



The MPI Message-passing Standard Practical use and implementation (VI)

SPD Course 23/03/2016 Massimo Coppola









Datatypes

REFINING DERIVED DATATYPES LAYOUT FOR COMPOSITION







Derived D.type Extent and composition



- For complex derived datatypes, extent plays an important role
 - Plain definition: distance between the first (smallest address) byte and last (largest address) byte used in memory
 - Actual use: the offset between two items of the given datatype when they are stored consecutively in memory
 - E.g. whenever a contiguous datatype is created or a communication buffer with more instances is used
- Setting extent manually (MPI1, MPI2>)
- Querying extent
- Examples with derived datatypes





Get extent



- Get the lower bound and extent of a datatype
 - By default, lower bound = lowest-address location of a datatype
 - Extent = distance from lower bound to highest address location used by the datatype







Modify extent



- Modify the lower bound and extent of a datatype
- Reset lower bound and extent of the datatype to new arbitrary values, for the sake of data structures composition





Retrieve original extent



- Retrieve the true lb and extend from a datatype
- MPI keeps those as it needs them for the actual packing and unpacking needed by communication







Intracommunicators

COLLECTIVE PRIMITIVES WITH COMMUNICATION AND COMPUTATION







Reduce



- int MPI_Reduce(
 const void* sendbuf, void* recvbuf, int count,
 MPI_Datatype datatype,
 MPI_Op op, int root, MPI_Comm comm)
- reduce operation across all processes of a Communicator
 - Reduces the elements in the same position of each process' buffer, leaving results in root's buffer
- count, datatype, op, root, comm arguments must match
 - If count == 1 we have a classical element-wise reduction
 - If count>1 we have several reductions at the same time
- As with any collective, the communication pattern is implementation dependent (but is op commutative ?)
- MPI provides most basic operators
 - Operators are associative
 - Operators may be commutative → potential optimizations
 - Note that: floating point op.s may be seen as non-commutative
 - Datatype must be compatible with op





MPI Scan/Reduce operators



- Arithmetic operations
 - MPI_MAX, MPI_MIN, MPI_SUM, MPI_PROD
 - Generally allowed on MPI integral and floating point types (including complex)
- Logic (L) and bit wise (B) operations
 - MPI_LAND, MPI_LOR, MPI_LXOR
 - Generally allowed on C integers and on logical types
 - MPI_BAND, MPI_BOR, MPI_BXOR
 - Generally allowed on C/Fortran integers







MPI_MINLOC and MAXLOC



- Operators defined on couples (value, index)
 - MPI_MAXLOC, MPI_MINLOC
 - Value is any integral or floating point type
 - Index is an integral type
 - chars used as integers require special attention
 - e.g. explicitly using MPI_SIGNED _CHAR / UNSIGNED_CHAR
- MINLOC: compute the global minimum of v and the index attached to it
- MAXLOC: compute the global maximum of v and the index attached to it
- Lexicographic order
 - when more values hit the minimum (maximum) the lower one is chosen
- Example application
 - pass (value, rank) to detect the rank of the process with the minimum/maximum value







MPI couple types



 These couple types are supported by the MPI_MINLOC and MPI_MAXLOC operators

- MPI_FLOAT_INT
- MPI LONG INT
- MPI_DOUBLE_INT
- MPI_SHORT_INT
- MPI_2INT

- struct { float, int }
- struct { long, int }
- struct { double, int }
- struct { short, int }
- struct { int, int }
- MPI_LONG_DOUBLE_INT
 - struct { long double, int }
 - this is an OPTIONAL type







Operators semantics



- Operators are called within the reduction collective by the instances of the MPI library of the processes of the program
- Each operators receives two local buffers and performs a reduction step on their contents
 - The buffers are possibly allocated by the library implementation as temporaries
 - Many operators are polymorphic, so they have to detect the type of data in the buffer
 - Datatype is a parameter passed from the collective down to the operator, but remember it is a handle
 - Easy case: MPI basic datatypes are globally known to MPI runtime and to the program
 - Besides, MPI standard operators are easy







MPI operators and computing-collective primitives



- MPI operators (including user-defined ones) are used by all MPI collectives performing distributed computation
 - MPI_REDUCE, MPI_ALLREDUCE,
 MPI_REDUCE_SCATTER_BLOCK,
 MPI_REDUCE_SCATTER,
 MPI_SCAN, MPI_EXSCAN
 - All the non blocking version of those collectives (since MPI-3)
 - MPI_REDUCE_LOCAL (special case actually designed for MPI implementers)







User-defined operators



- MPI allows you to define your own operators
 - They can apply to basic and user-defined datatypes
- What do you need to do
 - (Possibly) provide relevant datatype definitions
 - Provide MPI with a definition of the operator
 - a compiled function with a specific signature
 - the operator definition this is local to each process
 - Detect and recognize the MPI_Datatype within the operator code
 - To detect errors
 - If the operator needs to be polymorphic
 - Combine each couple of elements in the same position of the two input buffers
- Operator code can call **no** MPI communication primitives; only MPI_ABORT() in case of error







MPI_Op_create



- int MPI_Op_create(MPI_User_function* user_fn, int commute, MPI_Op* op)
- MPI primitive for defining operators
- Takes a user function pointer as first argument
- Can specify non-commuting operators
- Returns the operator handle
- MPI_Op_free allows to free operators





Operator signature



- row-wise combines data from two buffers
 - results are placed in the second buffer
- The datatype handle comes from the collective call (e.g. reduction) and may be unknown at compile time
 - For user-defined datatypes, polymorphic operators need to access a table of datatypes handles that are defined by the program





Example: rewriting MPI_SUM



```
/* this follows the MPI User function typedef */
void my sum op(void * b in, void * b inout,
               int * count, MPI Datatype * d) {
    if (d == MPI INT) {
       for (i=0; i<count; i++) {
           ((int*)b_inout)[i]+=((int *)b_in)[i]; }
    } else if (d == MPI FLOAT) {
       for (i=0; i<count; i++) {
           ((float*)b_inout)[i]+=((float *)b_in)[i];}
    } else MPI_Abort (MPI_COMM_WORLD, -12345);
MPI Op * op sum;
MPI Create op (* my sum op, MPI FALSE, op sum)
```

- Very limited example: it only accepts INT and FLOAT types
- Can call specialized functions in each case (code reuse, hardware acceleration)





Datatypes and polymorphic operators



- Check the datatype
 - Compare the received datatype handle to a list of allowed handles, execute proper code
 - Simple if/else error if only one type is allowed
- Check and switch for polymorphic op.s
 - Operators that can handle several datatypes should employ data structures that avoid any excessive comparison overhead
 - E.g. an hash-map (perfect hash?) associating handles with code (function pointers) implementing each case of use of the operator
 - The overhead is usually negligible with respect to the communication overhead of a reduction or scan







Example: rewriting MPI_SUM (II)



```
/* this follows the MPI User function typedef */
void my_sum_op(void * b_in, void * b inout,
               int * count, MPI Datatype * d) {
   int my type = my hashtable get(d);
   case (SUMOP INT T) : // MPI INT
       sumintarrays((int *)b in, (int*) b inout, count);
       break;
   case (SUMOP FLOAT T) : // MPI FLOAT
       sumfloatarrays((float *) b in, (float *) b inout, count);
           break;
   case (SUMOP_4BY4_FLOAT_T) : // user type example, 4*4 matrix
       sumfloatarrays((float *)b in, (float*)b inout, count*16);
       break;
   default: MPI Abort (MPI COMM WORLD, -12345);
```

- In the example we assume that
 - a hashtable is filled with custom values for each recognized datatype
 - if the type is not in the hashtable, the default value returned (unknown datatype) triggers the default case





Scan



- Applies a scan (parallel prefix) to the elements in corresponding position of the send buffers of the processes
- The scan works according to process rank
 - Process i receives the result of the combination of data from processes { 0, ... i }
 - Scan is the identity for process with rank 0
- If count>1 we have multiple scans within the same communication pattern
- With MPI_INPLACE in sendbuf, only the receive buffer is used





ExScan



- int MPI_Exscan(const void* sendbuf, void* recvbuf, int count,
 MPI_Datatype datatype, MPI_Op op,
 MPI_Comm comm)
- Same as MPI_Scan, but results are accumulated on the following process
 - Process with rank i gets the parallel prefix result of data contributed from processes 0.. i-1
 - Mnemonics: think of it as a "scan and shift"
 - Process 0 receives no data, and its receive buffer is not used by MPI_Exscan
 - MPI_IN_PLACE can be used in sendbuf





Other reduce collective operations



- MPI_Allreduce
 - Semantically equivalent to a reduce followed by a broadcast
 - May be implemented more efficiently, of course
- MPI_Reduce_scatter_block
 - Performs a reduction, then scatters the result buffer across the processes
 - Requires n*recvcount elements by each process, scatters the n blocks of recvcount elements of the result
 - Parameter recvcount is the number of elements received per process after the scatter
 - the overall reduction is computed on recvcount*N elements, where N is the communicator size.
- MPI_Reduce_scatter
 - Generalizes the scatter_block to a variable scatter (each process can receive a block of different size)
 - The recvcount is now an array of block sizes (the array is the same size as the communicator, see MPI_Scatterv)





References



With respect to MPI-3 standard

- Section 5.9 (Global reduction operators)
- You can skip reduce_local

