT &= \sum_{i=0}^{\infty} i \cdot CT_P \cdot r^i \cdot (1-r) \\
 &= \sum_{i=1}^{\infty} i \cdot CT_P \cdot r^i \cdot (1-r) \\
 &= CT_P \cdot r \cdot (1-r) \cdot \sum_{i=1}^{\infty} i \cdot r^{i-1} \\
 &= CT_P \cdot r \cdot (1-r) \cdot \frac{1}{(1-r)^2} \\
 &= \frac{r \cdot CT_P}{1-r}
 \end{aligned}$$

# Rework loop



Intuitively,

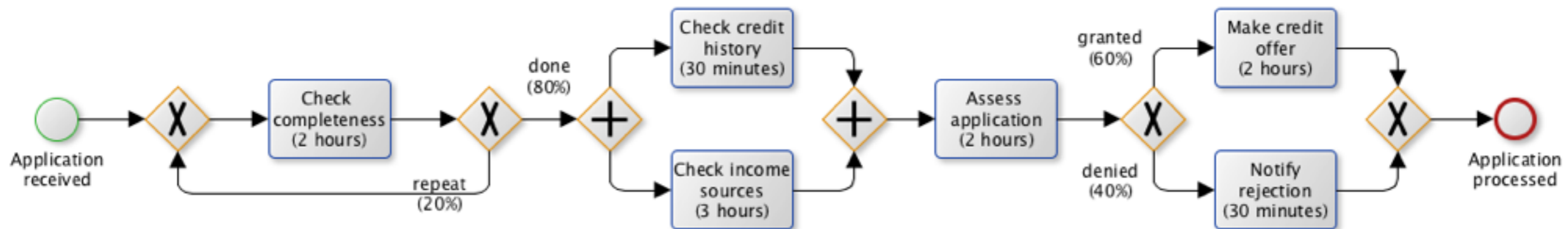
if  $\frac{CT_P}{1-r}$  is the average cycle time for reworking  $P$  one or more times,

$$\text{then } \frac{CT_P}{1-r} - CT_P = \frac{(1 - (1 - r)) \cdot CT_P}{1 - r} = \frac{r \cdot CT_P}{1 - r}$$

is the average cycle time for reworking  $P$  zero or more times

# Example

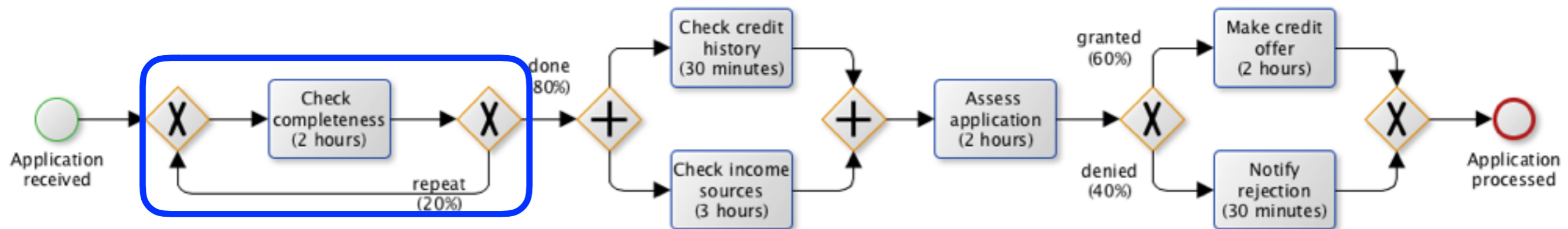
Compute the average cycle time CT of the process below



$$CT = \frac{2h}{1 - 0.2} + \max\{0.5h, 3h\} + 2h + 0.6 \cdot 2h + 0.4 \cdot 0.5h$$

# Example

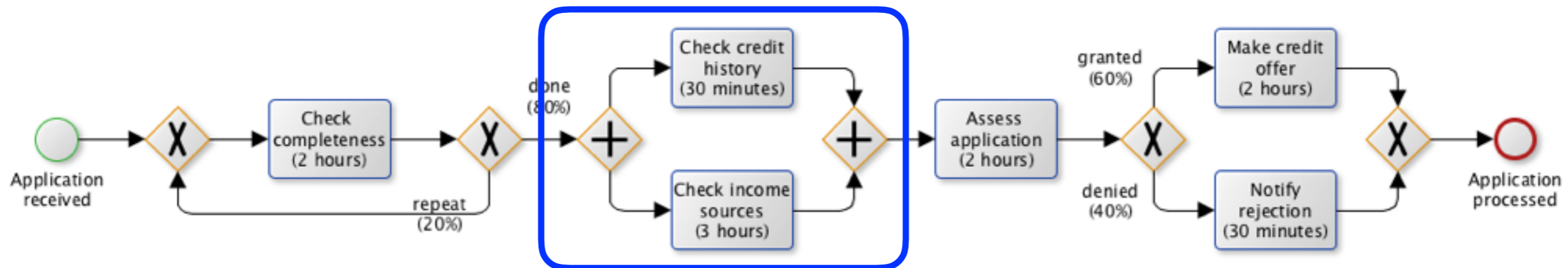
Compute the average cycle time CT of the process below



$$CT = \frac{2h}{1 - 0.2} + \max\{0.5h, 3h\} + 2h + 0.6 \cdot 2h + 0.4 \cdot 0.5h$$

# Example

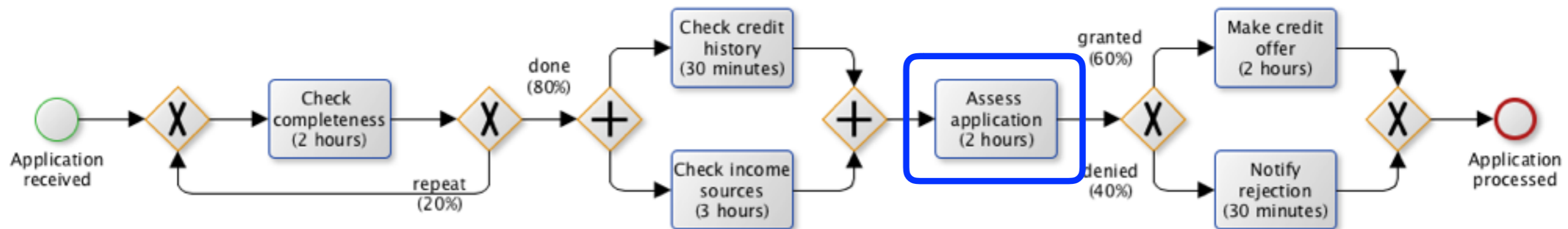
Compute the average cycle time CT of the process below



$$CT = \frac{2h}{1 - 0.2} + \max\{0.5h, 3h\} + 2h + 0.6 \cdot 2h + 0.4 \cdot 0.5h$$

# Example

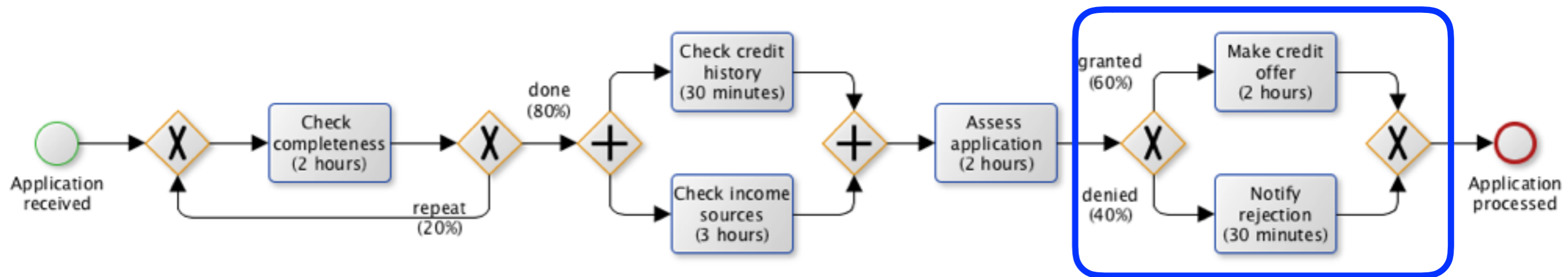
Compute the average cycle time CT of the process below



$$CT = \frac{2h}{1 - 0.2} + \max\{0.5h, 3h\} + 2h + 0.6 \cdot 2h + 0.4 \cdot 0.5h$$

# Example

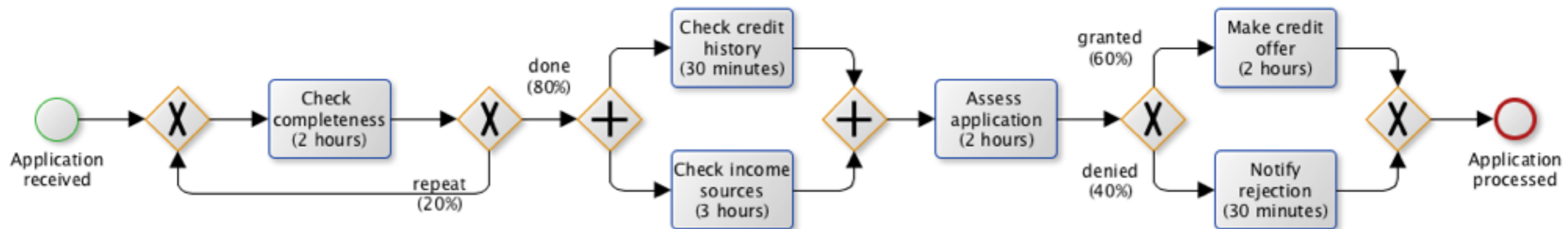
Compute the average cycle time CT of the process below



$$CT = \frac{2h}{1 - 0.2} + \max\{0.5h, 3h\} + 2h + 0.6 \cdot 2h + 0.4 \cdot 0.5h$$

# Example

Compute the average cycle time CT of the process below

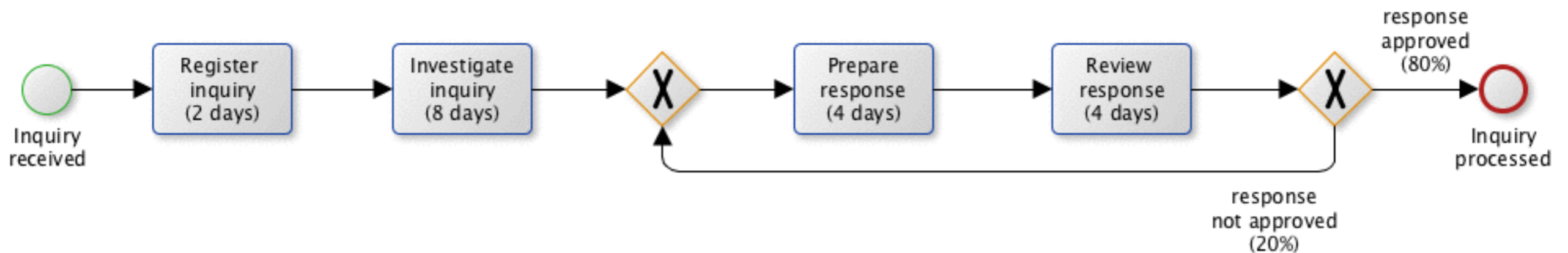


$$\begin{aligned}
 CT &= \frac{2h}{1 - 0.2} + \max\{0.5h, 3h\} + 2h + 0.6 \cdot 2h + 0.4 \cdot 0.5h \\
 &= \frac{2h}{0.8} + 3h + 2h + 1.2h + 0.2h \\
 &= 2.5h + 6.4h = 8.9h
 \end{aligned}$$



# Exercise

Compute the average cycle time CT of the process below



# Waiting vs processing

As mentioned at the beginning, the cycle time of an activity or a process can be divided into *waiting time* and *processing time*

## **Waiting time:**

is the portion of the cycle time where no work is being done to advance the process (e.g. time spent in transferring documents or waiting for an actor to perform the work)

## **Processing time:**

is the time that actors spend doing actual work

# Waiting vs processing

In most processes,  
the waiting time is a considerable portion of the cycle time!

For example,  
in many situations cases are processed in batches  
(e.g. applications, surveys)

and in many other cases actors are just not ready  
(e.g. supervisor approval, medical prescription)

# Theoretical cycle time

Assume that for each activity of the process both the processing time and the cycle time are known

Let **TCT** denote the **theoretical cycle time** of the process:  
this is computed in the same ways as CT,  
but using the processing time of activities  
(it is the amount of time a process would take on average  
if no waiting time was necessary)

# Cycle time efficiency

**Cycle time efficiency (CTE):**

is the ratio of processing time relative to the cycle time

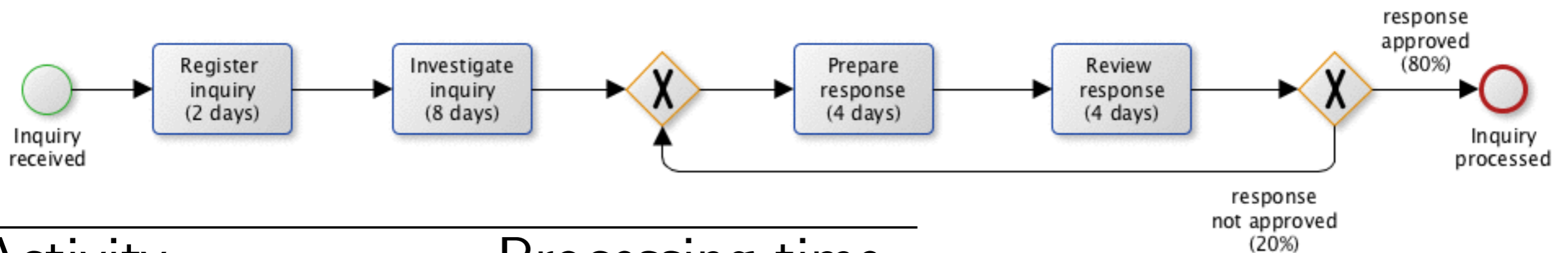
$$CTE = \frac{TCT}{CT}$$

A ratio close to 1 indicates that there is little room for improving the cycle time (unless radical changes in the process)

A ratio close to zero indicates that there a significant amount of room for improving cycle time (by reducing the waiting time)

# Exercise

Compute the TCT and CTE of the process below, given the processing times reported in the table (assume 1 day = 8 working hours)



Activity	Processing time
Register inquiry	30 minutes
Investigate inquiry	12 hours
Prepare response	4 hours
Review response	2 hours

# Limitation of flow analysis

# Pitfalls and limitation

The equations we have presented deals with *block-structured* process only: we cannot calculate the cycle time for any processes.

The equations exploit the average cycle time of activities: we need to estimate such values (interviewing stakeholders, inspecting logs)

Flow analysis does not account for the fact that a process can behave differently depending on the load: when the load goes up and the resources are constant, the waiting time increases.



# Little's law

# Arrival rate and Work-In-Process

Cycle time is directly related to two other important measures:

**Arrival rate  $\lambda$**  of a process:

is the average number of new instances of the process (i.e. cases) that are created per time unit

**Work-In-Process (WIP):**

is the average number of process instances (i.e. cases) that are active (i.e. not yet completed) at a given point in time

# Little's law

In a paper from 1954, operation research professor John Little assumed (without giving a proof) that the following equality holds in a *stable*\* system:

the long-term average number of customers (**WIP**)  
is equal to

the long-term average effective arrival rate ( $\lambda$ ) multiplied  
by the average time a customer spends in the system (**CT**)

$$\text{algebraically: } \mathbf{WIP} = \lambda \cdot \mathbf{CT}$$

\* *stable* means that the number of customers in the system is not increasing infinitely

$$WIP = \lambda \cdot CT$$

Little's law tell us that:

WIP increases

if the cycle time (CT) increases

or if the arrival rate ( $\lambda$ ) increases

(if the process slow down there will be more active cases and the faster new cases are created the higher will be the number of active instances)

If the arrival rate ( $\lambda$ ) increases and we want to keep WIP constant, then we must decrease the cycle time (CT) (i.e., we must work faster)

# A note on Little's law

The law is classically stated using different symbols

$$L = \lambda \cdot W$$

In a subsequent paper from 1961,  
John Little proved the equality  
later followed by simpler proofs in 1967, 1969, 1972

Since we can estimate WIP and  $\lambda$  by observing the system, we can use Little's law as an alternative way to calculate the average cycle time CT:

$$CT = \frac{WIP}{\lambda}$$

# Example

Assume there are 250 business days per year.

If the total number of applications received over the last year is 2500 we can infer that the average number of *applications per day* is 10 (i.e.  $\lambda=10$ ).

By sampling (e.g. checking every week), we observed that on average there were 200 *applications concurrently active* (i.e.  $WIP=200$ ).

$$CT = \frac{WIP}{\lambda} = \frac{200}{10} = 20 \text{ days}$$

# Exercise

A restaurant receives on average 1200 customers per day (from 10am to 10pm).

During peak times (12pm to 3pm, and 6pm to 9pm) the restaurant receives around 900 customers and, on average, 90 customers can be found in the restaurant at a given time.

At non-peak times, the restaurant receives 300 customers in total and, on average, 30 customers can be found in the restaurant at a given time.

What are the average times that a customer spends in the restaurant during peak/non-peak times?

# Exercise (continued)

The maximum capacity of the restaurant is sometimes reached during peak times.

The restaurant manager expects that the number of customers during peak times will increase slightly in the coming months.

What action can be taken to address this issue without investing in extending the building?



# Cost analysis

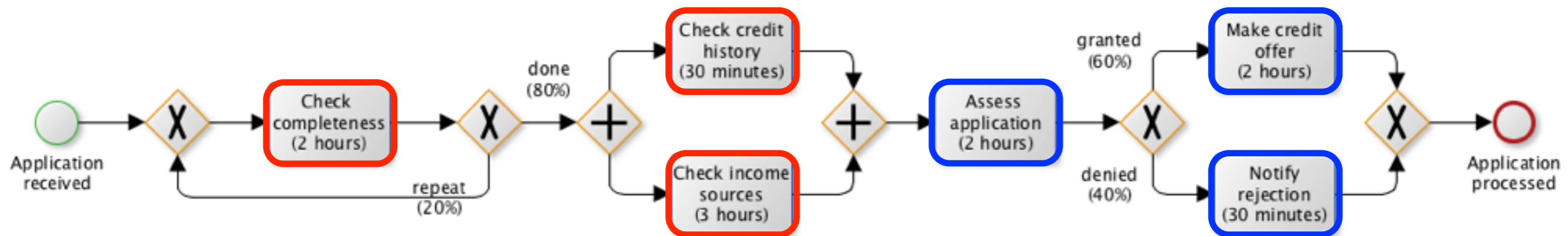
# Cost analysis

Analogously to the case of cycle time computation, flow analysis can be used to calculate other performance measures.

If we know the average cost of each activity, then we can calculate the average cost of the process more or less as we have just seen.

In fact the formulas for sequences, XOR-blocks and reworks are the same, but **for AND-blocks we need to take the sum** (instead of max)

# Example



Assume

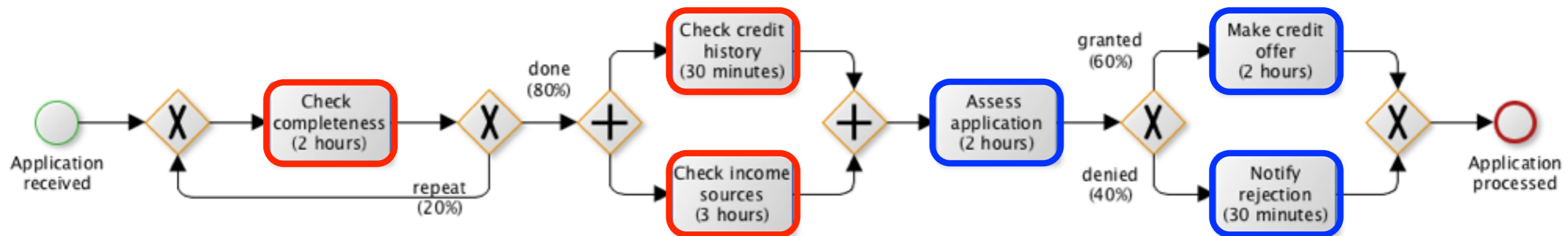
activities are annotated with processing time, "red" activities are performed by a clerk (hourly cost **25€**), while "blue" activities by a credit officer (hourly cost **50€**).

Assume also that

the bank is charged **1€** for each "credit history check".

What is the average cost of the process?

# Example (continued)

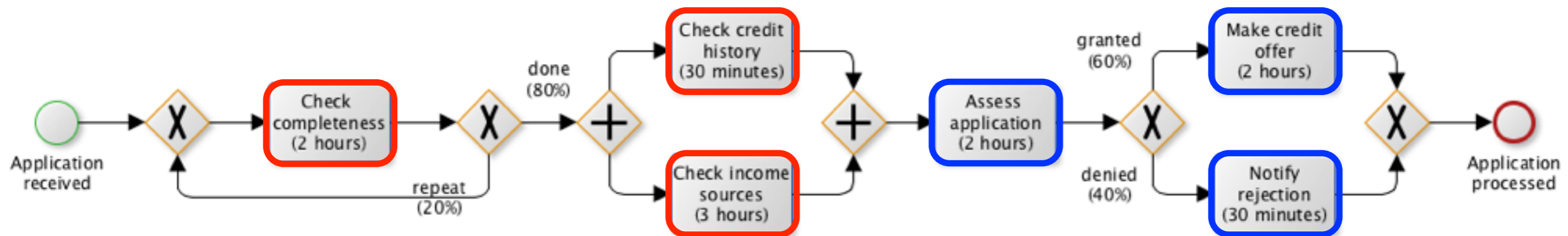


We can distinguish two kinds of costs:

**human resource costs:** can be calculated as the product of the (hourly) cost and the processing time of the task

**other costs:** fixed costs that are incurred by an execution of a task (not related to the time spent by human resources)

# Example (continued)

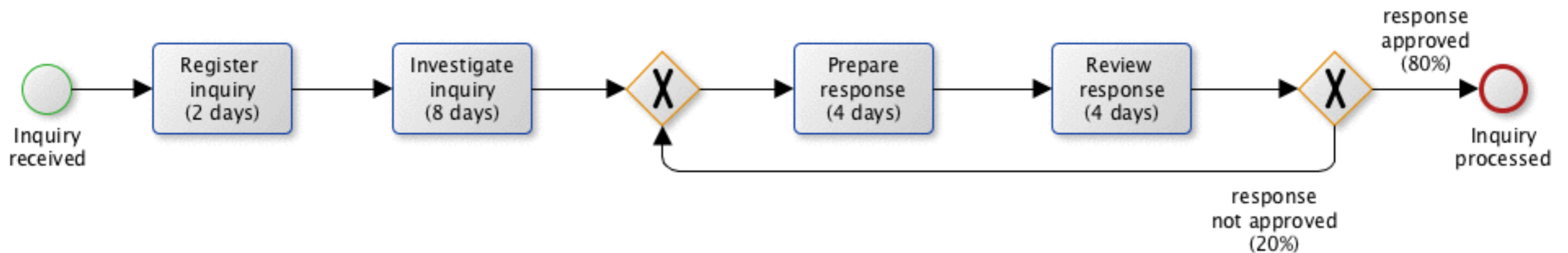


$$Cost = \frac{50}{1 - 0.2} + 13.5 + 75 + 100 + 0.6 \cdot 100 + 0.4 \cdot 25 = 321$$

Activity	Resource cost	Other cost	Total cost
Check completeness	$2 \cdot 25 = 50$		50
Check credit history	$0.5 \cdot 25 = 12.5$	1	13.5
Check income resources	$3 \cdot 25 = 75$		75
Assess application	$2 \cdot 50 = 100$		100
Make credit offer	$2 \cdot 50 = 100$		100
Notify rejection	$0.5 \cdot 50 = 25$		25

# Exercise

Compute the average total cost of the process below



Activity	Resource	Resource hourly cost	Processing time
Register inquiry	clerk	25 euros	30 minutes
Investigate inquiry	advisor	50 euros	12 hours
Prepare response	senior advisor	75 euros	4 hours
Review response	counselor	100 euros	2 hours