

• current = ϕ

In the nodes are indicated $\pi(u)$, whereas outside there is u .

| min π | max π | $c(u,v)$ | u | v |
|-----------|-----------|----------|-----|-----|
| 1 | 5 | 3 | C | D |
| 2 | 5 | 1 | B | C |
| 2 | 6 | 2 | A | B |
| 2 | 7 | 3 | B | F |
| 4 | 6 | 6 | A | E |
| 5 | 7 | 5 | C | F |
| 6 | 7 | 4 | A | F |

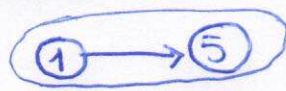
this is actually a priority queue



So we sort it first by min π then by $c(u,v)$.

Extract min Q \rightarrow $\langle 1, 5, 3, C, D \rangle$ $\rightarrow u=1, v=5$

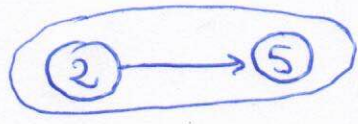
since current $\neq 1 = u \rightarrow$ current = 1, ~~relinkTo~~ relinkTo = 5
output $\langle C, D \rangle$



then two nodes are merged as indicated

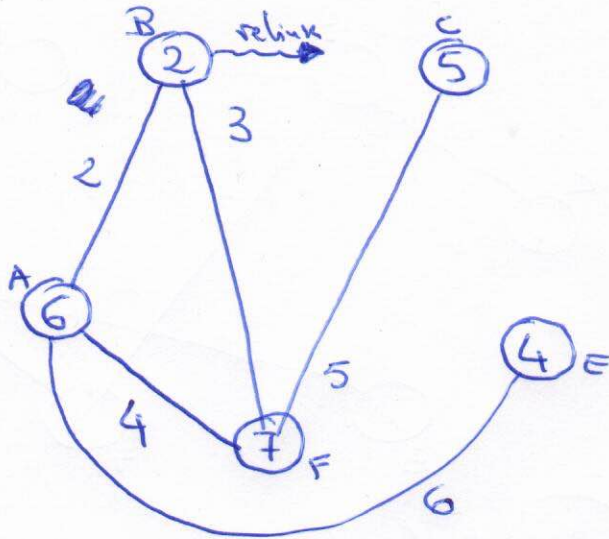
Extract min Q \rightarrow $\langle 2, 5, 1, B, C \rangle$ $\rightarrow u=2, v=5$

since current $\neq 2 = u \rightarrow$ current = 2, relinkTo = 5
output $\langle B, C \rangle$



all edges incident in 2 will be moved to 5 by the next steps.

We draw the graph for simplicity at this point:



| min π | max π | c | u | v |
|-----------|-----------|---|---|---|
| 2 | 6 | 2 | A | B |
| 2 | 7 | 3 | B | F |
| 4 | 6 | 6 | A | E |
| 5 | 7 | 5 | C | F |
| 6 | 7 | 4 | A | F |

current = 2, relink to = 5

The algorithm at this point proceeds as follows:

min Q \rightarrow $\langle 2, 6, 2, A, B \rangle \rightarrow u=2, v=6$

since current = 2 \Rightarrow we have to relink the edge changing $u \rightarrow$ relink to

Q.insert (5, 6, 2, A, B)

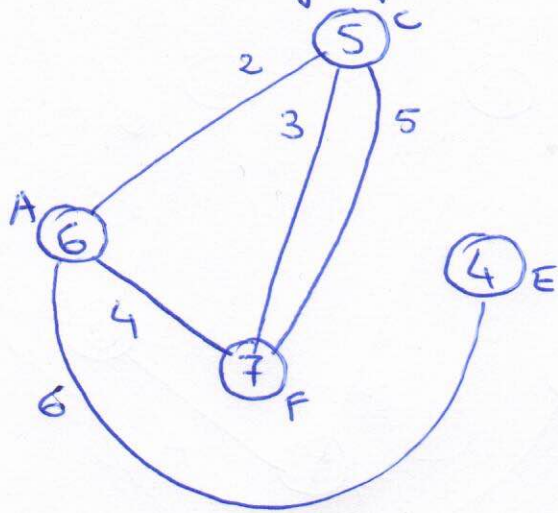
Before plotting the next graph and priority queue, let us do another step, since the next edge is also outgoing from 2 and thus it'll be relink to 5;

min Q \rightarrow $\langle 2, 7, 3, B, F \rangle \rightarrow u=2, v=7$

since current = 2 = u \Rightarrow we have to relink 2-7 to 5-7

Q.insert (5, 7, 3, B, F)

We draw the graph of this point, and the current PQ



| min π | max π | c | u | v |
|-----------|-----------|---|---|---|
| 4 | 6 | 6 | A | E |
| 5 | 6 | 2 | A | B |
| 5 | 7 | 3 | B | F |
| 5 | 7 | 5 | C | F |
| 6 | 7 | 4 | A | F |

current = 2, relinkTo = 5

the algorithm at this point does:

min Q \rightarrow $\langle 4, 6, 6, A, E \rangle \rightarrow u = 4, v = 6$

since current $\neq 4 = v \rightarrow$ we have to ~~relink~~ change



current = 4
relinkTo = 6

output $\langle A, E \rangle$

This way $\textcircled{4} \rightsquigarrow \textcircled{6}$ hence every edge linked to $\textcircled{4}$ has to be relinked to $\textcircled{6}$, if any (we actually do not have them).

notice that

min Q is extracted and thus removed from Q

min Q \rightarrow $\langle 5, 6, 2, A, B \rangle \Rightarrow u = 5, v = 6$

since current $\neq 5 = u \rightarrow$ current = 5, relinkTo = 6

This way $\textcircled{5} \rightsquigarrow \textcircled{6}$

hence every edge incident in $\textcircled{5}$ will be relinked to $\textcircled{6}$

output $\langle A, B \rangle$

extract min Q

min Q \rightarrow $\langle 5, 7, 3, B, F \rangle \rightarrow u=5, v=7$

since current = 5 = u \Rightarrow we have to rethink



Q. insert (6, 7, 3, B, F)

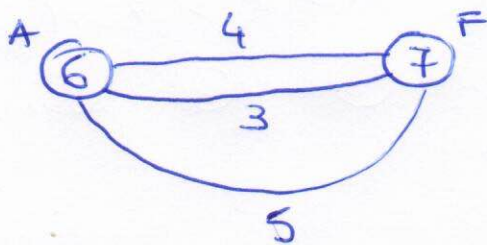
min Q \rightarrow $\langle 5, 7, 5, C, F \rangle \rightarrow u=5, v=7$

since current = 5 = u \Rightarrow we have to rethink



Q. insert (6, 7, 5, C, F)

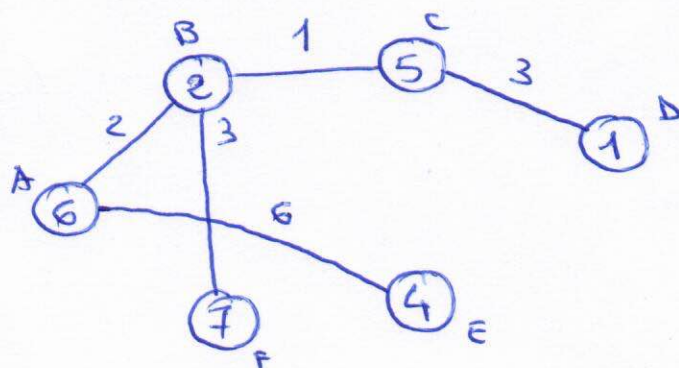
We draw the current graph and PQ



| min | max | c | u | v |
|-----|-----|---|---|---|
| 6 | 7 | 3 | B | F |
| 6 | 7 | 4 | A | F |
| 6 | 7 | 5 | C | F |

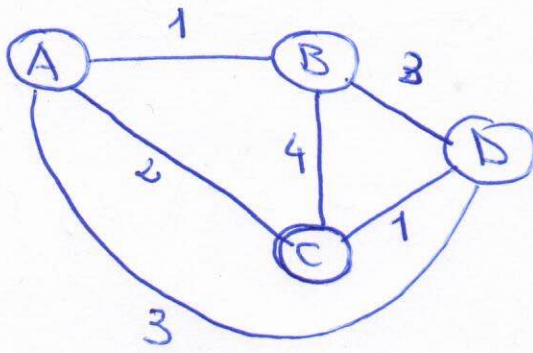
If $M=2$ then we can solve MST in memory of this graph hence output $\langle B, F \rangle$ which corresponds to

$\langle 6, 7, 3, B, F \rangle$
 \uparrow
 min weight



MST

Esercizio 2



current = ϕ
 π = Identical Perm.

| min π | max π | cost | u | v |
|-----------|-----------|------|---|---|
| A | B | 1 | A | B |
| A | C | 2 | A | C |
| A | D | 3 | A | D |
| B | D | 3 | B | D |
| B | C | 4 | B | C |
| C | D | 1 | C | D |

min Q $\rightarrow \langle A, B, 1, A, B \rangle \rightarrow u = A, v = B$

since current $\neq A = u \rightarrow$ current = A, relinkTo = B
 output $\langle A, B \rangle$
 extract min

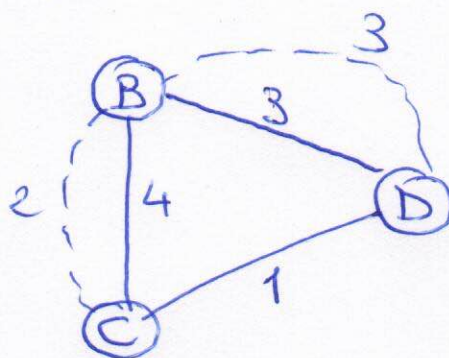
min Q $\rightarrow \langle A, C, 2, A, C \rangle \rightarrow u = A, v = C$

since current = A = u \Rightarrow Q.insert(B, C, 2, A, C)
 [we have relinked $\textcircled{A} \rightarrow \textcircled{B}$]

min $Q \rightarrow \langle A, D, 3, A, D \rangle \rightarrow u=A, v=D$

since current = $A = u \Rightarrow Q.inset(B, D, 3, A, D)$
 [we have relinked $A \rightarrow B$]

The new graph is:



0-----0
relinked edge

| min π | max π | cost | u | v |
|-----------|-----------|------|---|---|
| B | C | 2 | A | C |
| B | C | 4 | B | C |
| B | D | 3 | B | D |
| B | D | 3 | A | D |
| C | D | 1 | C | D |

goes here



min $Q \rightarrow \langle B, C, 2, A, C \rangle \Rightarrow u=B, v=C$

since current $\neq B = u \rightarrow$ current = B, relink to = C
 output $\langle A, C \rangle$
 extract min

This goes *

min $Q \rightarrow \langle B, C, 4, B, C \rangle \Rightarrow u=B, v=C$

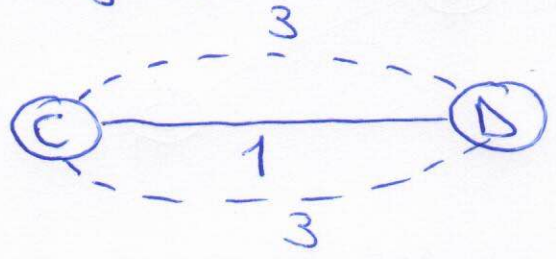
since current = $B = u \Rightarrow$ ~~current = C, relink to = C~~
 no op because
 relink to = $C = v$.

$\min Q \rightarrow \langle B, D, 3, B, D \rangle \Rightarrow u = B, v = D$
since $\text{current} = B = u \Rightarrow Q.\text{insert}(C, D, 3, B, D)$
 =

$\min Q \rightarrow \langle B, D, 3, A, D \rangle \Rightarrow u = B, v = D$
since $\text{current} = B = u \Rightarrow Q.\text{insert}(C, D, 3, A, D)$

* MOVED EDGE

The current graph and priority queue is:



| min π | max π | cost | u | v |
|-----------|-----------|------|---|---|
| C | D | 1 | C | D |
| C | D | 3 | A | D |
| C | D | 3 | B | D |

If $M=2$ then we apply Kruskal and thus
 output $\langle C, D, 1, C, D \rangle$

