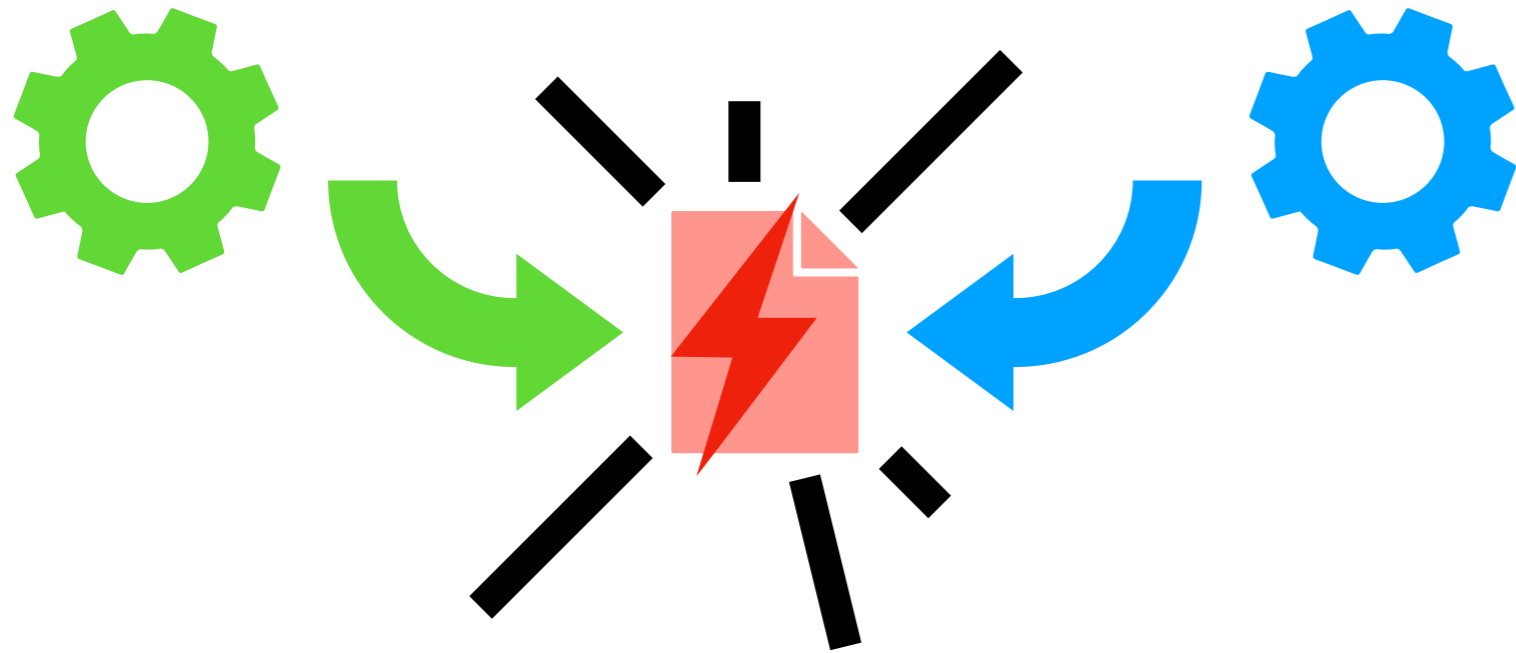


Linguaggi di Programmazione



Roberta Gori

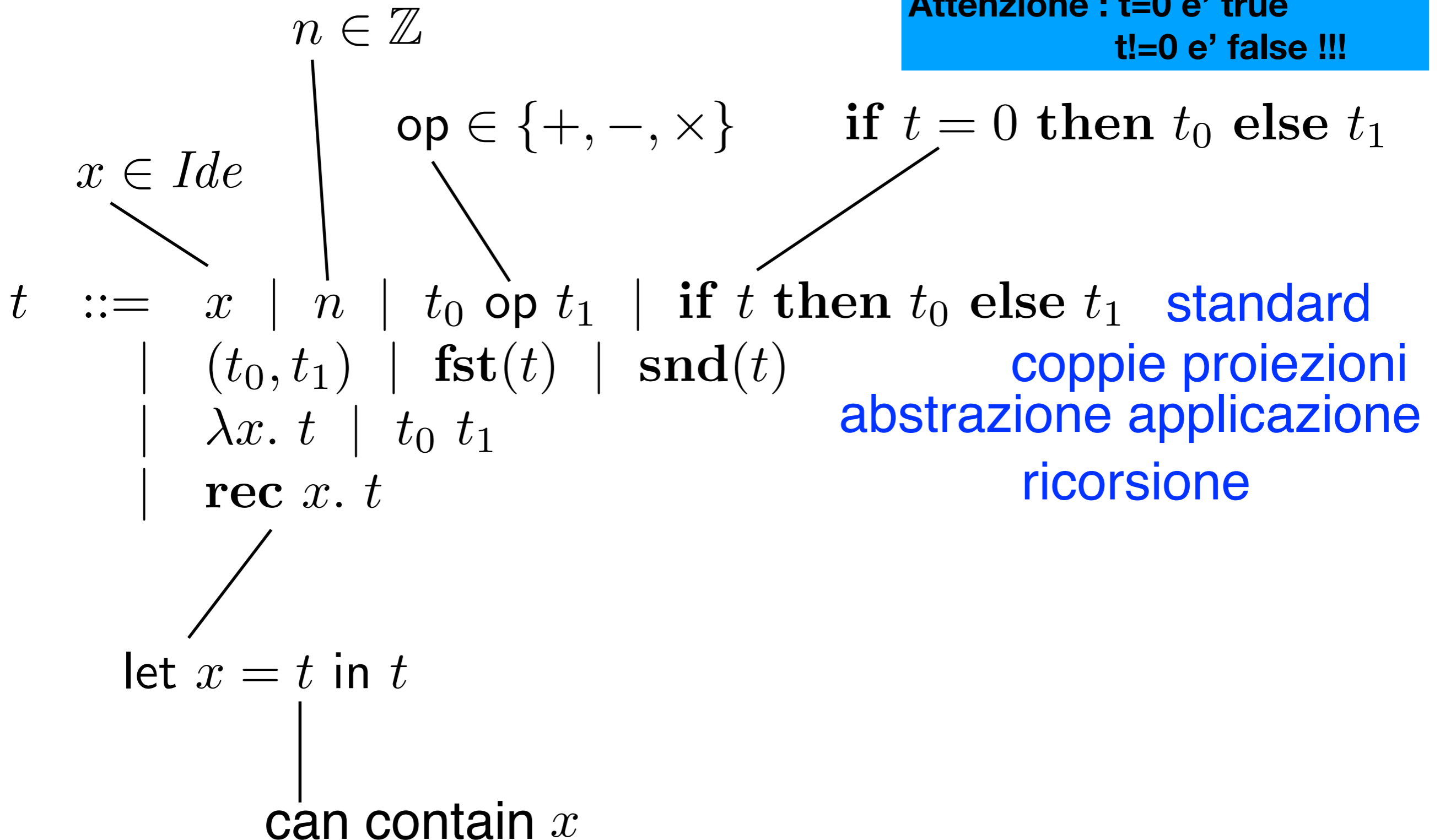
HOFPL Sintassi e Tipi -7.1

HOFL pre-termini

(linguaggio funzionale di ordine superiore)

Sintassi HOFL

Attenzione : $t=0$ e' true
 $t!=0$ e' false !!!



Esercizio

rec f . λx . **if** x **then** 1 **else** $x \times (f (x - 1))$

quale e' il significato del precedente pre-termine?

Definisce la funzione fattoriale

Esercizio

rec *rep.* $\lambda n. \lambda f. \lambda x.$ **if** n **then** x
else f (*rep* $(n - 1)$ f x)

quale e' significato del precedente pre-termini?

$$\text{rep } n \ f \ x = f^n \ x$$

Esercizio

$$\lambda x. \left(\left(\begin{array}{l} \text{rec } f. \lambda y. \text{ if } (x - y) \text{ then } 0 \\ \text{else if } (x + y) \text{ then } 1 \\ \text{else } f (y + 1) \end{array} \right) 0 \right)$$

quale e' significato del precedente pre-termine?

maggiore o uguale a 0

Esercizio (da consegnare)

assumiamo $true = 0$

$false = \text{qls } n \neq 0$

riempire al posto dei puntini (in HOFL)

$or \triangleq \lambda n. \lambda m. \dots$

$and \triangleq \lambda n. \lambda m. \dots$

$not \triangleq \lambda m. \dots$

$implies \triangleq \lambda n. \lambda m. \dots$

$iff \triangleq \lambda n. \lambda m. \dots$

Pre-termini

$t ::= x \mid n \mid t_0 \text{ op } t_1 \mid \text{if } t \text{ then } t_0 \text{ else } t_1$
 $\mid (t_0, t_1) \mid \text{fst}(t) \mid \text{snd}(t)$
 $\mid \lambda x. t \mid t_0 t_1$
 $\mid \text{rec } x. t$

Perche' sono chiamati pre-termini?

$x + 1$ ✓

✗ $1 + (0, 5)$

if x **then** $x + 1$ **else** $x - 1$ ✓

✗ $2 \times \lambda x. x$

$(0, \lambda x. x)$ ✓

abbiamo bisogno
di un sistema di tipi

✗ $3 \lambda x. x + 1$

fst $(0, \lambda x. x)$ ✓

✗ **fst** (3)

$(\lambda x. x + 1) 3$ ✓

✗ **if** x **then** $\lambda x. x$ **else** (x, x)

rec $f. \lambda x. x + (f 0)$ ✓

✗ **rec** $f. \lambda x. f + x$

tipi HOFL

Tipi

$$t ::= x \mid n \mid t_0 \text{ op } t_1 \mid \text{if } t \text{ then } t_0 \text{ else } t_1$$
$$\mid (t_0, t_1) \mid \text{fst}(t) \mid \text{snd}(t)$$
$$\mid \lambda x. t \mid t_0 t_1$$
$$\mid \text{rec } x. t$$

quali tipi?

infinite combinazioni!

coppie

int

int * *int*

funzioni

int → *int*

int * (*int* → *int*)

(*int* * *int*) → *int*

int * (*int* * *int*)

(*int* → *int*) → *int*

(*int* → *int*) * (*int* → *int*)

(*int* → *int*) → (*int* → (*int* * *int*))

Sintassi dei tipi

$\tau ::= int \mid \tau_0 * \tau_1 \mid \tau_0 \rightarrow \tau_1$

\mathcal{T}

insieme di tutti i tipi

assumiamo variabili tipate

$Ide = \{Ide_\tau\}_{\tau \in \mathcal{T}}$

$\hat{\cdot} : Ide \rightarrow \mathcal{T}$

\hat{x} denota il tipo di x

Type judgements

formula: $t : \tau$ si legge "ha tipo"

sono assegnati ai pre-termini
usando un insieme di regole di inferenza
(induzione strutturale sulla sintassi HOFL)

Sistema di tipi

$$\frac{}{x : \hat{x}} \quad \frac{}{n : int} \quad \frac{t_0 : int \quad t_1 : int}{t_0 \text{ op } t_1 : int} \quad \frac{t : int \quad t_0 : \tau \quad t_1 : \tau}{\text{if } t \text{ then } t_0 \text{ else } t_1 : \tau}$$

$$\frac{t_0 : \tau_0 \quad t_1 : \tau_1}{(t_0, t_1) : \tau_0 * \tau_1}$$

$$\frac{t : \tau_0 * \tau_1}{\text{fst}(t) : \tau_0}$$

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$$\frac{t_1 : \tau_0 \rightarrow \tau_1 \quad t_0 : \tau_0}{t_1 t_0 : \tau_1}$$

$$\frac{x : \tau \quad t : \tau}{\text{rec } x. t : \tau}$$

Termini ben formati

$$t ::= x \mid n \mid t_0 \text{ op } t_1 \mid \text{if } t \text{ then } t_0 \text{ else } t_1$$
$$\mid (t_0, t_1) \mid \text{fst}(t) \mid \text{snd}(t)$$
$$\mid \lambda x. t \mid t_0 t_1$$
$$\mid \text{rec } x. t$$
$$\tau ::= \text{int} \mid \tau_0 * \tau_1 \mid \tau_0 \rightarrow \tau_1$$

\mathcal{T} insieme di tutti i tipi

un pre-termine t e' ben formato se $\exists \tau \in \mathcal{T}. t : \tau$

cioe' se possiamo assegnargli un tipo
anche detto ben tipato o *tipabile*

T_τ insieme di tutti i termini ben formati di tipo τ

Controllo dei tipi Church Type Theory

Church Type Theory

- Le variabili sono etichettate con tipi (dichiarati)
- Deduciamo il tipo dei termini per induzione strutturale cioè **usando le regole di inferenza in modo bottom-up**

Esempio

$$\frac{}{x : \widehat{x}} \quad \frac{}{n : int} \quad \frac{t_0 : int \quad t_1 : int}{t_0 \text{ op } t_1 : int} \quad \frac{t : int \quad t_0 : \tau \quad t_1 : \tau}{\text{if } t \text{ then } t_0 \text{ else } t_1 : \tau}$$

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$fact \triangleq \text{rec } f : int \rightarrow int. \lambda x : int. \text{if } x \text{ then } 1 \text{ else } x \times (f (x - 1))$

$fact : int \rightarrow int$

Esempio

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$$f : int \rightarrow int \quad \lambda x. \text{if } x \text{ then } 1 \text{ else } (x \times (f(x - 1))) : int \rightarrow int$$

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Esempio

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Esempio

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scritto in maniera più semplice

$fact \stackrel{\text{def}}{=} \mathbf{rec} \underbrace{f}_{int \rightarrow int} . \lambda \underbrace{x}_{int} . \mathbf{if} x \mathbf{then} 1 \mathbf{else} x \times (f (x - 1))$

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scritto in maniera più semplice

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=

$fact \triangleq \mathbf{rec} f : int \rightarrow int. \lambda x : int. \mathbf{if} x \mathbf{then} 1 \mathbf{else} x \times (f (x - 1))$

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Esempio

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The diagram shows the following structure with brackets:

- A bracket under x and 1 is labeled int .
- A bracket under $f(x - 1)$ is labeled int .
- A larger bracket under the entire expression $x \times (f(x - 1))$ is labeled int .

Esempio

$fact \triangleq \mathbf{rec} f : int \rightarrow int. \lambda x : int. \mathbf{if} x \mathbf{then} 1 \mathbf{else} x \times (f (x - 1))$

scritto in maniera più semplice

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The diagram illustrates the type annotations for the factorial function. Brackets are used to group parts of the expression and indicate their types: f is of type $int \rightarrow int$, x is of type int , the condition x is of type int , the then branch 1 is of type int , the else branch x is of type int , the recursive call $f(x-1)$ is of type int , and the final multiplication result is of type int .

Esempio

$fact \triangleq \mathbf{rec} f : int \rightarrow int. \lambda x : int. \mathbf{if} x \mathbf{then} 1 \mathbf{else} x \times (f (x - 1))$

scritto in maniera più semplice

$fact \stackrel{\text{def}}{=} \mathbf{rec} \underbrace{f}_{int \rightarrow int} . \lambda \underbrace{x}_{int} . \mathbf{if} \underbrace{x}_{int} \mathbf{then} \underbrace{1}_{int} \mathbf{else} \underbrace{x}_{int} \times \left(\underbrace{f}_{int \rightarrow int} \left(\underbrace{x}_{int} - \underbrace{1}_{int} \right) \right)$

Esempio

$fact \triangleq \mathbf{rec} f : int \rightarrow int. \lambda x : int. \mathbf{if} x \mathbf{then} 1 \mathbf{else} x \times (f (x - 1))$

scritto in maniera più semplice

$fact \stackrel{\text{def}}{=} \mathbf{rec} \underbrace{f}_{int \rightarrow int} . \lambda \underbrace{x}_{int} . \mathbf{if} \underbrace{x}_{int} \mathbf{then} \underbrace{1}_{int} \mathbf{else} \underbrace{x}_{int} \times \left(\underbrace{f}_{int \rightarrow int} \left(\underbrace{\underbrace{x}_{int} - \underbrace{1}_{int}}_{int} \right) \right)$

Esempio

$fact \triangleq \mathbf{rec} f : int \rightarrow int. \lambda x : int. \mathbf{if} x \mathbf{then} 1 \mathbf{else} x \times (f (x - 1))$

scritto in maniera più semplice

$fact \stackrel{\text{def}}{=} \mathbf{rec} \underbrace{f}_{int \rightarrow int} . \lambda \underbrace{x}_{int} . \mathbf{if} \underbrace{x}_{int} \mathbf{then} \underbrace{1}_{int} \mathbf{else} \underbrace{x}_{int} \times \left(\underbrace{f}_{int \rightarrow int} \left(\underbrace{x}_{int} - \underbrace{1}_{int} \right) \right) : int \rightarrow int$

Inferenza di tipi

Curry Type theory

Curry Type Theory

- I tipi delle variabili non devono necessariamente essere (dichiarati)
- Inferiamo il tipo dei termini e' inferito usando le regole per derivare vincoli di tipo (equazioni di tipo) le cui soluzioni (tramite unificazione) definiscono il tipo principale

Esempio

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2)))$$

intuitivamente

$$t \ 0 \equiv (0, (t \ 2)) \equiv (0, (2, (t \ 4))) \equiv \dots \equiv (0, (2, (4, \dots)))$$

sequenza di tutti i numeri pari

possiamo digitare sequenze di interi di lunghezza fissa
non abbiamo un tipo per sequenze di lunghezza qualsiasi/infinita

Esempio

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2)))$$

Haskell

```
Prelude> let p x = (x, p (x+2))
```

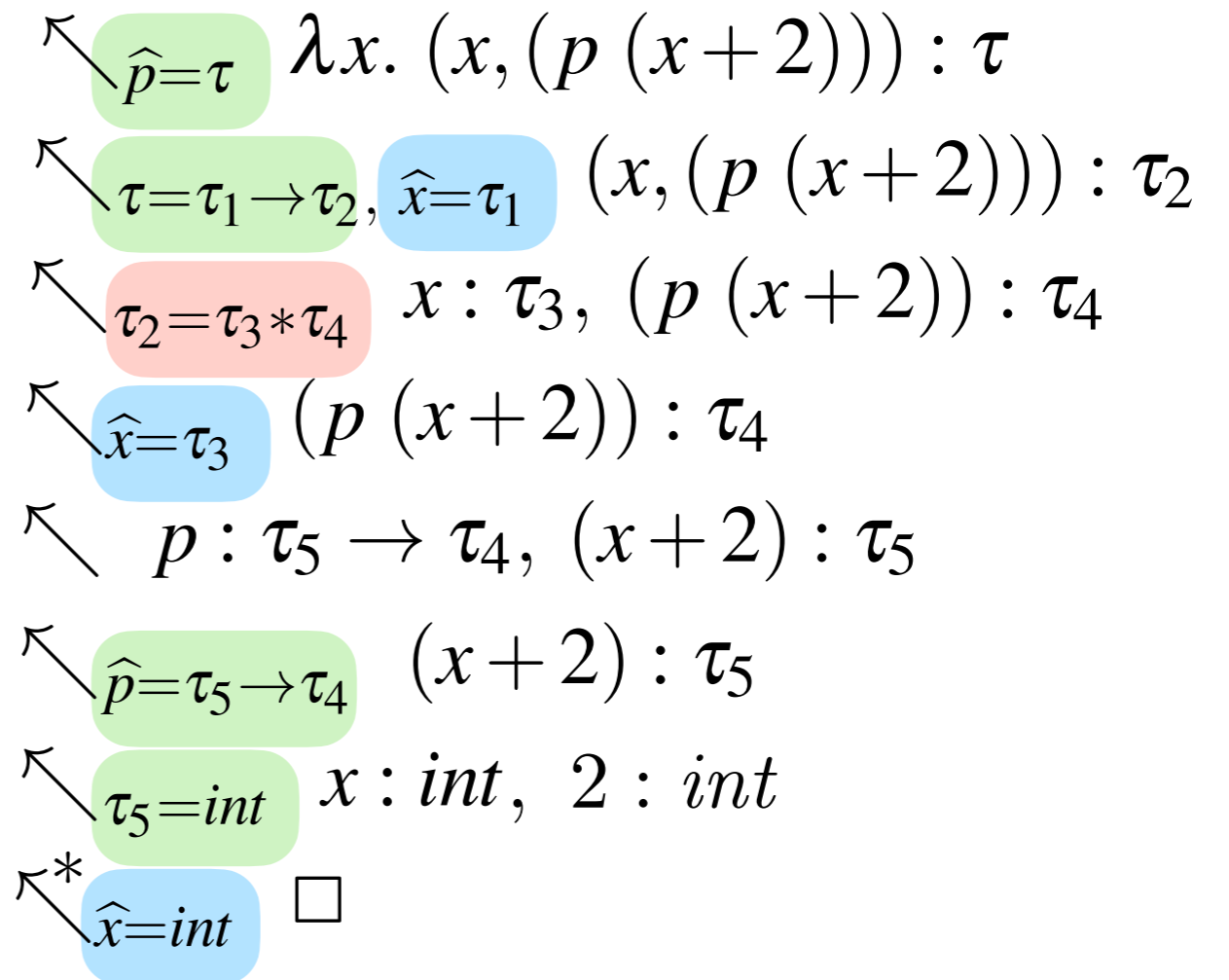
```
<interactive>:...:5: error:
```

- Occurs check: cannot construct the infinite type: b ~ (t, b)
Expected type: t -> b
Actual type: t -> (t, b)
- Relevant bindings include
p :: t -> b (bound at <interactive>:...:5)

Esempio

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2)))$$

$$t = \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2))) : \tau$$



$$\left. \begin{array}{l} \hat{x} = \tau_1 \\ \hat{x} = \tau_3 \\ \hat{x} = int \end{array} \right\} \tau_1 = \tau_3 = int$$

$$\left. \begin{array}{l} \hat{p} = \tau = \tau_1 \rightarrow \tau_2 \\ \hat{p} = \tau_5 \rightarrow \tau_4 \end{array} \right\} \begin{array}{l} \tau_1 = \tau_5 = int \\ \tau_2 = \tau_4 \end{array}$$

$$\left. \begin{array}{l} \tau_2 = \tau_4 \\ \tau_2 = \tau_3 * \tau_4 \end{array} \right\} \text{fail! (occur check)}$$

$$\frac{}{x : \hat{x}} \quad \frac{}{n : int} \quad \frac{t_0 : int \quad t_1 : int}{t_0 \text{ op } t_1 : int} \quad \frac{t : int \quad t_0 : \tau \quad t_1 : \tau}{\text{if } t \text{ then } t_0 \text{ else } t_1 : \tau}$$

$$\frac{t_0 : \tau_0 \quad t_1 : \tau_1}{(t_0, t_1) : \tau_0 * \tau_1} \quad \frac{t : \tau_0 * \tau_1}{\mathbf{fst}(t) : \tau_0} \quad \frac{t : \tau_0 * \tau_1}{\mathbf{snd}(t) : \tau_1}$$

$$\frac{x : \tau_0 \quad t : \tau_1}{\lambda x. t : \tau_0 \rightarrow \tau_1} \quad \frac{t_1 : \tau_0 \rightarrow \tau_1 \quad t_0 : \tau_0}{t_1 \ t_0 : \tau_1}$$

$$\frac{x : \tau \quad t : \tau}{\mathbf{rec} \ x. t : \tau}$$

Esempio

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2)))$$

scritto in maniera più semplice

$$t = \mathbf{rec} \ p. \ \lambda \underline{x}. \ (\underline{x}, (p \ (x + \underline{2})))$$

\square
int

Esempio

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2)))$$

scritto in maniera più semplice

$$t = \mathbf{rec} \ p. \ \lambda \underline{x}. \ (\underline{x}, (p \ (\underline{x} + \underline{2})))$$

int *int*

Esempio

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2)))$$

scritto in maniera più semplice

$$t = \mathbf{rec} \ p. \ \lambda \underset{\text{int}}{\underline{x}}. \ (\underset{\text{int}}{\underline{x}}, (p \ (\underset{\text{int}}{\underline{x}} + \underset{\text{int}}{\underline{2}})))$$

Esempio

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2)))$$

scritto in maniera più semplice

$$t = \mathbf{rec} \ p. \ \lambda \underset{\substack{\square \\ int}}{x}. \ (\underset{\substack{\square \\ int}}{x}, \ (\ p \ (\underset{\substack{\square \\ int}}{x} + \underset{\substack{\square \\ int}}{2})))$$

$\underbrace{\hspace{10em}}_{int}$

Esempio

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2)))$$

scritto in maniera più semplice

$$t = \mathbf{rec} \ p. \ \lambda \underset{\text{int}}{\underline{x}}. \ \left(\underset{\text{int}}{\underline{x}}, \left(\underset{\text{int} \rightarrow \tau_4}{\underline{p}} \left(\underset{\text{int}}{\underline{x}} + \underset{\text{int}}{\underline{2}} \right) \right) \right)$$

$\underbrace{\hspace{10em}}_{\text{int}}$

Esempio

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2)))$$

scritto in maniera più semplice

$$t = \mathbf{rec}_{\substack{\square \\ int \rightarrow \tau_4}} \ p. \ \lambda_{\substack{\square \\ int}} \ x. \ (\underbrace{x}_{int}, (\underbrace{p}_{\substack{\square \\ int \rightarrow \tau_4}} (\underbrace{x}_{int} + \underbrace{2}_{int}))))$$

Esempio

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2)))$$

scritto in maniera più semplice

$$t = \mathbf{rec} \underset{\text{int} \rightarrow \tau_4}{\underbrace{p}}. \ \lambda \underset{\text{int}}{\underbrace{x}}. \ (\underset{\text{int}}{\underbrace{x}}, (\underset{\text{int} \rightarrow \tau_4}{\underbrace{p}} \ (\underset{\text{int}}{\underbrace{x}} + \underset{\text{int}}{\underbrace{2}})))$$

τ_4

Esempio

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2)))$$

scritto in maniera più semplice

$$t = \mathbf{rec} \underset{\text{int} \rightarrow \tau_4}{\boxed{p}}. \ \lambda \underset{\text{int}}{\boxed{x}}. \ (\underset{\text{int}}{\boxed{x}}, (\underset{\text{int} \rightarrow \tau_4}{\boxed{p}} \ (\underbrace{\underset{\text{int}}{\boxed{x}} + \underset{\text{int}}{\boxed{2}}}_{\text{int}})))$$

τ_4

$\text{int} * \tau_4$

Esempio

$$t \stackrel{\text{def}}{=} \mathbf{rec} \ p. \ \lambda x. \ (x, (p \ (x + 2)))$$

scritto in maniera più semplice

$$t = \mathbf{rec} \ p. \ \lambda \underset{\substack{\square \\ int}}{x}. \ \left(\underset{\substack{\square \\ int}}{x}, \left(\underset{\substack{\square \\ int \rightarrow \tau_4}}{p} \left(\underset{\substack{\square \\ int}}{x} + \underset{\substack{\square \\ int}}{2} \right) \right) \right)$$

$$(int \rightarrow (int * \tau_4))$$

$$(int \rightarrow (int * \tau_4)) = (int \rightarrow \tau_4) \Rightarrow \tau_4 = (int * \tau_4) \quad \text{fallisce (occur check)}$$

Esercizio

inferire il tipo del termine sotto

$$\frac{}{x : \widehat{x}} \quad \frac{}{n : int} \quad \frac{t_0 : int \quad t_1 : int}{t_0 \text{ op } t_1 : int} \quad \frac{t : int \quad t_0 : \tau \quad t_1 : \tau}{\text{if } t \text{ then } t_0 \text{ else } t_1 : \tau}$$

$$\frac{t_0 : \tau_0 \quad t_1 : \tau_1}{(t_0, t_1) : \tau_0 * \tau_1} \quad \frac{t : \tau_0 * \tau_1}{\text{fst}(t) : \tau_0} \quad \frac{t : \tau_0 * \tau_1}{\text{snd}(t) : \tau_1}$$

$$\frac{x : \tau_0 \quad t : \tau_1}{\lambda x. t : \tau_0 \rightarrow \tau_1} \quad \frac{t_1 : \tau_0 \rightarrow \tau_1 \quad t_0 : \tau_0}{t_1 t_0 : \tau_1}$$

$$\frac{x : \tau \quad t : \tau}{\text{rec } x. t : \tau}$$

$\text{rec } rep. \lambda n. \lambda f. \lambda x. \text{if } \underline{n} \text{ then } \underline{x}$
 $\text{else } \underline{f} (\underline{rep} (\underline{n - 1}) \underline{f} \underline{x})$

$$\tau_0 \rightarrow \tau_0$$

$$(\tau_1 \rightarrow \tau_0) \rightarrow \tau_0 \rightarrow \tau_0$$

$$int \rightarrow (\tau_1 \rightarrow \tau_0) \rightarrow \tau_0 \rightarrow \tau_0$$

$$n : int$$

$$x : \tau_0$$

$$f : \tau_1 \rightarrow \tau_0$$

$$rep : int \rightarrow (\tau_1 \rightarrow \tau_0) \rightarrow \tau_0 \rightarrow \tau_1 \quad \rightarrow \quad \tau_0 = \tau_1$$

Esercizio

$$\frac{}{x : \widehat{x}} \quad \frac{}{n : int} \quad \frac{t_0 : int \quad t_1 : int}{t_0 \text{ op } t_1 : int} \quad \frac{t : int \quad t_0 : \tau \quad t_1 : \tau}{\text{if } t \text{ then } t_0 \text{ else } t_1 : \tau}$$

$$\frac{t_0 : \tau_0 \quad t_1 : \tau_1}{(t_0, t_1) : \tau_0 * \tau_1}$$

$$\frac{t : \tau_0 * \tau_1}{\text{fst}(t) : \tau_0}$$

$$\frac{t : \tau_0 * \tau_1}{\text{snd}(t) : \tau_1}$$

inferire il tipo del termine sotto

$$\frac{x : \tau_0 \quad t : \tau_1}{\lambda x. t : \tau_0 \rightarrow \tau_1}$$

$$\frac{t_1 : \tau_0 \rightarrow \tau_1 \quad t_0 : \tau_0}{t_1 t_0 : \tau_1}$$

$$\frac{x : \tau \quad t : \tau}{\text{rec } x. t : \tau}$$

$$\lambda x. \left(\left(\begin{array}{l} \text{rec } f. \lambda y. \text{if } \underline{(x - y)} \text{ then } \underline{0} \\ \text{else if } \underline{(x + y)} \text{ then } \underline{1} \\ \text{else } \underline{f (y + 1)} \end{array} \right) \underline{0} \right)$$

$$int \rightarrow int$$

$$int$$

$$int \rightarrow int$$

int

x : int

y : int

f : int → int

Substituzioni capture-avoiding (di nuovo)

Variabili libere

$$\text{fv}(n) \stackrel{\text{def}}{=} \emptyset$$

$$\text{fv}(x) \stackrel{\text{def}}{=} \{x\}$$

$$\text{fv}(t_0 \text{ op } t_1) \stackrel{\text{def}}{=} \text{fv}(t_0) \cup \text{fv}(t_1)$$

$$\text{fv}(\mathbf{if } t \mathbf{ then } t_0 \mathbf{ else } t_1) \stackrel{\text{def}}{=} \text{fv}(t) \cup \text{fv}(t_0) \cup \text{fv}(t_1)$$

$$\text{fv}((t_0, t_1)) \stackrel{\text{def}}{=} \text{fv}(t_0) \cup \text{fv}(t_1)$$

$$\text{fv}(\mathbf{fst}(t)) \stackrel{\text{def}}{=} \text{fv}(t)$$

$$\text{fv}(\mathbf{snd}(t)) \stackrel{\text{def}}{=} \text{fv}(t)$$

$$\text{fv}(\lambda x. t) \stackrel{\text{def}}{=} \text{fv}(t) \setminus \{x\}$$

$$\text{fv}((t_0 \ t_1)) \stackrel{\text{def}}{=} \text{fv}(t_0) \cup \text{fv}(t_1)$$

$$\text{fv}(\mathbf{rec } x. t) \stackrel{\text{def}}{=} \text{fv}(t) \setminus \{x\}$$

Substituzioni

$$n[t/x] = n$$

$$y[t/x] \stackrel{\text{def}}{=} \begin{cases} t & \text{if } y = x \\ y & \text{if } y \neq x \end{cases}$$

$$(t_0 \text{ op } t_1)[t/x] \stackrel{\text{def}}{=} t_0[t/x] \text{ op } t_1[t/x] \quad \text{with op} \in \{+, -, \times\}$$

$$(\text{if } t' \text{ then } t_0 \text{ else } t_1)[t/x] \stackrel{\text{def}}{=} \text{if } t'[t/x] \text{ then } t_0[t/x] \text{ else } t_1[t/x]$$

$$(t_0, t_1)[t/x] \stackrel{\text{def}}{=} (t_0[t/x], t_1[t/x])$$

$$\text{fst}(t')[t/x] \stackrel{\text{def}}{=} \text{fst}(t'[t/x])$$

$$\text{snd}(t')[t/x] \stackrel{\text{def}}{=} \text{snd}(t'[t/x])$$

$$(t_0 t_1)[t/x] \stackrel{\text{def}}{=} (t_0[t/x] t_1[t/x])$$

$$(\lambda y. t')[t/x] \stackrel{\text{def}}{=} \lambda z. (t'[z/y][t/x]) \quad \text{for } z \notin \text{fv}(\lambda y. t') \cup \text{fv}(t) \cup \{x\}$$

$$(\text{rec } y. t')[t/x] \stackrel{\text{def}}{=} \text{rec } z. (t'[z/y][t/x]) \quad \text{for } z \notin \text{fv}(\text{rec } y. t') \cup \text{fv}(t) \cup \{x\}$$

I tipi sono rispettati

$$\text{TH. } \begin{array}{l} x_0 : \tau_0 \\ t_0 : \tau_0 \end{array} \quad t : \tau \quad \Rightarrow \quad t^{[t_0 / x_0]} : \tau$$

prova omessa
(per induzione strutturale)